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Model-based Test Case Generation for (Dynamic) Software Product Lines

The 6th Meeting on Feature-oriented Software Development
2014, May 04-07 – Schloss Dagstuhl, Germany



DFG Priority Programme 1593
Design for Future – Managed Software Evolution

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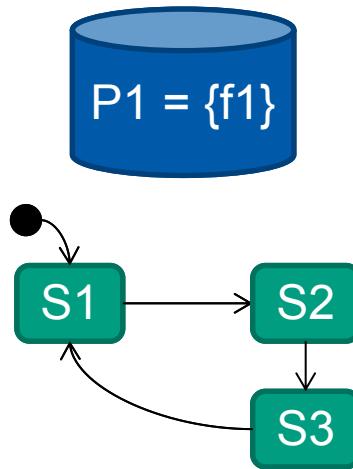
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Testing Software Product Lines



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Products



Test Models

Test Suites

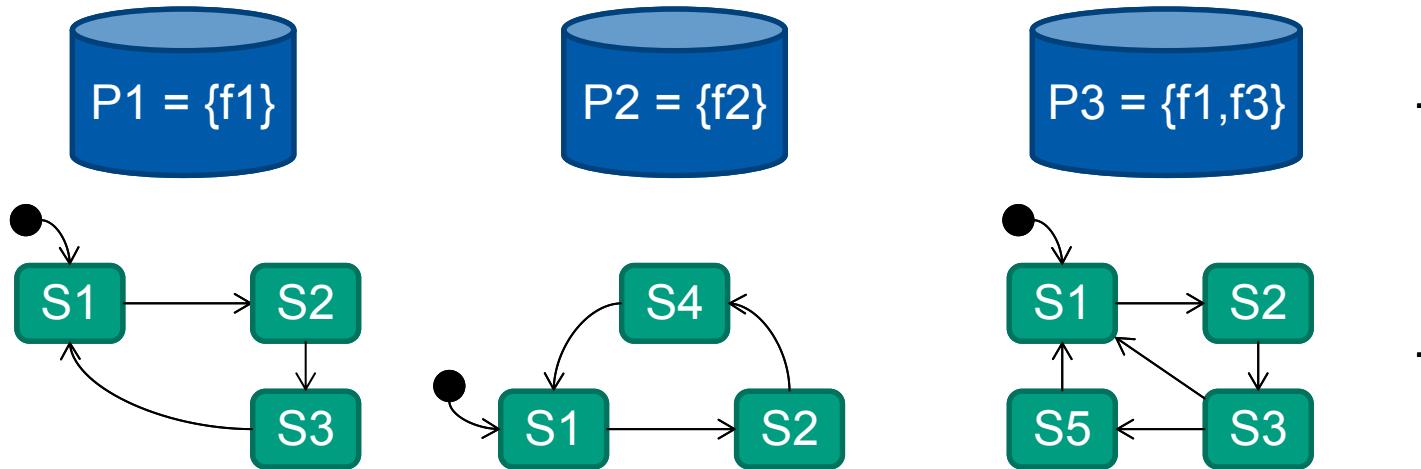


[Cichos et al. 2011]

Testing Software Product Lines



Products



Test Models

Generate test cases for every product individually is inefficient!

Test Suites

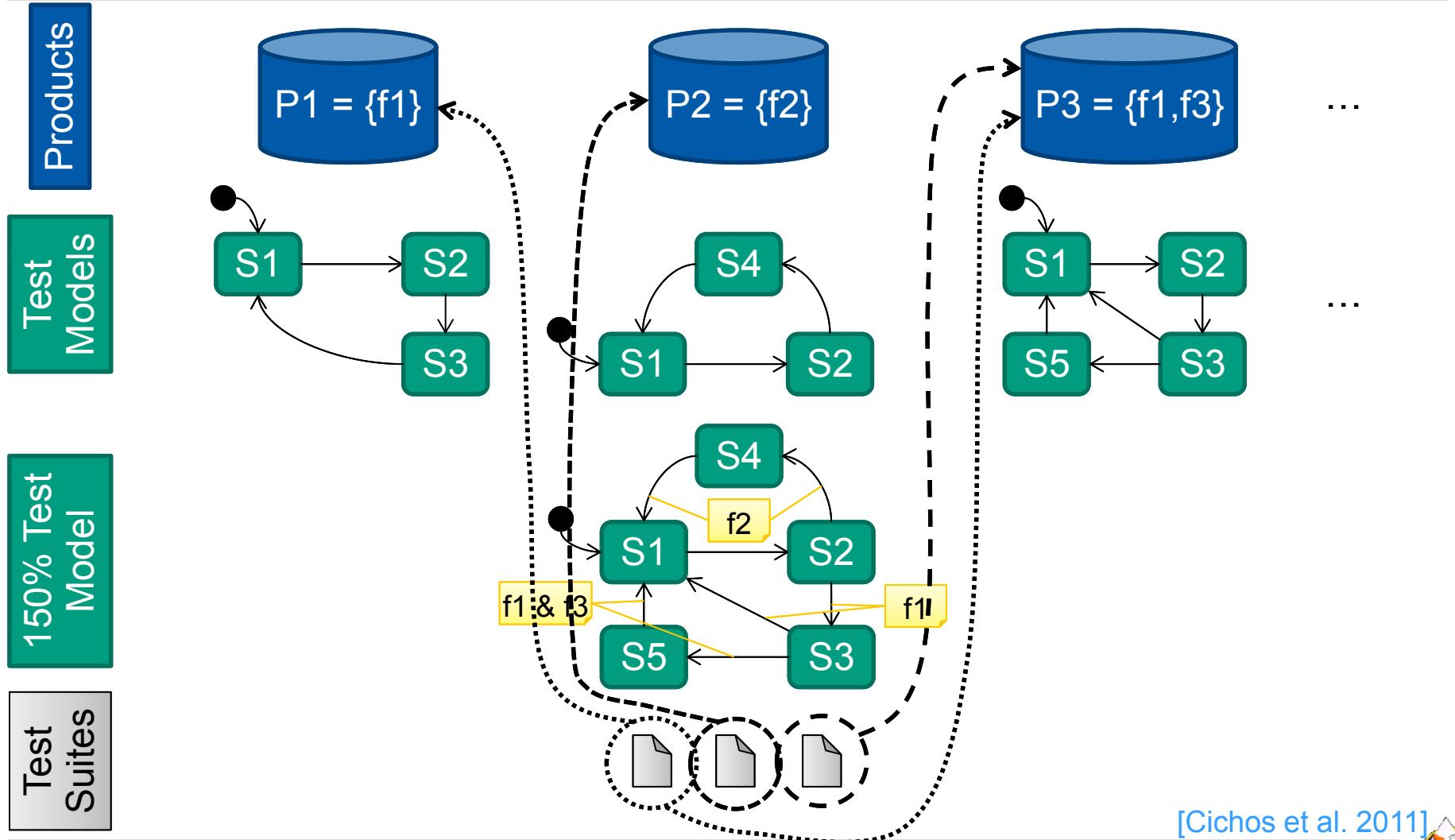


[Cichos et al. 2011]

