

Model-Based Design  
with  
**SCADE Suite®**



# Training Resources

Day 2

# The SCADE Suite workflow

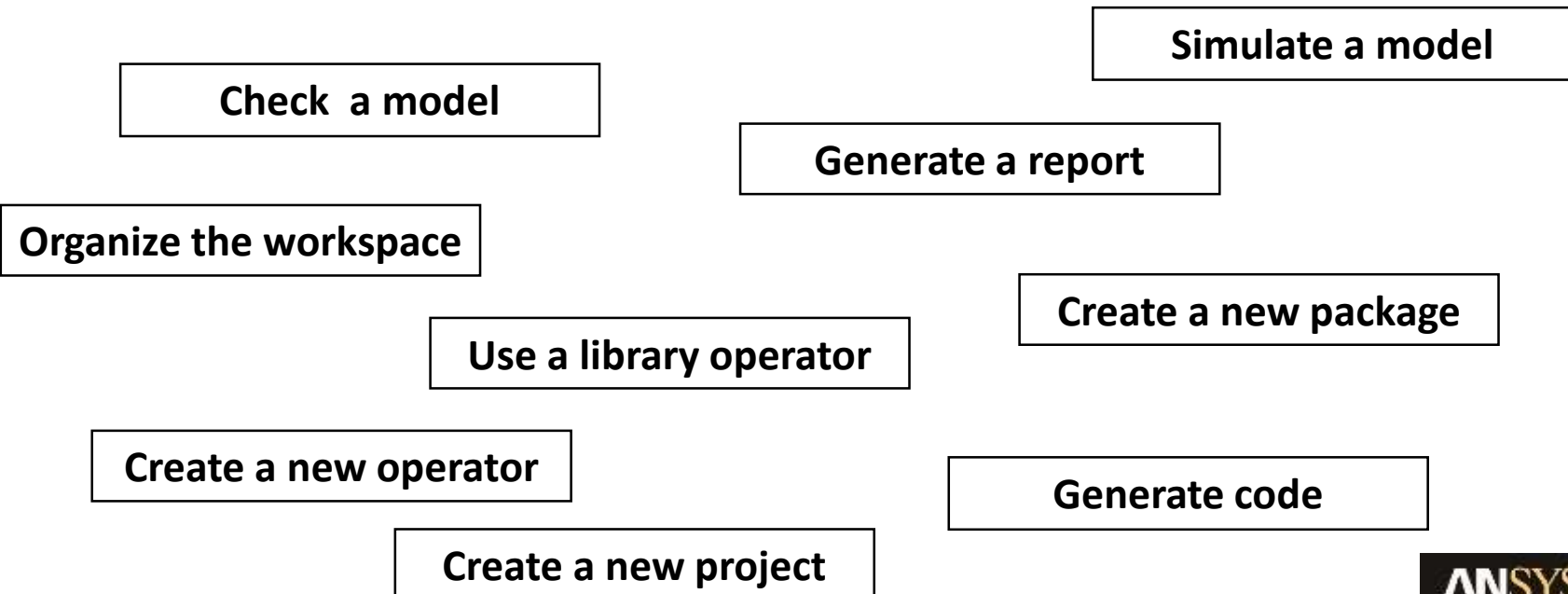
## Objective:

Remember the basic SCADE Suite workflow

Time: 10 min

## Requirements:

Organize the following steps in the correct order:



# AGENDA

## Type declaration

Constant declaration

Operator and interface declaration

Control flow declaration



How to use common types  
such as floats, integers, or  
booleans?

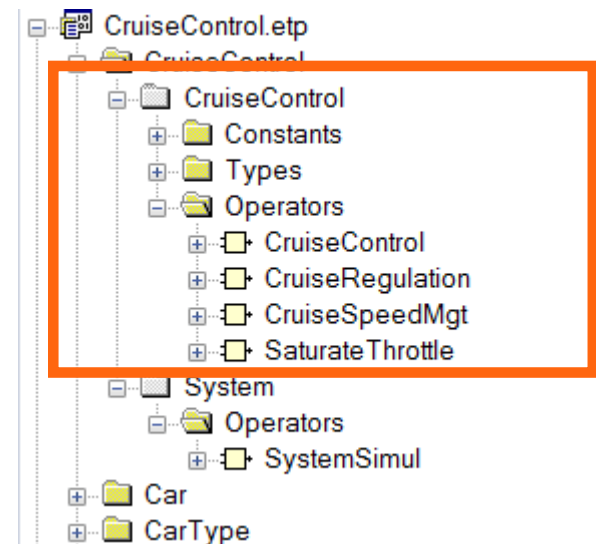
How to implement user-  
defined types to declare the  
system interface?

# New Package Creation

A package allows to define a namespace and group elements together

Select the project / package where the package must be created:

- Create a package by clicking on:
- Give it a name



# Software Engineering Package Rules (Ada)

Ada

Find below the best practices to implement a Scade design for Ada code generation:

- Each declaration is encapsulated into a package
- Sensors, imported and non-imported elements are not mixed in the same package
- No top-down dependency (see next slides)
- No package circular dependency (see next slides)

# Description of Package Dependencies (Ada)

Ada

P1 and P2 are packages:

Element of P1 depends on P2 declarations

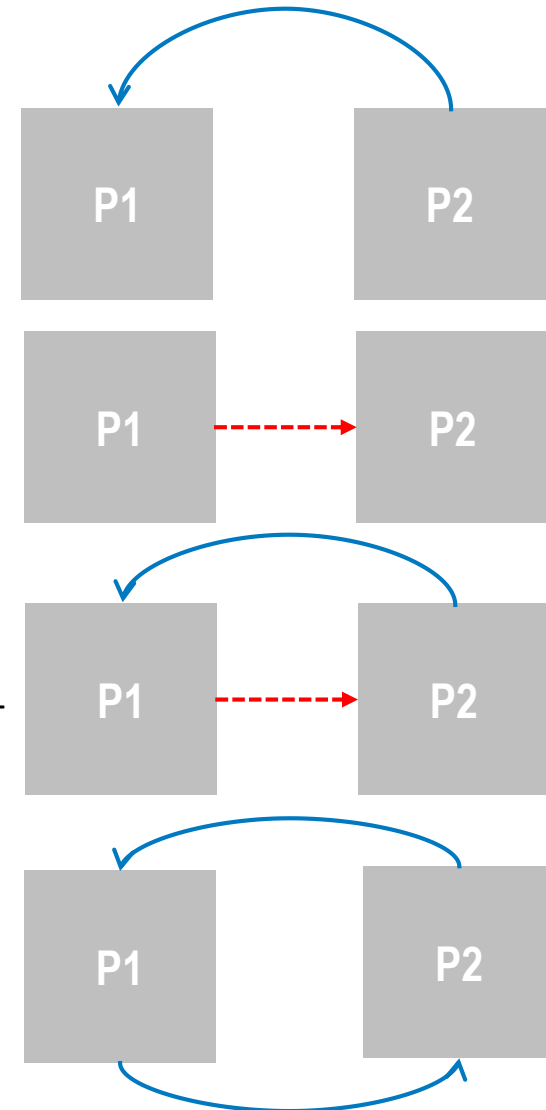
P2 is sub-package of P1

Top-down dependency:

- Element of P1 depends top-down on P2 sub-package declarations

Circular dependency between P1 and P2:

- Element of P1 depends on P2 declarations
- Element of P2 depends on P1 declarations



# Conditions for Direct Translation into Ada Code

Ada

All global Scade declarations belong to package(s)

A package containing imported elements (including sensors) does not contain other non-imported elements at the same time

No package contains a global declaration that depends on one of its sub-packages

- But referring to a sibling package is allowed

There are no circular dependencies between packages



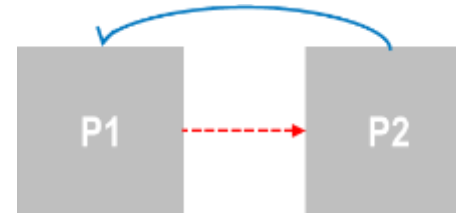
When following constraints are not fulfilled, additional packages ensure the generated code remains SPARK-compliant

A warning message systematically informs users when this occurs:

- Non encapsulated declarations are available:
  - ➔ *Kcg\_Top* package is generated in the code to gather declarations
- A package P contains imported elements and / or sensors mixed with non-imported elements:
  - ➔ *Kcg\_Imported* and / or *Kcg\_Sensor*, sub-packages of P, are generated in the code to separate declarations

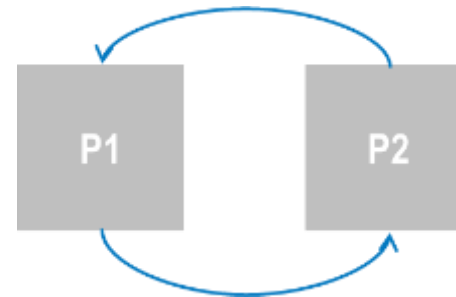
# Ada Code Generation

Ada



Top-down dependency exists:

➔ Declarations of P1 are generated into *Kcg\_sub*, sub-package of P1

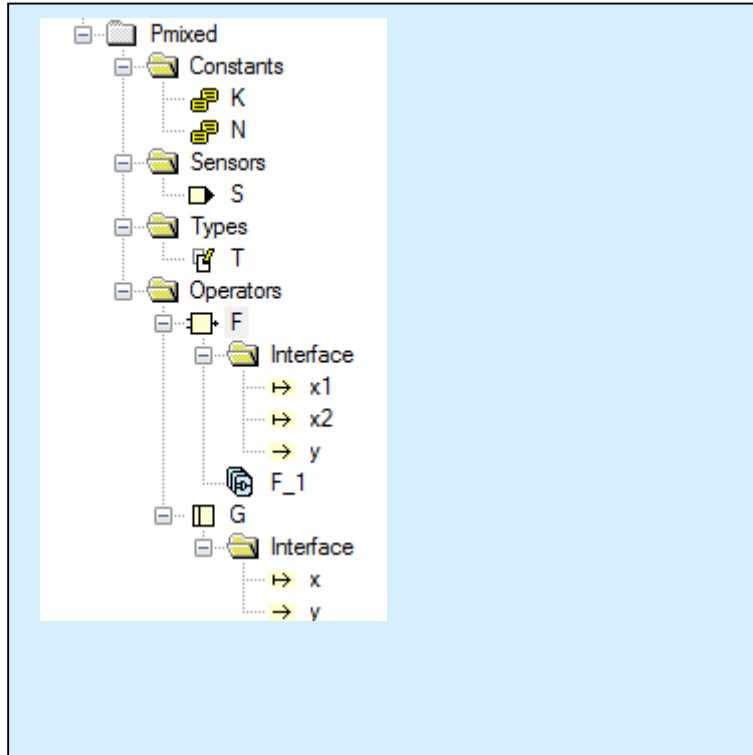


Circular dependencies exist:

➔ Each declaration belonging to the circular dependencies is generated into a new package named *Kcg\_<package>\_<declaration>*

# Mixed Elements (Ada)

Ada



```

package Pmixed
const K : int16 = 42;
const imported N : int16;

sensor S : float64 ;

type imported T;

function F (x1 , x2 : bool) returns (y : bool) ...
function imported G (x : T) returns (y : T);

end;
```

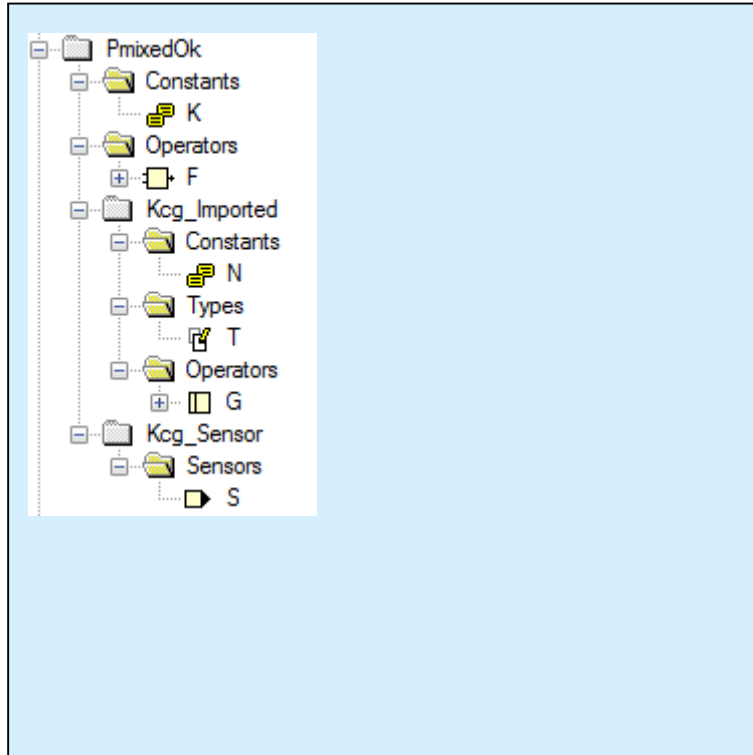
**Code generation: following message is raised to warn that additional packages are produced in the Ada code**

```

Pmixed::: Introducing additional package
Pmixed::: new package "Kcg_Imported" to gather imported declarations
Pmixed::: Introducing additional package
Pmixed::: new package "Kcg_Sensor" to gather sensor declarations
```