Testing Software Product Lines



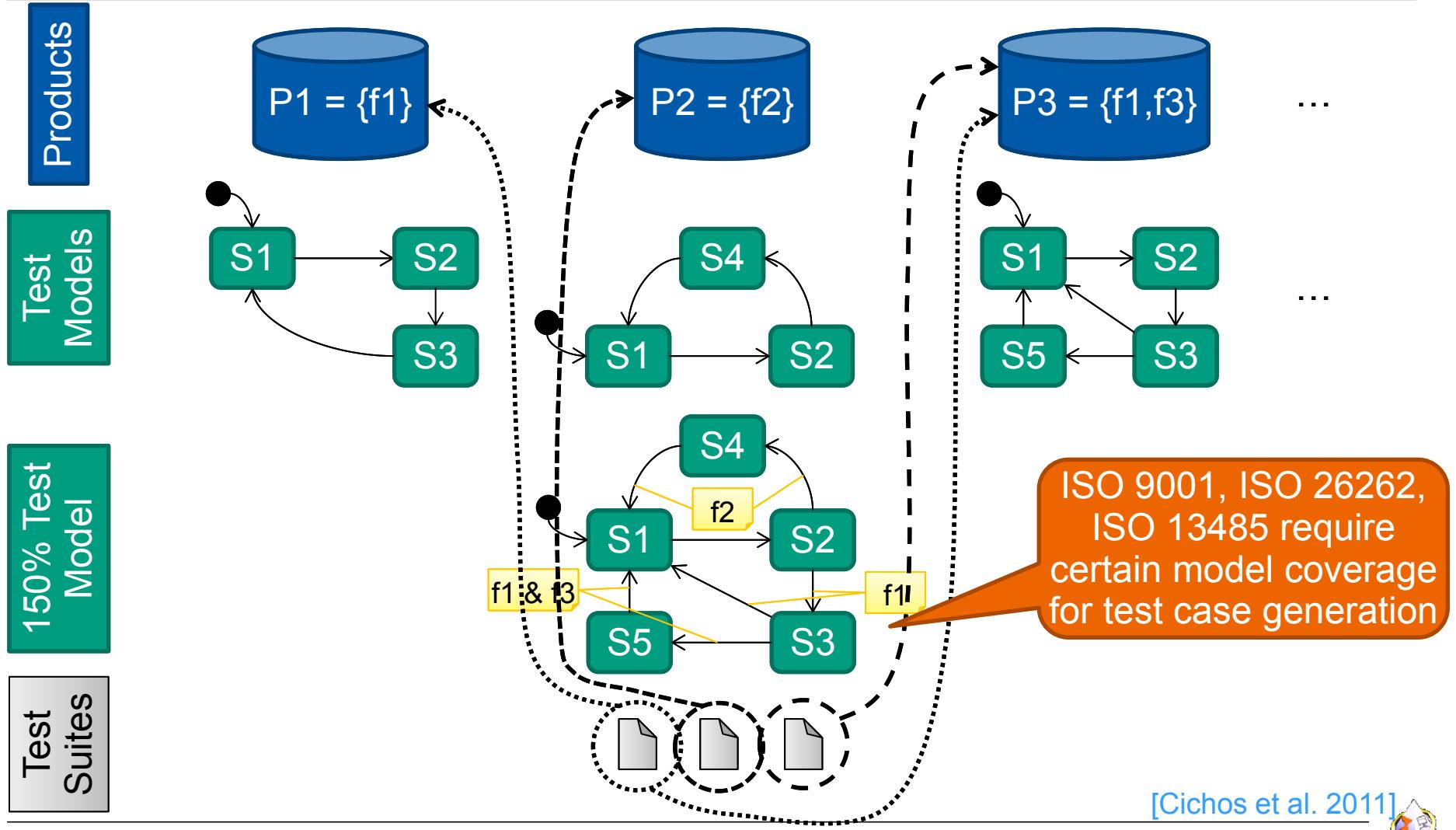
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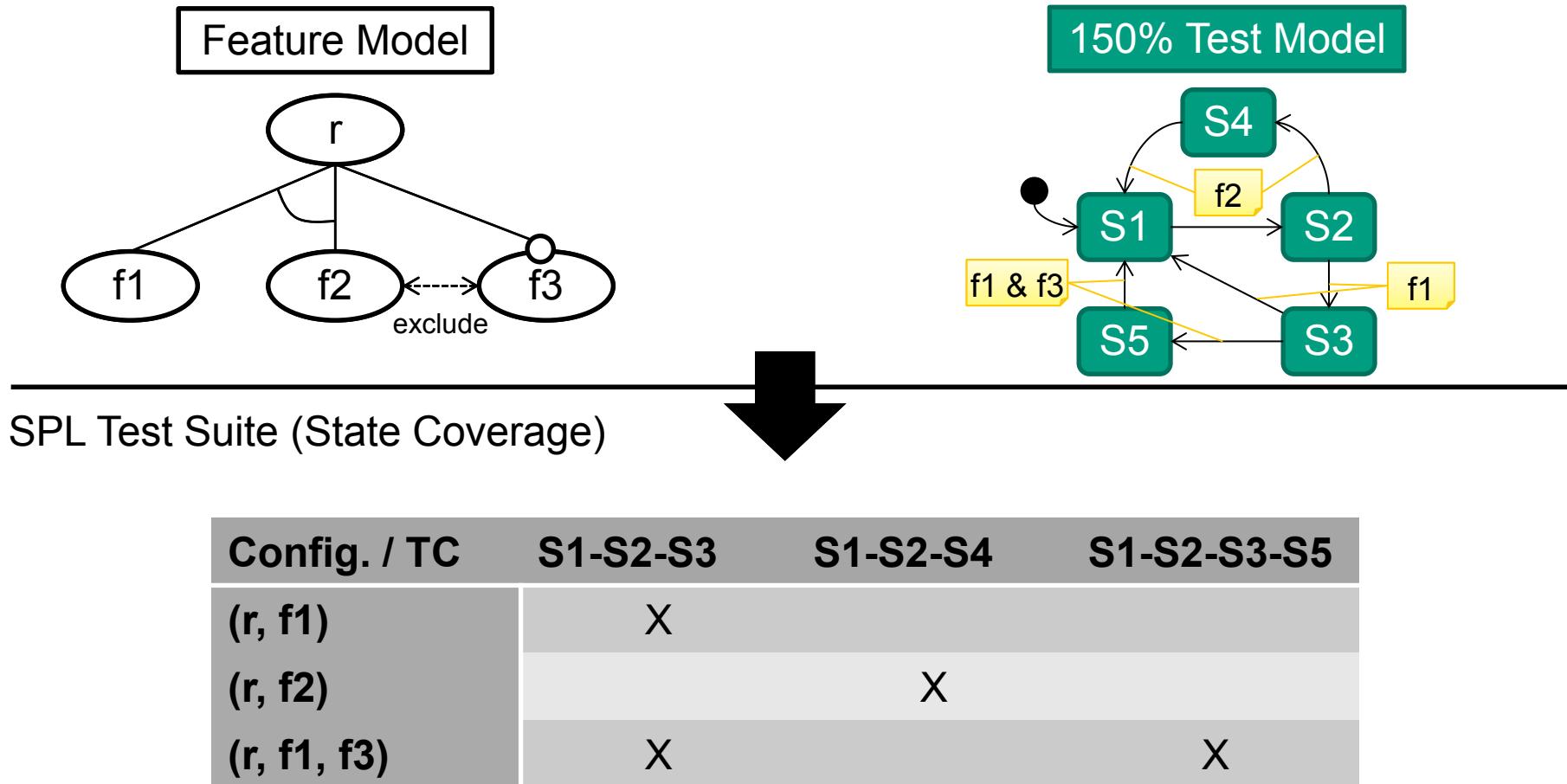
Testing Software Product Lines



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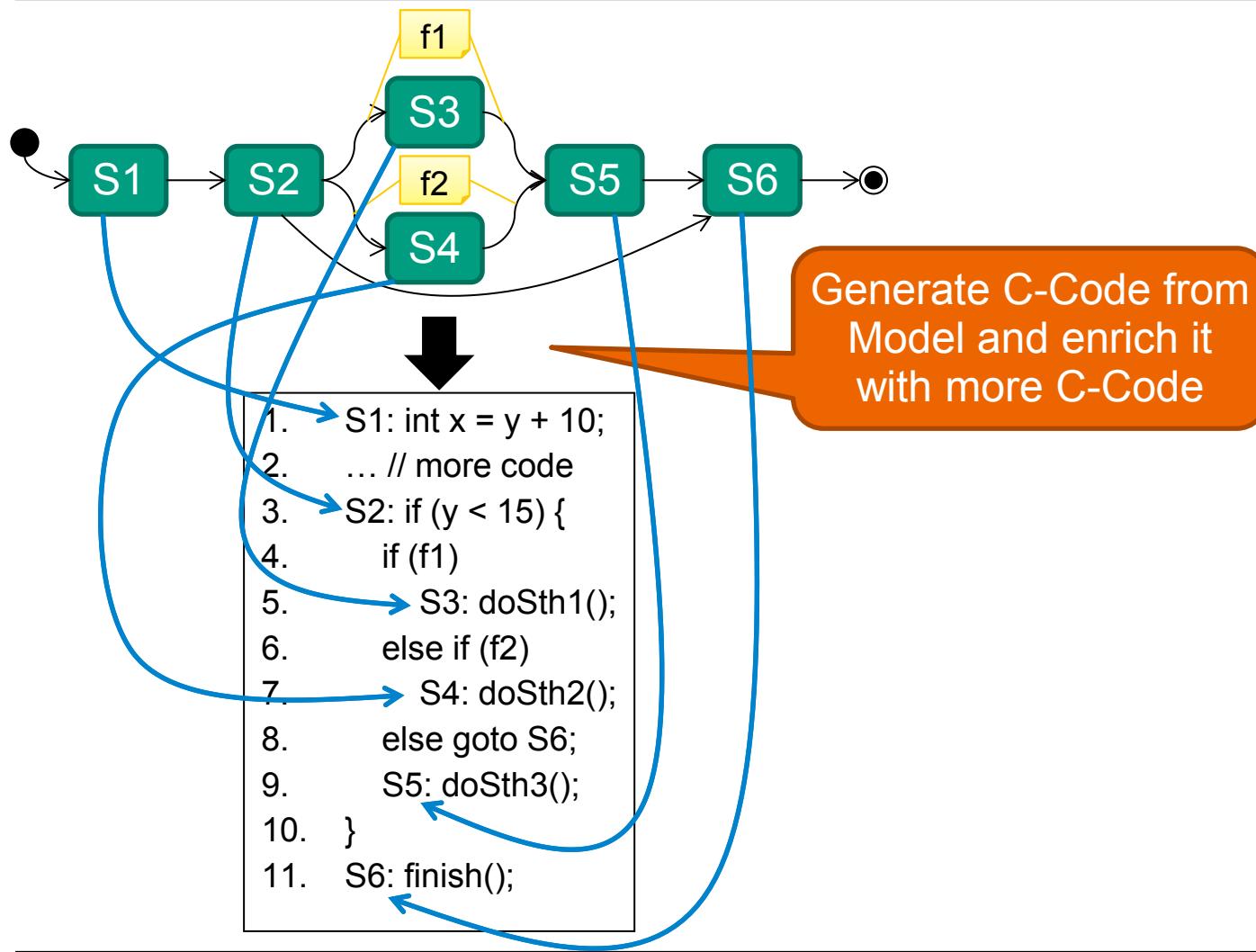