# Mixed Elements (Ada)

Ada



```
package PmixedOk
  const K : int16 = 42;
  function F (x1 , x2 : bool) returns (y : bool) ...
```

```
package Kcg_Imported
  const imported N : int16;
  type imported T;
  function imported G (x : T) returns (y : T);
end;
```

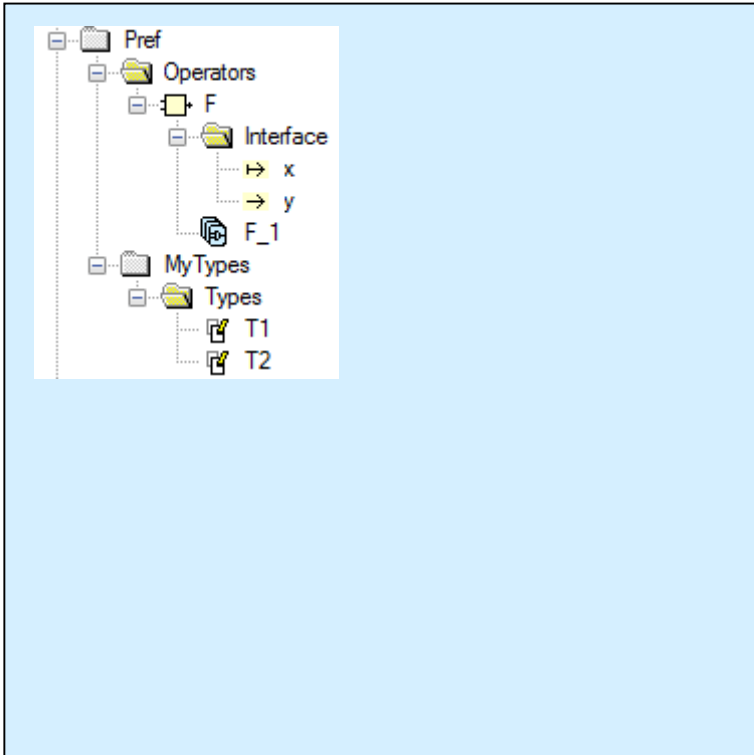
```
package Kcg_Sensor
  sensor S : float64;
end;

end;
```

Direct code generation without additional packages

# References to Sub-Packages (Ada)

Ada



```
package Pref
function F (x : MyTypes :: T1)
  returns (y : MyTypes :: T2) ..

package MyTypes
  type T1 = int32;
  type T2 = T1 ^ 10;
end;

end;
```

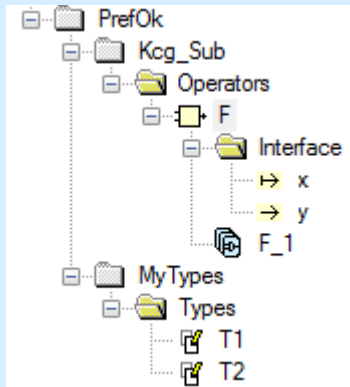
**Code generation: following message is raised to warn that one additional package is produced in the Ada code**

Pref::: Introducing additional package

Pref::: new package "Kcg\_Sub" to handle references to sub-packages

# References to Sub-Packages (Ada)

Ada



```
package PrefOk
  package Kcg_Sub
    function F (x : MyTypes :: T1)
      returns (y : MyTypes :: T2) ...
    end;
```

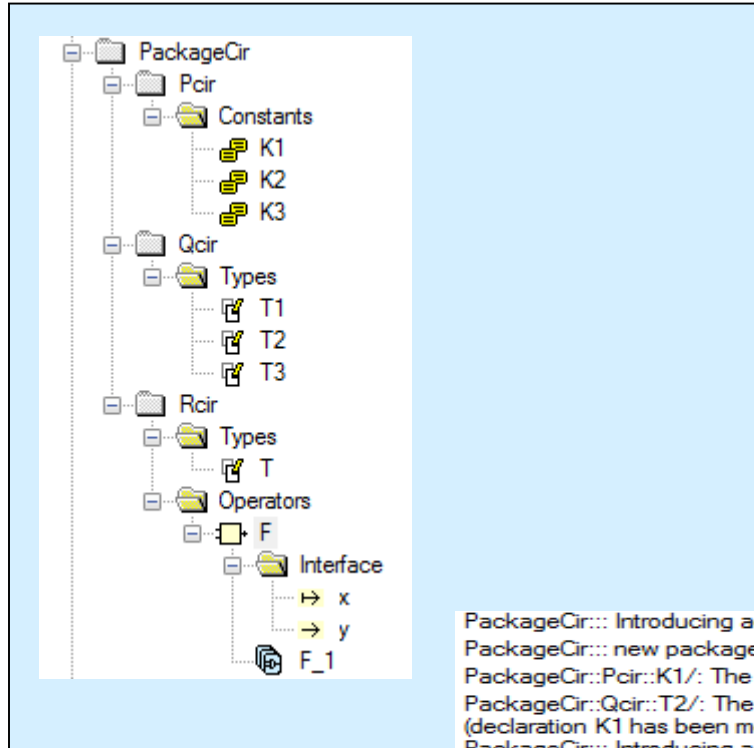
```
package MyTypes
  type T1 = int32
  type T2 = T1 ^ 10;
end;
```

```
end;
-- sibling reference is ok
```

Direct code generation without additional packages

# Circular Dependencies (Ada)

Ada



**Code generation:  
additional packages  
are produced in the  
Ada code**

```

package Pcir
  const K1 : Qcir:: T1 = 42;
  const K2 : int32 = K1 + 1;
  const K3 : int32 = 50;
end;
  
```

```

package Qcir
  type T1 = int32;
  type T2 = T1 ^ Pcir:: K1;
  type T3 = float32 ;
end;
  
```

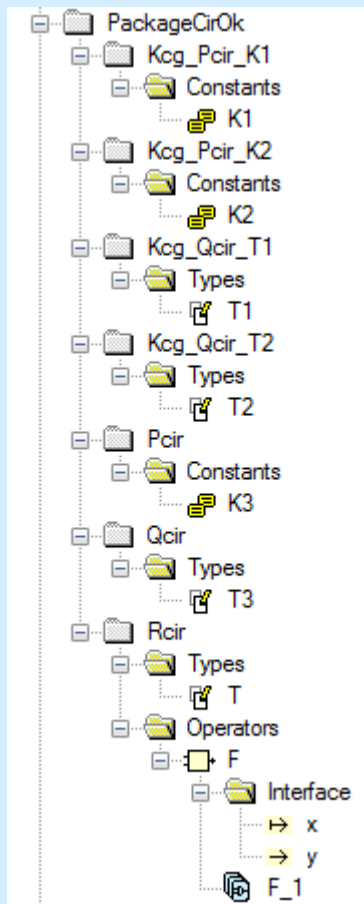
```

package Rcir
  type T = Qcir:: T3 ^ 10;
  
```

PackageCir::: Introducing additional package  
 PackageCir::: new package Kcg\_Pcir\_K1 to break circular dependencies between packages:  
 PackageCir::Pcir::K1/: The declaration of PackageCir::Pcir::K1/ depends on PackageCir::Qcir::T1/  
 PackageCir::Qcir::T2/: The declaration of PackageCir::Qcir::T2/ depends on PackageCir::Pcir::K1/  
 (declaration K1 has been moved from package Pcir to package Kcg\_Pcir\_K1)  
 PackageCir::: Introducing additional package  
 PackageCir::: new package Kcg\_Pcir\_K2 to break circular dependencies between packages:  
 PackageCir::Pcir::K2/: The declaration of PackageCir::Pcir::K2/ depends on PackageCir::Pcir::K1/  
 PackageCir::Pcir::K1/: The declaration of PackageCir::Pcir::K1/ depends on PackageCir::Qcir::T1/  
 PackageCir::Qcir::T2/: The declaration of PackageCir::Qcir::T2/ depends on PackageCir::Pcir::K1/  
 (declaration K2 has been moved from package Pcir to package Kcg\_Pcir\_K2)  
 PackageCir::: Introducing additional package  
 PackageCir::: new package Kcg\_Qcir\_T1 to break circular dependencies between packages:  
 PackageCir::Qcir::T2/: The declaration of PackageCir::Qcir::T2/ depends on PackageCir::Pcir::K1/  
 PackageCir::Pcir::K1/: The declaration of PackageCir::Pcir::K1/ depends on PackageCir::Qcir::T1/  
 (declaration T1 has been moved from package Qcir to package Kcg\_Qcir\_T1)  
 PackageCir::: Introducing additional package  
 PackageCir::: new package Kcg\_Qcir\_T2 to break circular dependencies between packages:  
 PackageCir::Qcir::T2/: The declaration of PackageCir::Qcir::T2/ depends on PackageCir::Pcir::K1/  
 PackageCir::Pcir::K1/: The declaration of PackageCir::Pcir::K1/ depends on PackageCir::Qcir::T1/  
 (declaration T2 has been moved from package Qcir to package Kcg\_Qcir\_T2)

# Circular Dependencies (Ada)

Ada



**Direct code generation without additional packages**

```

package Kcg_Pcir_K1
  const K1 : Kcg_Qcir_T1 :: T1 = 42;
end;

package Kcg_Pcir_K2
  const K2 : int32 = Kcg_Pcir_K1 :: K1 + 1;
end;

package Kcg_Qcir_T1
  type T1 = int32;
end;

package Kcg_Qcir_T2
  type T2 = Kcg_Qcir_T1 :: T1 ^ Kcg_Pcir_K1 :: K1;
end;

package Pcir
  const K3 : int32 = 42;
end;

package Qcir
  type T3 = float32 ;
end;

package Rcir
  type T = Qcir:: T3 ^ 10;
  function F (x : Qcir:: T3) returns (y : T) ...
end;
  
```



# Ada Code Generation

Ada

Function/procedures are translated into their own *.adb* file, with separate body feature

Boolean operators can be short-circuited or not

Predefined Scade types are defined in *Kcg\_Config.ads* file located in `<installation>\SCADE\include\Ada`

# Data Types

The Scade language is strongly typed:

- All the variables are explicitly typed
- Users can define as many types as necessary
- Conversion between values of different types are explicit

**Scade predefined types:**

- Boolean: **bool** and character: **char**
- Signed Integer: **int8**, **int16**, **int32**, **int64**
- Unsigned Integer: **uint8**, **uint16**, **uint32**, **uint64**
- Floating point: **float32**, **float64**



# Data Types

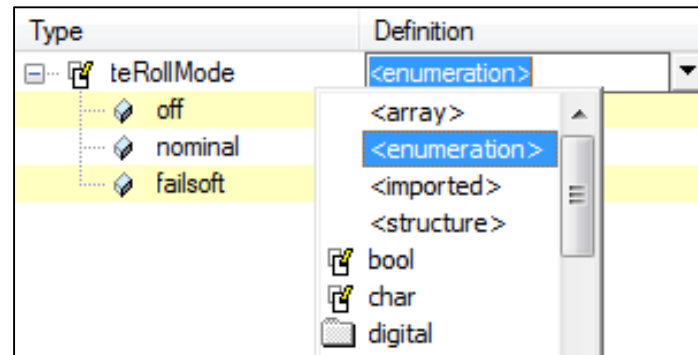
User defined types:

- **Equivalent types:**
  - **Speed = Velocity**
  - **Euro = Currency**
  - Enhances model readability
- **Enumerated types:**
  - **Colors = enum {blue, green, red}**
- **Structured types:**
  - Data structure: corresponds to usual structures or records
    - **Ts = {x: int8, y: float32}**
  - Array: always fixed size, the first index is 0
    - **Vector = float64<sup>3</sup>**
    - **Table = float32<sup>5</sup><sup>5</sup>**

# New Type Creation

Select the package where the type must be created:

- Create a new type by clicking on 
- Select the type definition from the Definition column:
  - If the type definition needs a new definition element, click on 
  - Or right click on the type and select Insert > Definition element



# New Type Creation

Type	Definition
array_3_3r	<array> ^3 array_3r
array_3r	<array> ^3 float64
degree	float64
kelvin	float64
Sx	float32
Sy	float32
Sz	float32
tColorInputs	<enumeration> eRed eYellow eGreen
tComplex	<structure> re im

scalar

enumeration

array^3^3

array^3

structure

# AGENDA

Type declaration

**Constant declaration**

Operator and interface declaration

Control flow declaration



# Scade Constants

Constants are data having the same value throughout the program execution

Constants can appear as:

- 3.14
- A textual expression:  $4 * C1$
- A name defined by a constant:  $C1 = 2.5$  or  $C2 = 3.0 * C1$

Used for:

- Constant value definitions
- Parameters or settings of Scade constructs such as the size of an array



# Literals

Literals  can be defined as:

## A constant literal:

- `true`, `false`: Boolean
- `1.0` (floating-value) is different from `1` (integer)
- `10_ui32` has `uint32` type
- `10_i16` has `int16` type
- `10.0_f32` has `float32` type
- Quote for char constants: `'a'`

## An expression of constants:

- `C1` and `C2`
- `a * c + b`

Expressions used to define constants must be statically defined at compilation time

# Named Constants

Constants are always typed:

By named types:

- `MyCons : MyType = [3, 2];` (with `MyType = uint32^2`)

Or anonymous types:

- `MyCons : uint32^2 = [3, 2];`


Used in graphical and textual Scade

# New Constant Creation

Constant	Type	Value
Ki	float64	0.5
Kp	char	8.113
PedalsMin	digital	3.0
RegulThrottleMax	filters	45.0
SpeedInc	float32	2.5
SpeedMax	float64	150.0
SpeedMin	int16	30.0
ZeroPercent	int32	0.0
ZeroSpeed	int64	0.0

The image shows a screenshot of a software interface for creating constants. A dropdown menu is open for the 'Type' column, showing a list of data types: char, digital, filters, float32, float64 (highlighted), int16, int32, int64, int8, and linear. The 'Constant' column lists various parameters like Ki, Kp, PedalsMin, RegulThrottleMax, SpeedInc, SpeedMax, SpeedMin, ZeroPercent, and ZeroSpeed. The 'Value' column shows numerical values for each parameter.

Select the package where the constant must be created:

- Create a constant by clicking on: 
- Give it a name and select its type: click on cell and browse in the displayed list
- Fill the appropriate fields according to the type definition

# New Constant Creation

Constant	Type	Value
BackgroundColor	tColorInputs	eGreen
cMode	char ^10	"Active "
ComplexA	tComplex	{re : 10.0, im : (-10.0)}
ComplexB	tComplex	
re	float64	5.5
im	float64	15.7
IdentityMatrix	array_3_3r	
0	array_3r	
0	float64	1.0
1	float64	0.0
2	float64	0.0
1	array_3r	[0.0, 1.0, 0.0]
2	array_3r	[0.0, 0.0, 1.0]
PosX	Sx	124.4
PosY	Sy	-4.589
PosZ	Sz	7.25
VectInit	array_3r	
0	float64	2.0
1	float64	2.0
2	float64	2.0