SPL Test Suite Generation

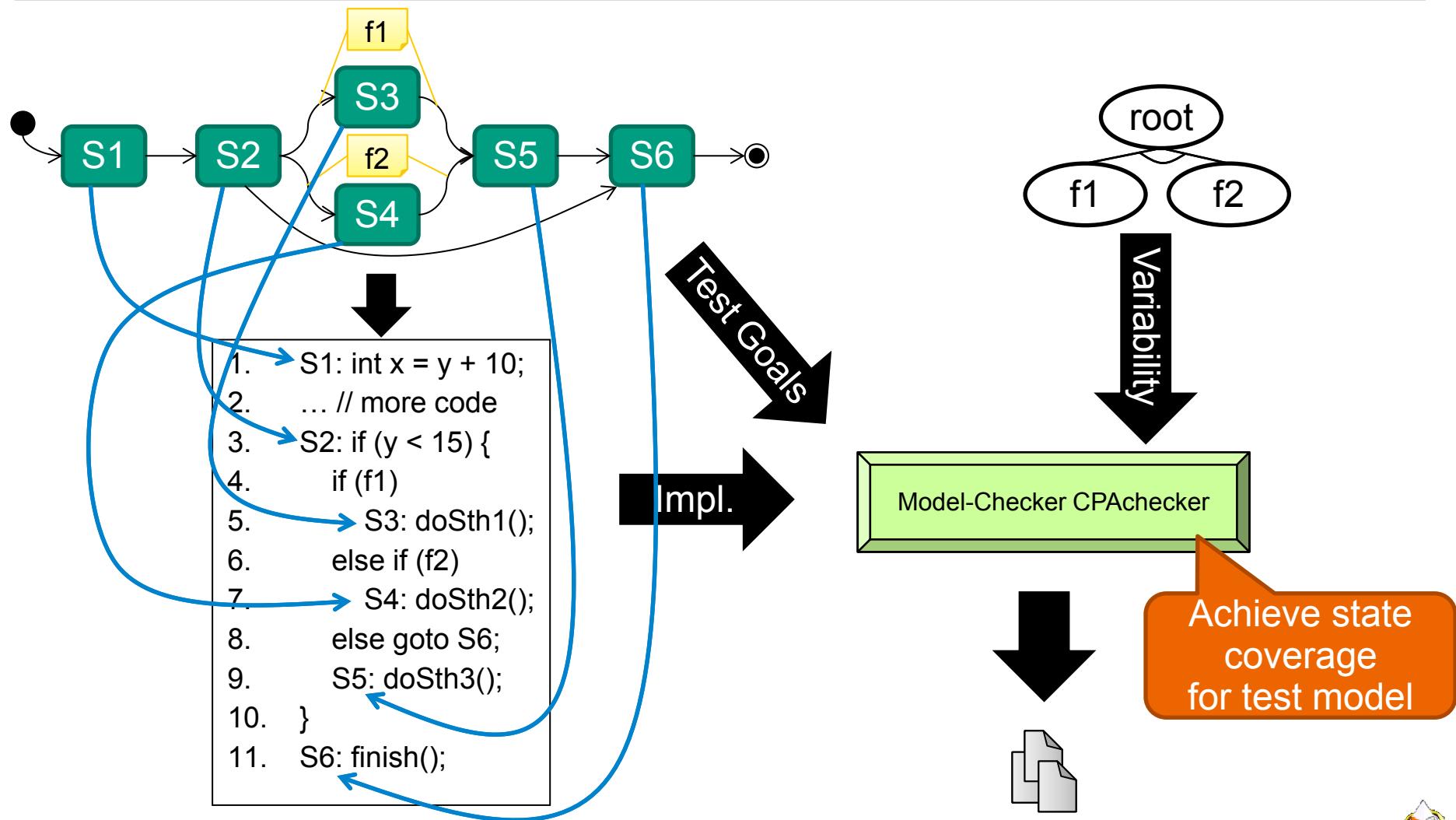


[Cichos et al. 2011], [Cichos et al. 2012], [Lochau et al. 2014]

Using Symbolic Model Checking for Test Case Generation



Using Symbolic Model Checking for Test Case Generation

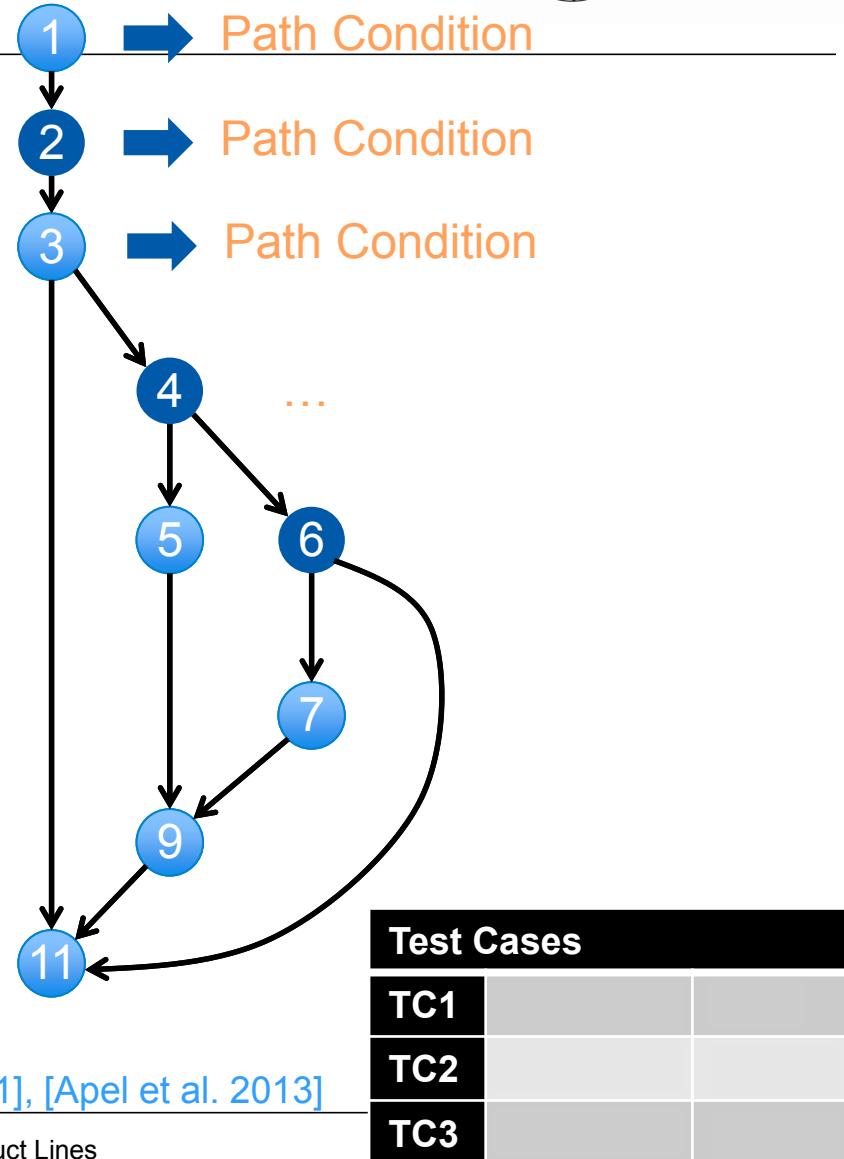


Test Case Derivation from Counterexamples

```

1. S1: int x = y + 10; ← Assignment
2. ... // more code
3. S2: if (y < 15) ← Assumption
4.   if (f1)
5.     S3: doSth1();
6.   else if (f2)
7.     S4: doSth2();
8.   else goto S6;
9.   S5: doSth3();
10. }
11. S6: finish();
  
```

	(r,f1)	(r,f2)
S1		
S2		
S3		
S4		
S5		
S6		



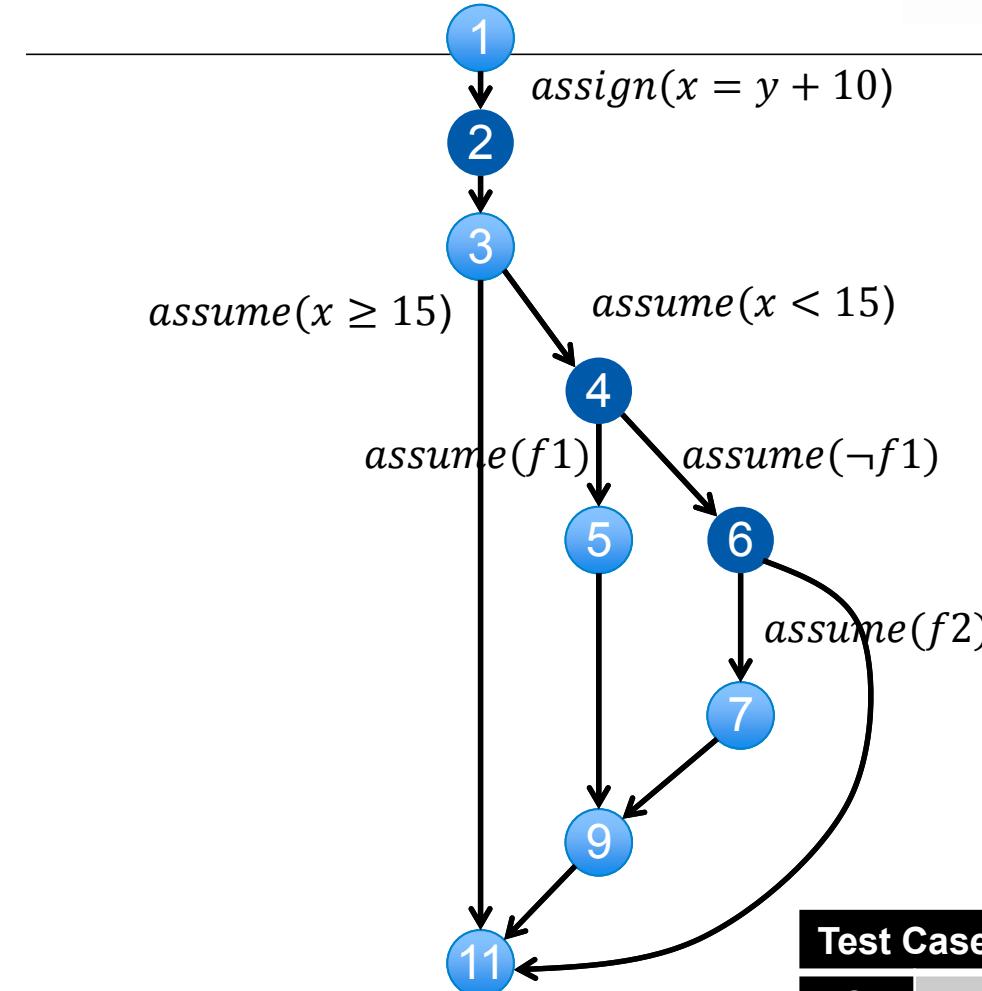
[Beyer et al. 2004], [Rhein et al. 2011], [Apel et al. 2013]

Test Case Derivation from Counterexamples

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```

	(r,f1)	(r,f2)
S1		
S2		
S3		
S4		
S5		
S6		



Test Cases		
TC1		
TC2		
TC3		

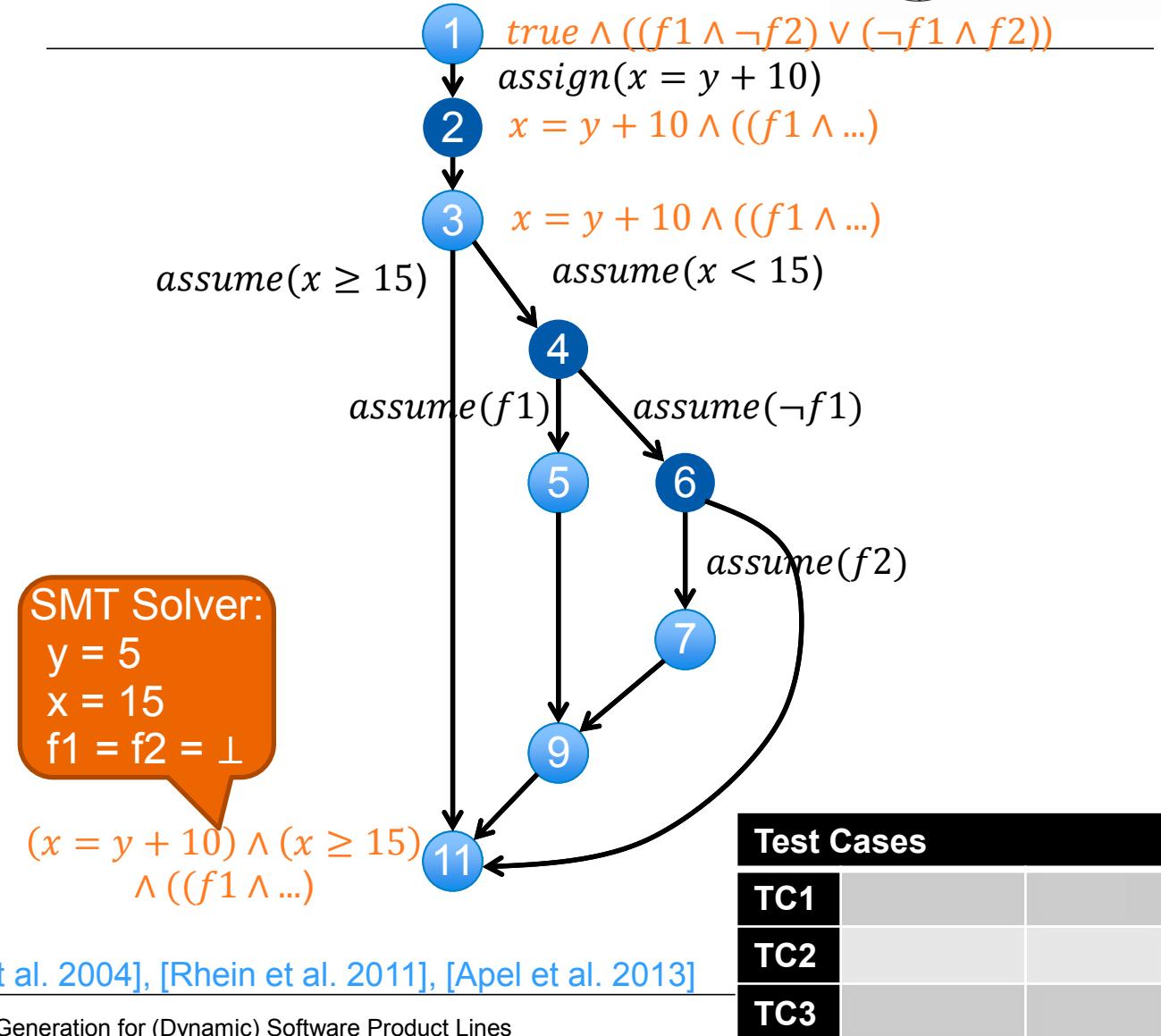
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	(r,f1)	(r,f2)
S1		
S2		
S3		
S4		
S5		
S6		

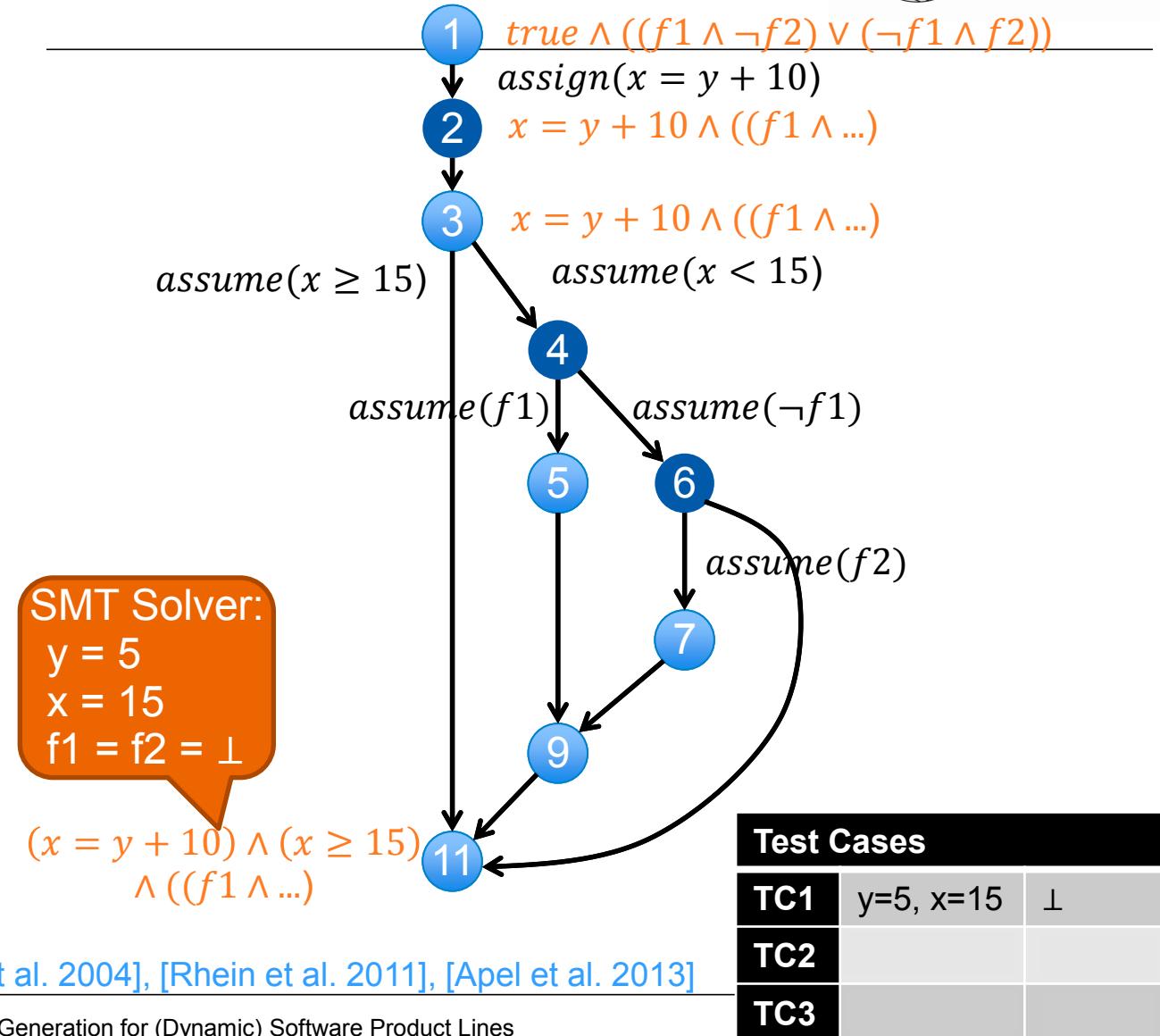


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	(r,f1)	(r,f2)
S1	TC1	TC1
S2	TC1	TC1
S3		
S4		
S5		
S6	TC1	TC1