enumeration

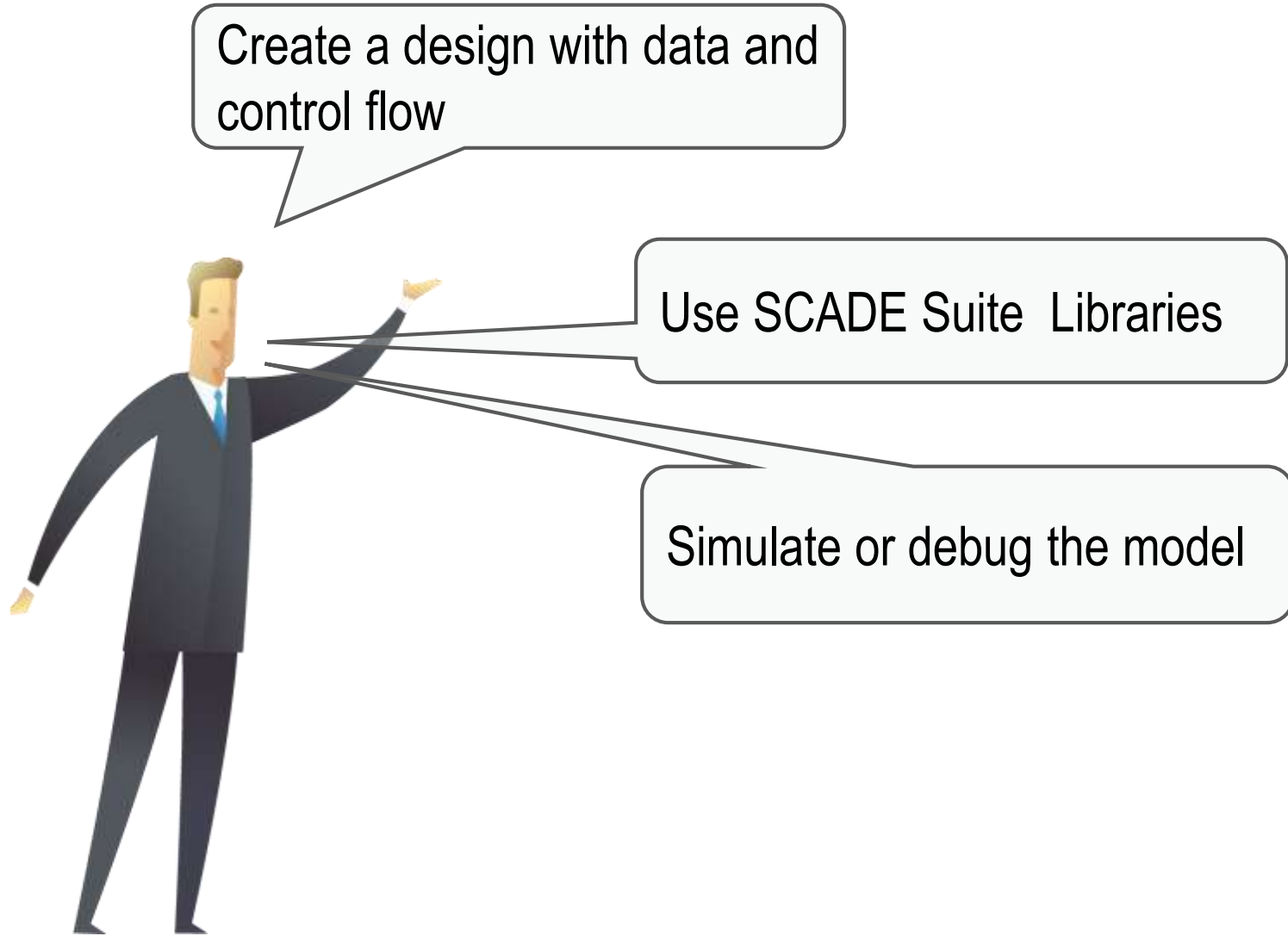
structure

array^3^3

scalar

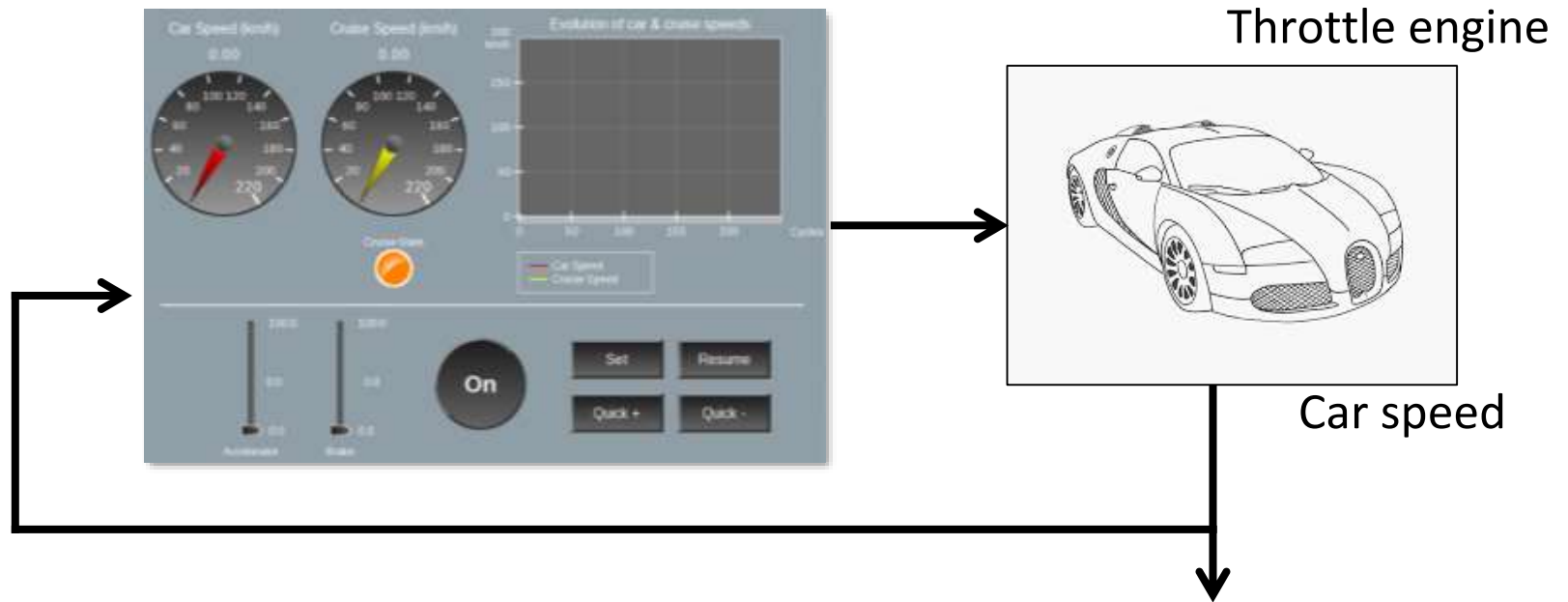
array^3

# Lab Objectives



# The System to Create

Build a system of a Cruise Control to drive the car speed in order to maintain a constant cruise speed



During the Lab, you will create bricks and assemble them to obtain the whole application

# Requirements

The software requirements of Cruise Control design are defined in *\Day 2\Labs66\Requirements*

The labs refer these requirements. However, you should be able to do the design without needing to read this document

# Lab 1 (1/2)

Lab Support p.8-21

## Objective:

Create a new type and new constants

Time: 20 min

## Requirements:

Create a new SCADE Suite Project and a package:

**Project Name:** `Cruise Control`

**SCADE Suite library:** `Car.etp` (located at `Q:\Labs\Prerequisites\Lab 1\Libraries\Car`)

**Package Name:** `CruiseControl`

Create a new type enumerate `teCruiseState` (*defined in requirement CC\_HLR\_OUT\_03*)

Type	Definition
<code>teCruiseState</code>	<code>enum {OFF, INT, STDBY, ON}</code>



# Lab 1 (2/2)

Create the constants (*defined in requirement CC\_HLR\_CCP\_01 to CC\_HLR\_CCP\_07*)

Name	Type	Value
KI	float64	0.5
KP	float64	8.113
PEDALS_MIN	CarType::tPercent	3.0
REGUL_THROTTLE_MAX	CarType::tPercent	45.0
SPEED_INC	CarType::tSpeed	2.5
SPEED_MAX	CarType::tSpeed	150.0
SPEED_MIN	CarType::tSpeed	30.0
ZERO_PERCENT	CarType::tPercent	0.0
ZERO_SPEED	CarType::tSpeed	0.0