Test Case Derivation from Counterexamples

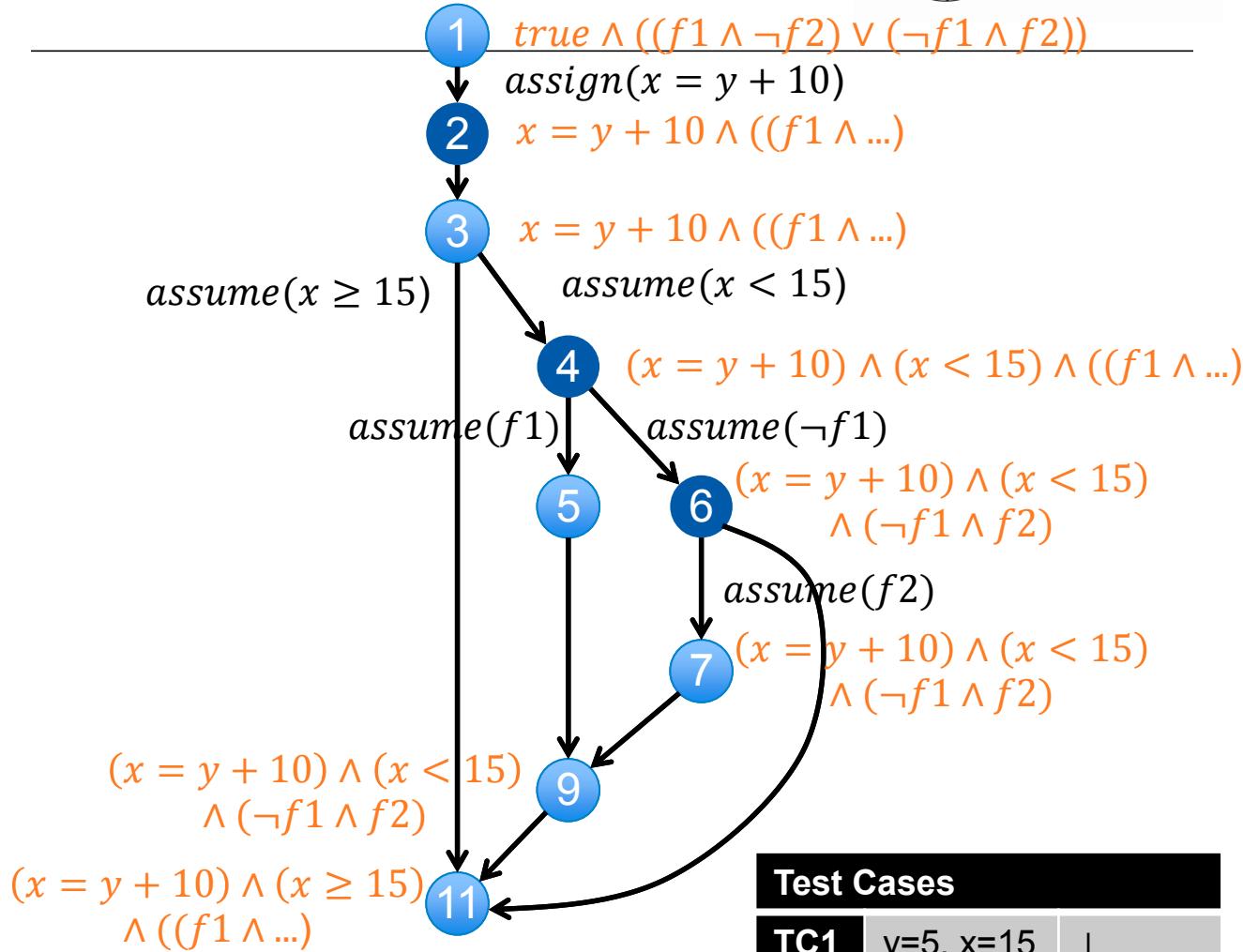


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	(r,f1)	(r,f2)
S1	TC1	TC1
S2	TC1	TC1
S3		
S4		
S5		
S6	TC1	TC1



[Beyer et al. 2004], [Rhein et al. 2011], [Apel et al. 2013]

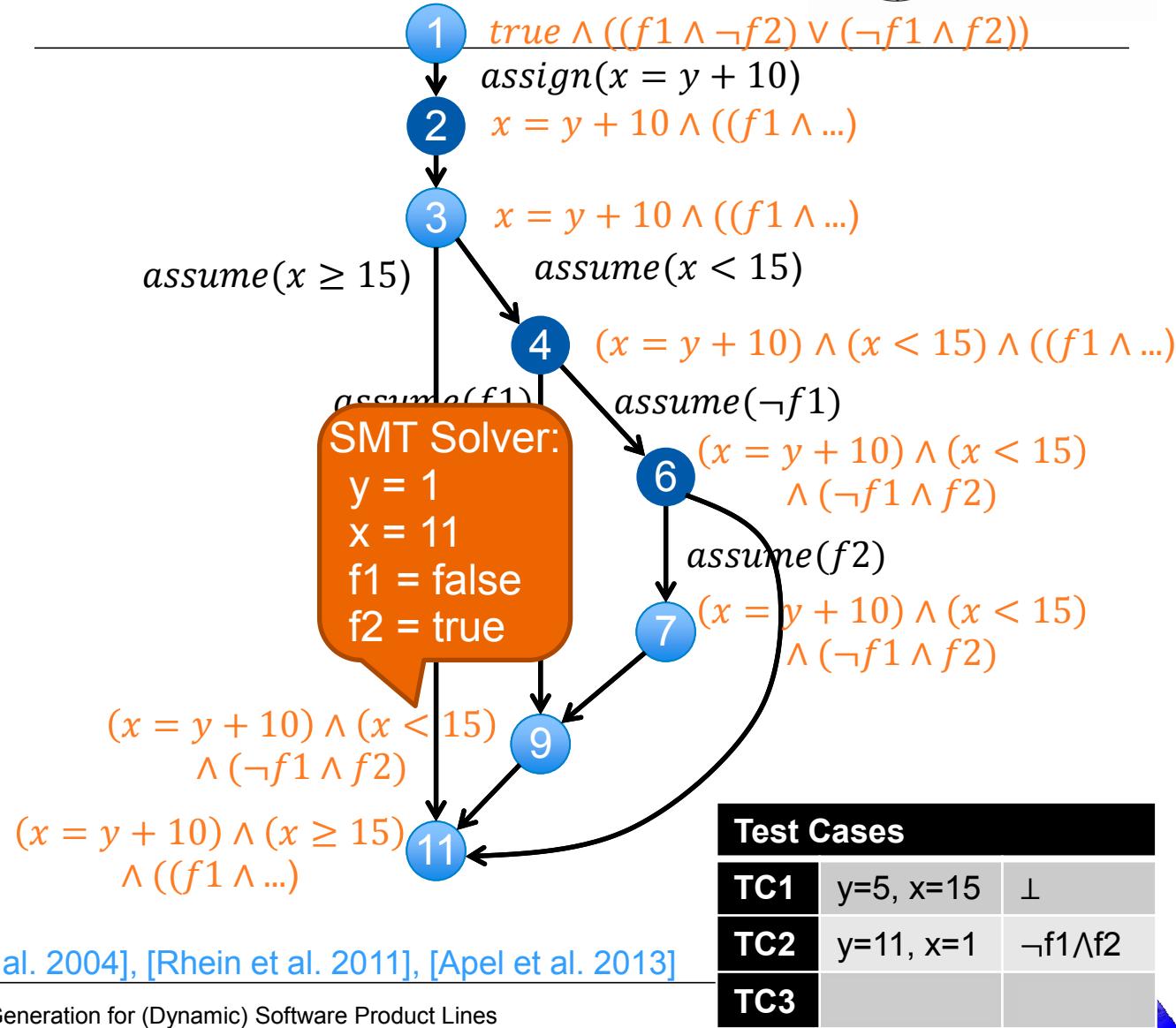
Test Cases		
TC1	y=5, x=15	⊥
TC2		
TC3		

Test Case Derivation from Counterexamples

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1. S1: int x = y + 10;
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	(r,f1)	(r,f2)
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S2	TC1	TC1
S3		
S4		TC2
S5		TC2
S6	TC1	TC1



[Beyer et al. 2004], [Rhein et al. 2011], [Apel et al. 2013]

Test Case Derivation from Counterexamples

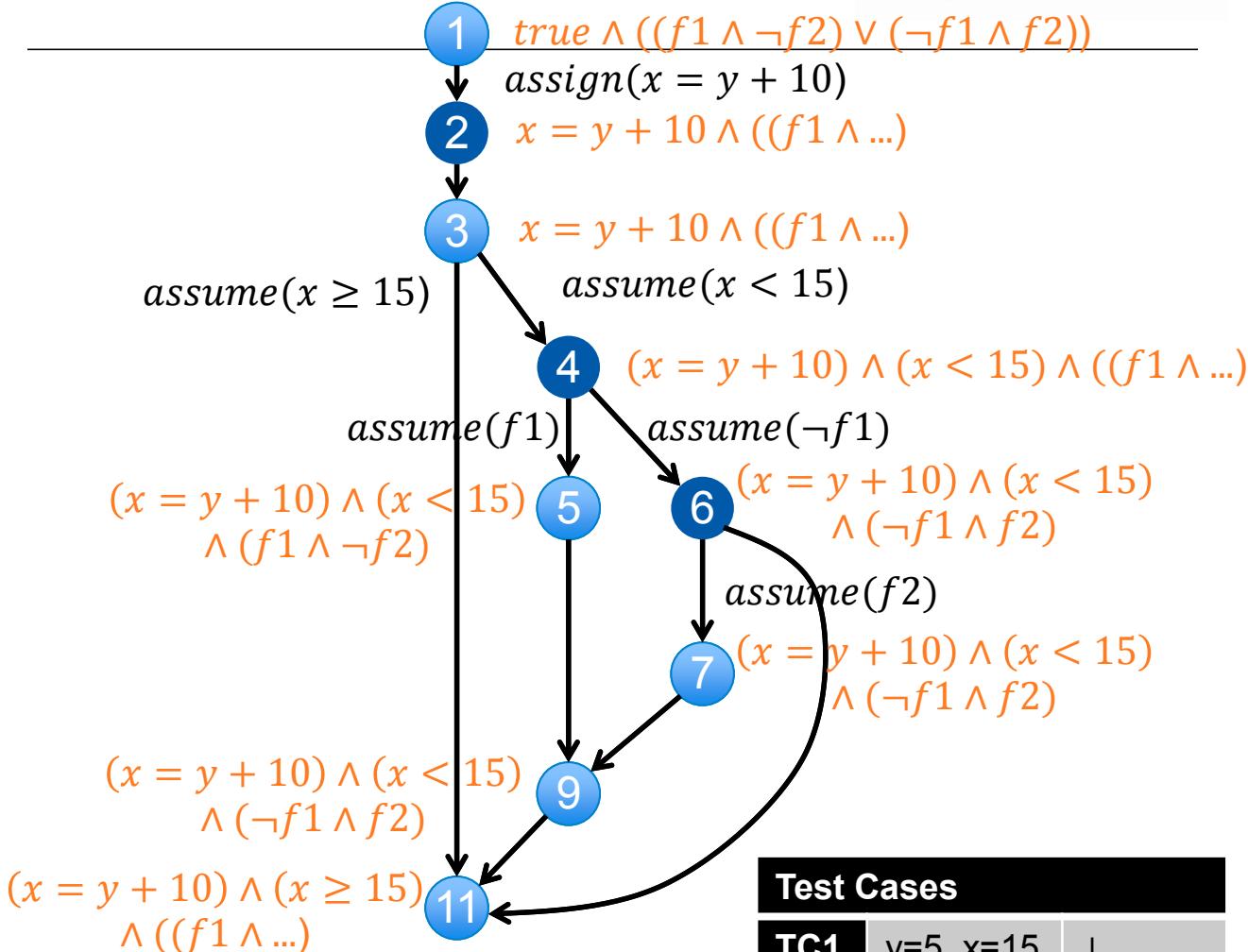


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S1	TC1	TC1
S2	TC1	TC1
S3		
S4		TC2
S5		TC2
S6	TC1	TC1



[Beyer et al. 2004], [Rhein et al. 2011], [Apel et al. 2013]

Test Cases		
TC1	y=5, x=15	\perp
TC2	y=11, x=1	$\neg f1 \wedge f2$
TC3		

Test Case Derivation from Counterexamples

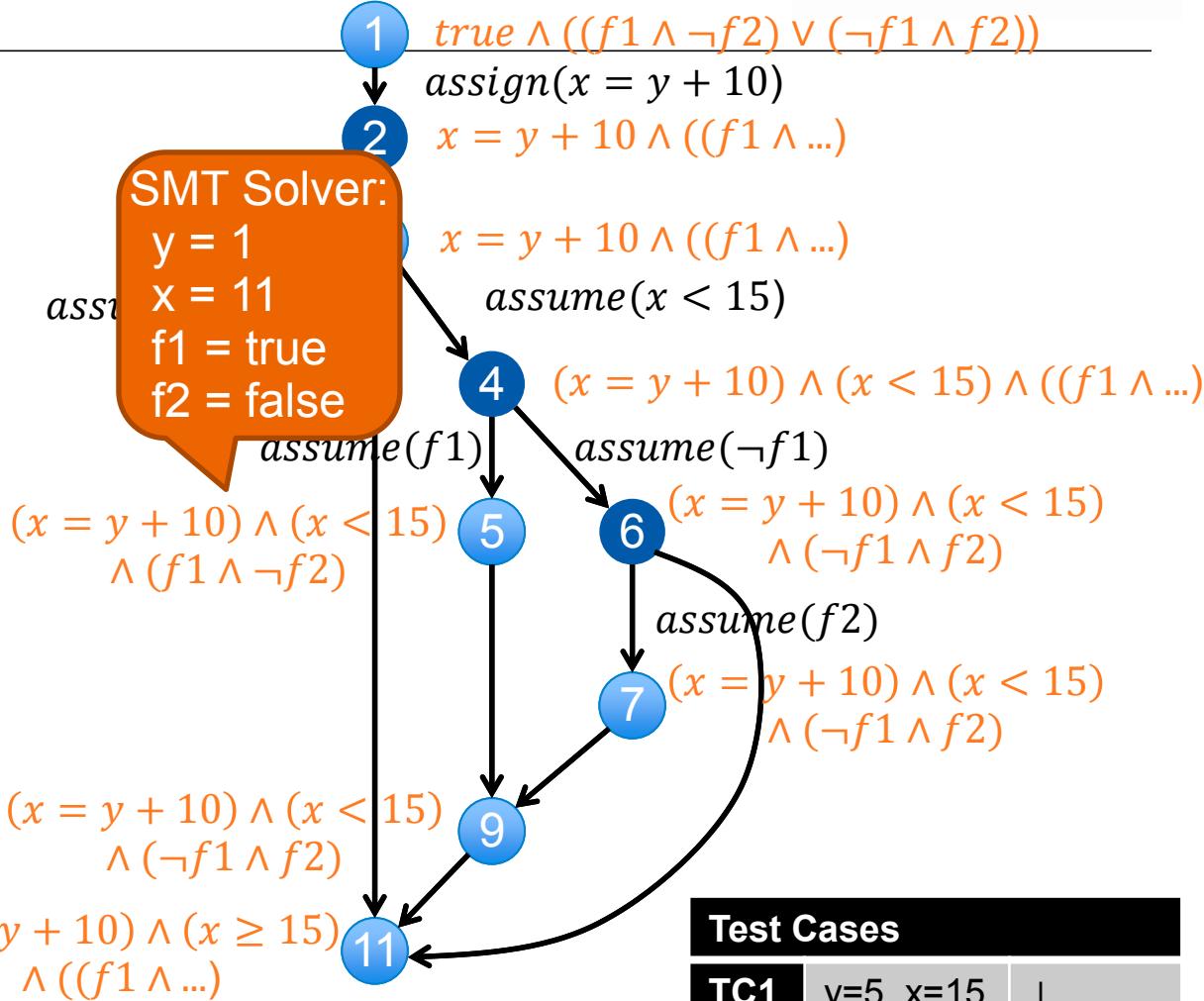


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	(r,f1)	(r,f2)
S1	TC1	TC1
S2	TC1	TC1
S3	TC3	
S4		TC2
S5		TC2
S6	TC1	TC1



Test Cases		
TC1	$y=5, x=15$	\perp
TC2	$y=11, x=1$	$\neg f1 \wedge f2$
TC3	$y=11, x=1$	$f1 \wedge \neg f2$

Test Case Derivation from Counterexamples

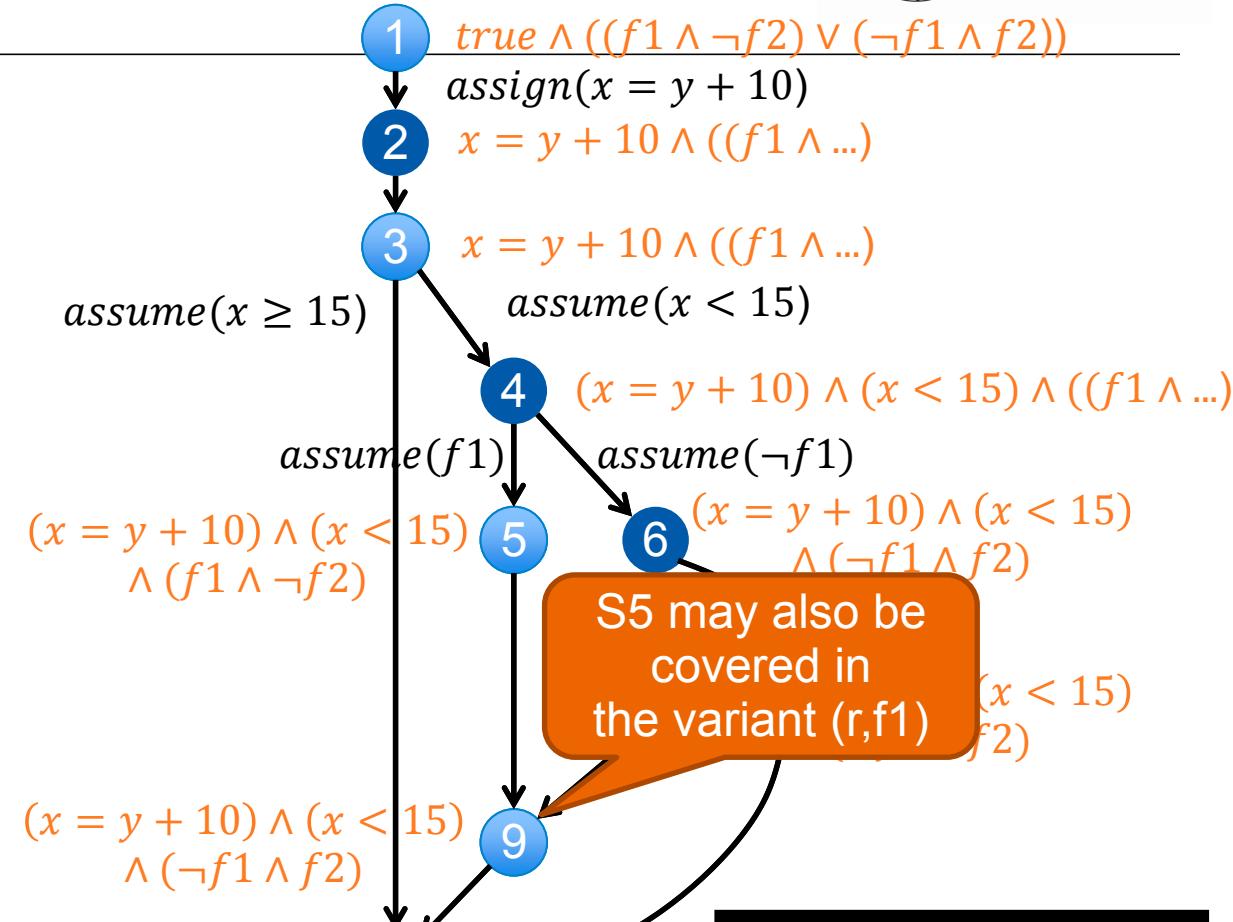


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	(r,f1)	(r,f2)
S1	TC1	TC1
S2	TC1	TC1
S3	TC3	
S4		TC2
S5		TC2
S6	TC1	TC1



A BDD analysis checks if a location may be reached by a path with another feature configuration.