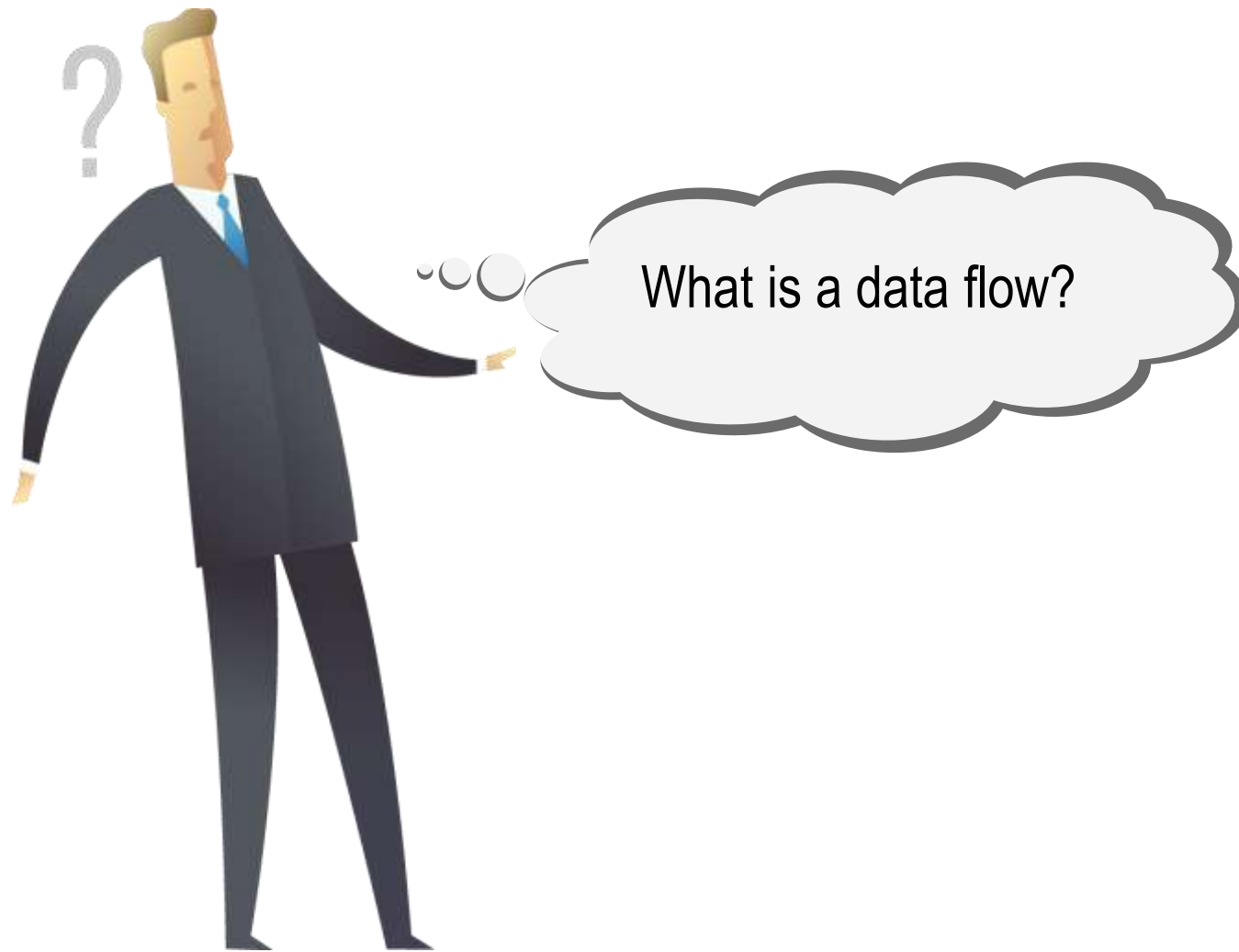
# AGENDA

Type declaration

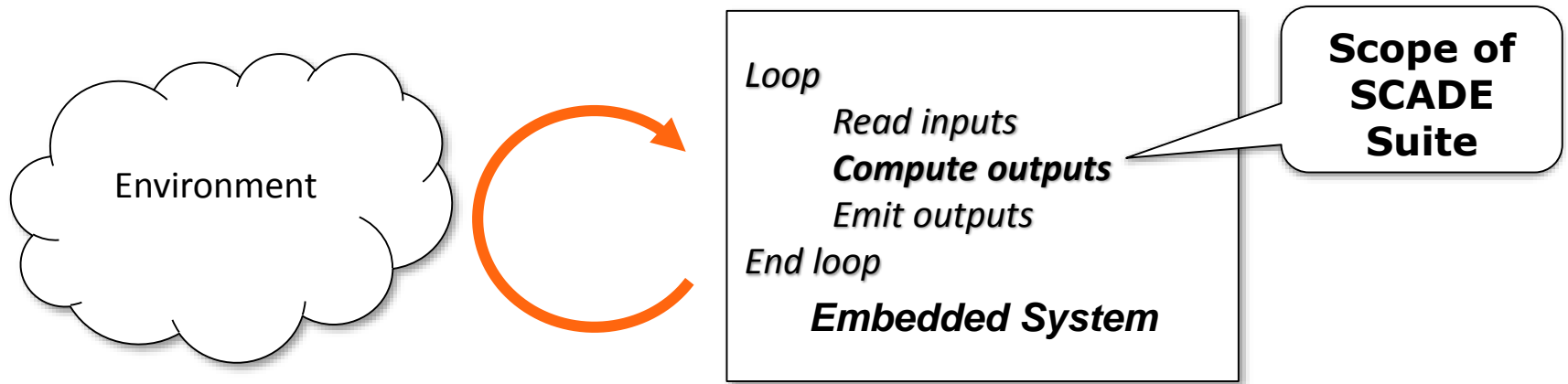
Constant declaration

**Operator and interface declaration**

Control flow declaration



# Scade Works on Data Flows



- A SCADE Suite model is a cycle-based computation model where the program examines its inputs and provides the outputs in the same cycle
- Inputs and outputs are data flows
- A flow has a single definition and must always be defined
- Flows are streams of values, that are infinite sequences of values of a given type