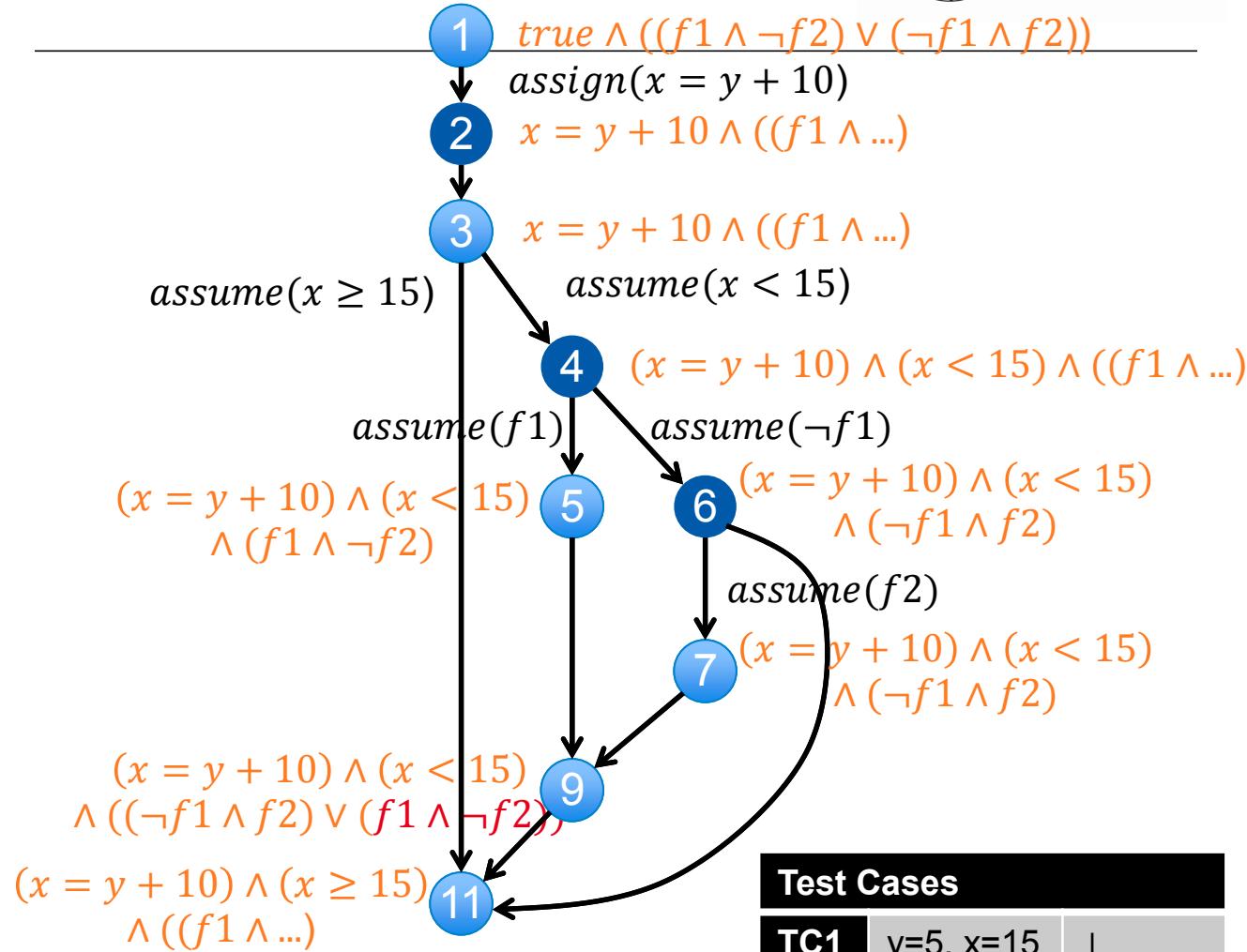
[Beyer]

Test Cases		
TC1	y=5, x=15	\perp
TC2	y=11, x=1	$\neg f1 \wedge f2$
TC3	y=11, x=1	$f1 \wedge \neg f2$

Test Case Derivation from Counterexamples

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	(r,f1)	(r,f2)
S1	TC1	TC1
S2	TC1	TC1
S3	TC3	
S4		TC2
S5		TC2
S6	TC1	TC1
10	5/8/2014	Model-p



[Beyer et al. 2004], [Rhein et al. 2011], [Apel et al. 2013]

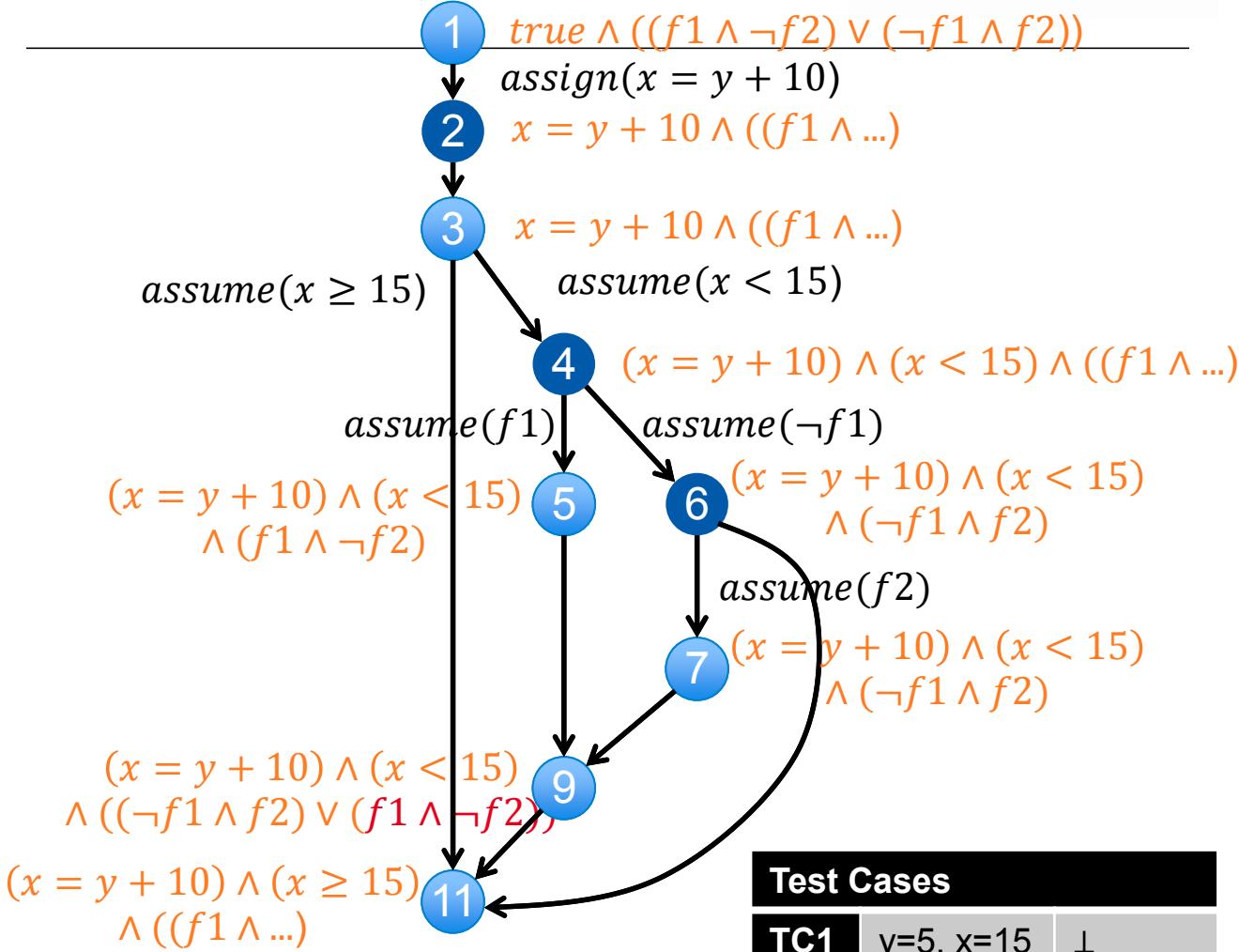
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	(r,f1)	(r,f2)
S1	TC1	TC1
S2	TC1	TC1
S3	TC3	
S4		TC2
S5	TC3	TC2
S6	TC1	TC1



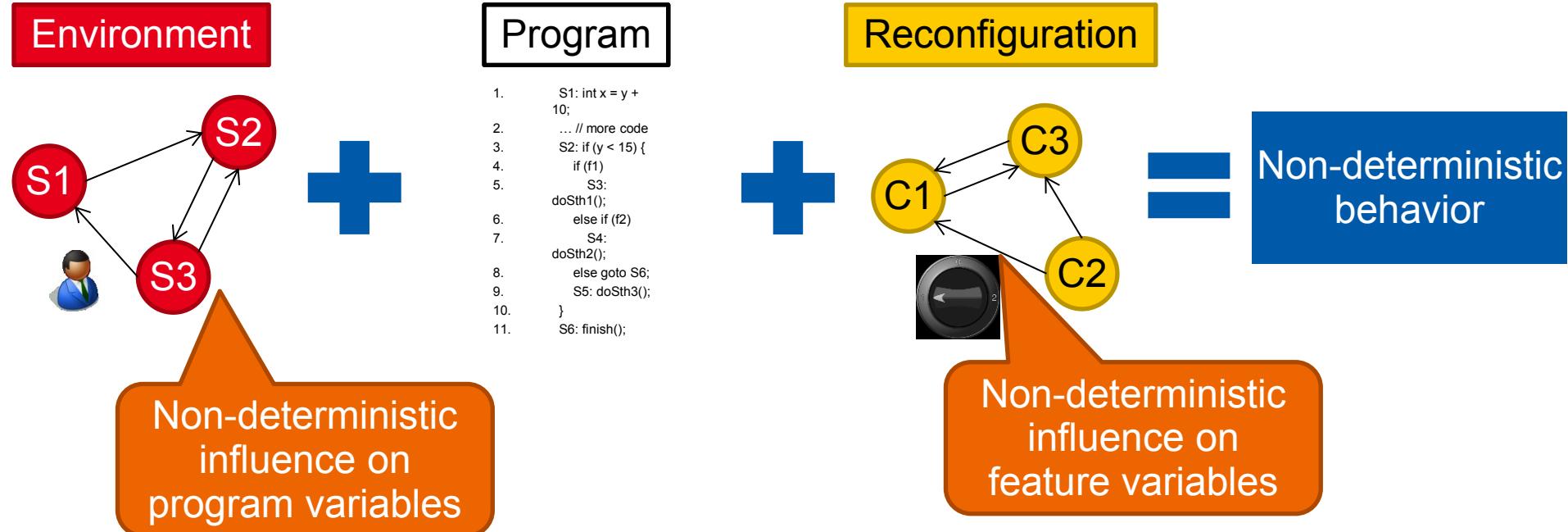
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TC3	y=11, x=1	\perp

Testing Dynamic Software Product Lines

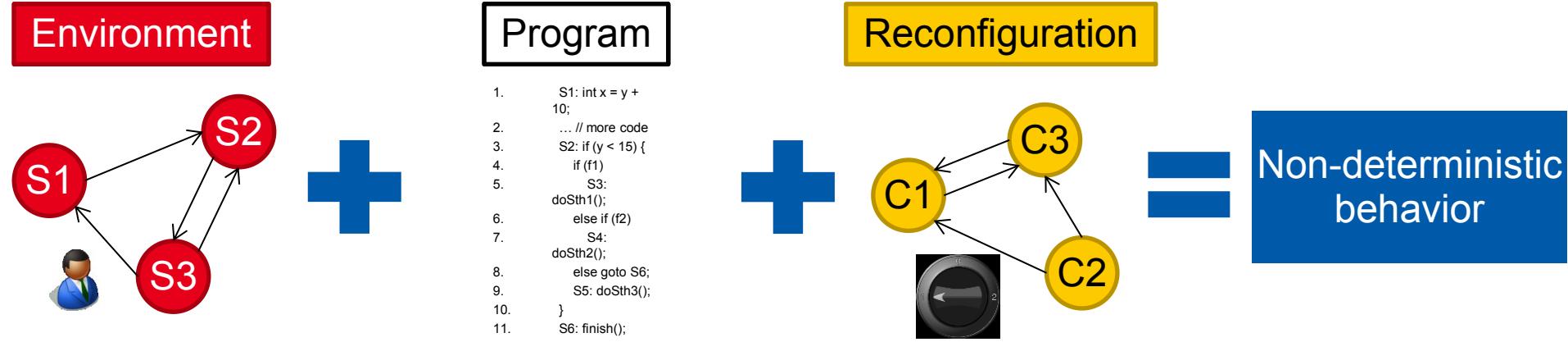


- **Testing SPLs:** test that every derivable variant is correctly implemented
- **Testing DSPLs:** test that all possible reconfigurations are correct



Testing Dynamic Software Product Lines

- **Testing SPLs:** test that every derivable variant is correctly implemented
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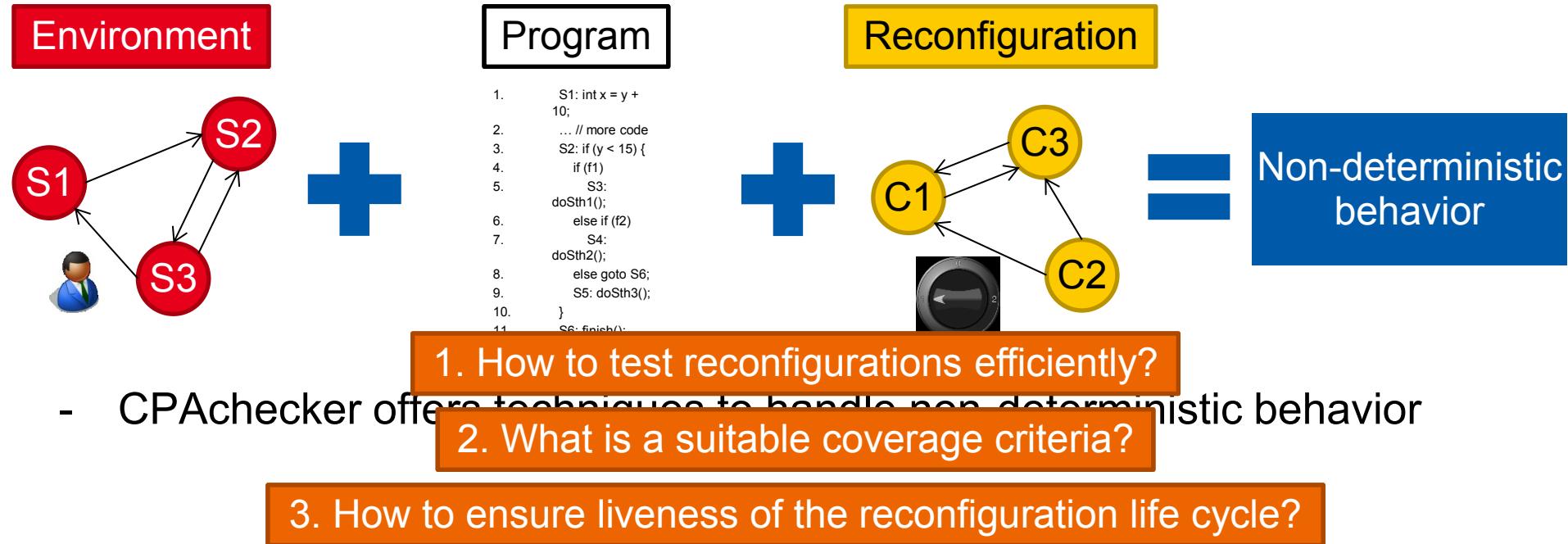


- CPAchecker offers techniques to handle non-deterministic behavior

Testing Dynamic Software Product Lines



- **Testing SPLs:** test that every derivable variant is correctly implemented
- **Testing DSPLs:** test that all possible reconfigurations are correct



List of Literature



- [Cichos et al. 2011] H. Cichos, S. Oster, M. Lochau, A. Schürr : *Model-Based Coverage-Driven Test Suite Generation for Software Product Lines*. MoDELS 2011: 425-439
- [Cichos et al. 2012] H. Cichos, M. Lochau, S. Oster, A. Schürr : *Reduktion von Testsuiten für Software-Produktlinien*. Software Engineering 2012: 143-154
- [Lochau et al. 2014] M. Lochau, J. Bürdek, S. Lity, M. Hagner, C. Legat, U. Goltz, A. Schürr : *Applying Model-based Software Product Line Testing Approaches to the Automation Engineering Domain*. AT Journal 2014 (submitted)
- [Holzer et al. 2011] A. Holzer, M. Tautschnig, C. Schallhart, H. Veith : *An Introduction to Test Specification in FQL*. In S. Barner, I. Harris, D. Kroening, and O. Raz, editors, *Hardware and Software: Verification and Testing*, volume 6504 of *Lecture Notes in Computer Science*, pages 9–22. Springer Berlin Heidelberg, 2011.
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- [Holzmann 1997] G. J. Holzmann : *The Model Checker Spin*, IEEE Trans. on Software Engineering, Vol. 23, No. 5, May 1997, pp. 279-295.



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- [Beyer et al. 2004] D. Beyer, A. J. Chlipala, T. A. Henzinger, R. Jhala, and R. Majumdar. *Generating Test from Counterexamples*. In *Proc. of the 26th ICSE*, pages 326–335, 2004.
- [Rhein et al. 2011] A. von Rhein, S. Apel, F. Raimondi : *Introducing Binary Decision Diagrams in the Explicit-State Verification of Java Code*. Java Pathfinder Workshop (co-located with ASE'11), 2011.
- [Apel et al. 2013] S. Apel, A. von Rhein, P. Wendler, A. Größlinger, D. Beyer : *Strategies for Product-Line Verification: Case Studies and Experiments*. In Proceedings of the IEEE/ACM International Conference on Software Engineering (ICSE), pages 482–491. IEEE Computer Society, 2013.