# Scade Works on Data Flows

If  $v$  is a data flow expression, its values sequence are  $v_1, v_2, \dots$

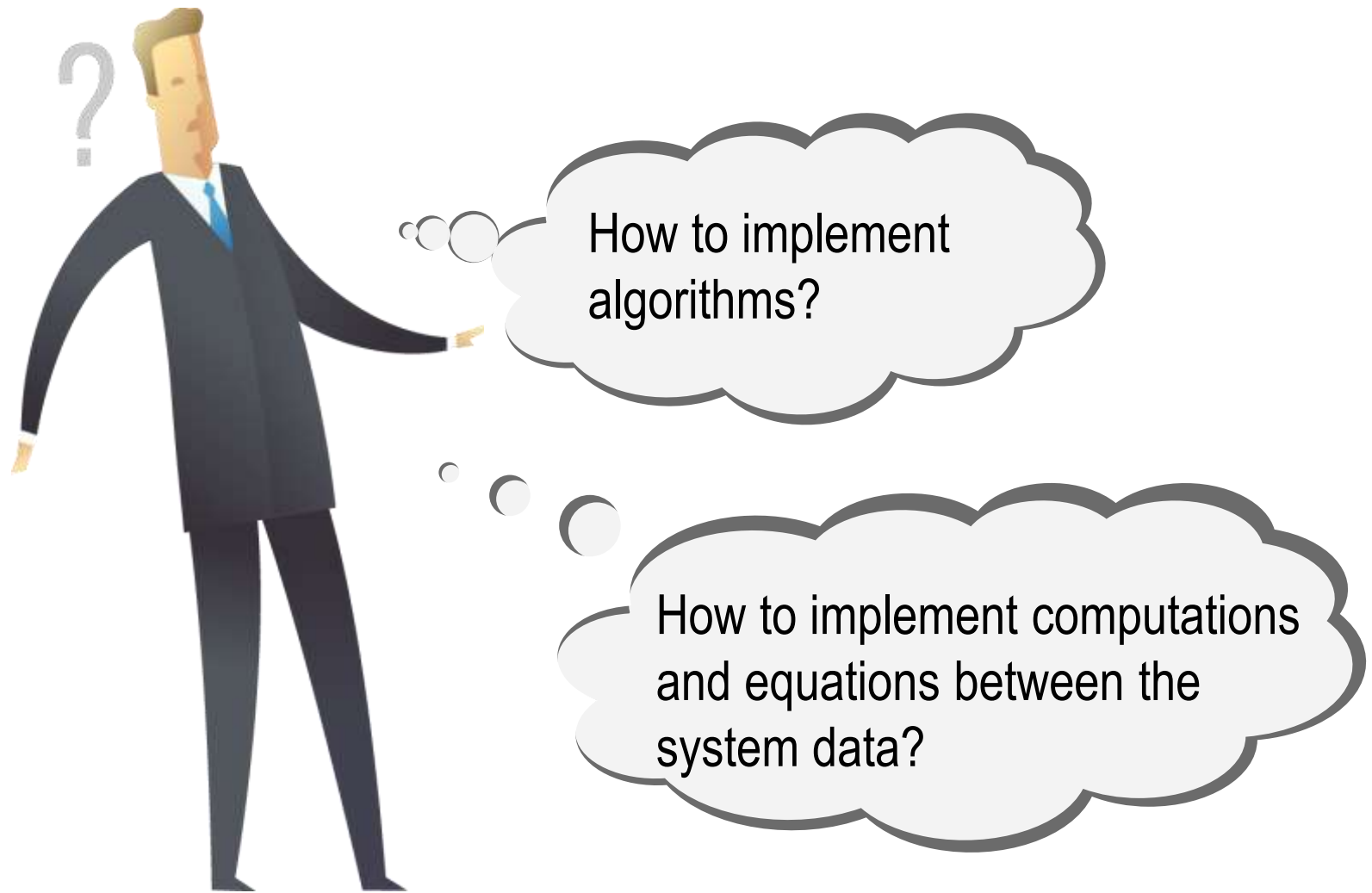
$3$  denotes the constant flow  $3, 3, \dots$

$a + b$  denotes the flow  $a_1 + b_1, a_2 + b_2, \dots$

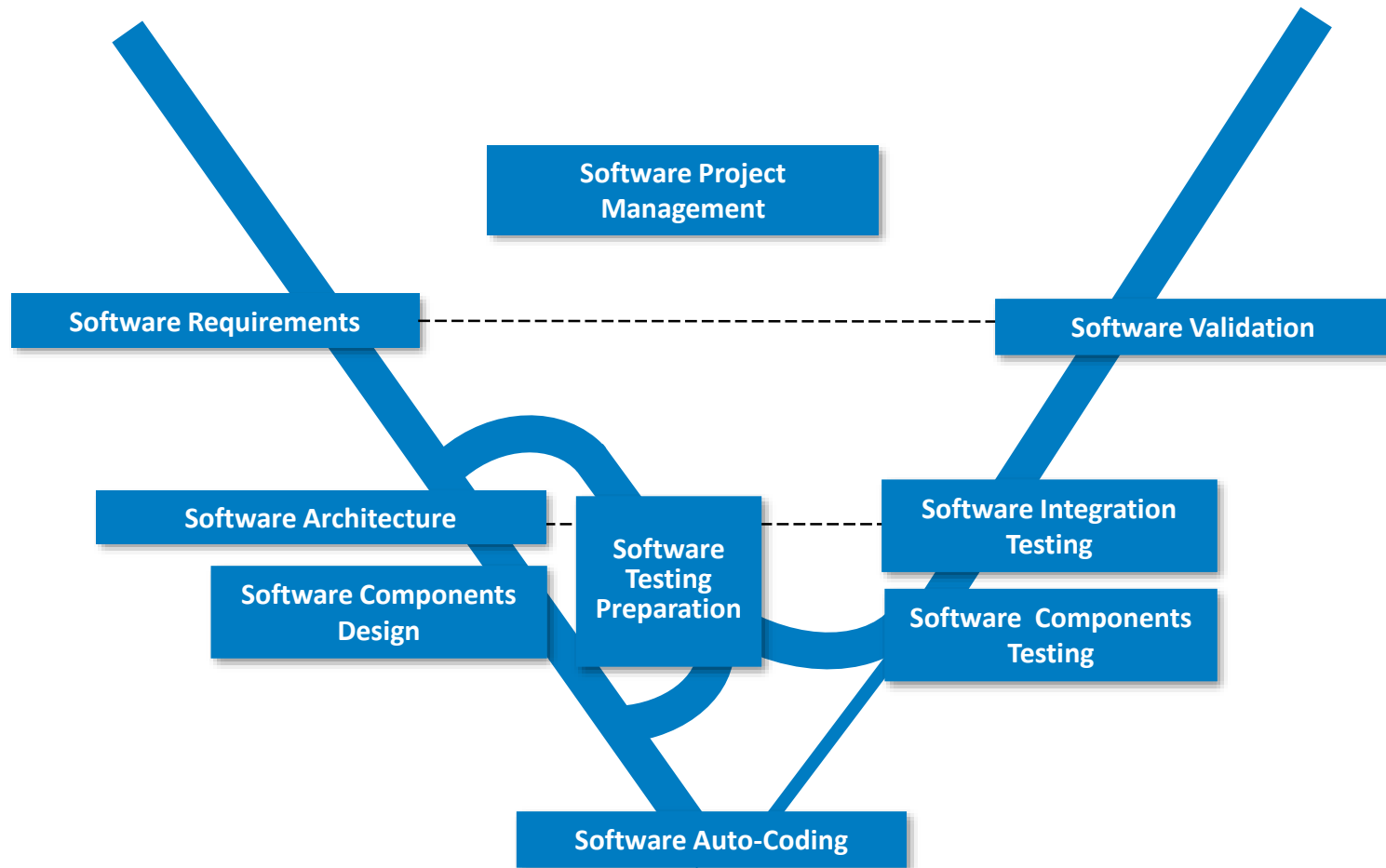
Outputs computation may require input and output values of previous cycles:

The relationship between inputs and outputs is time-dependent

Cycle	1	2	3	4	...
3	3	3	3	3	...
a	a1	a2	a3	a4	...
b	b1	b2	b3	b4	...
a+b	a1+b1	a2+b2	a3+b3	a4+b4	...
pre(a)	nil	a1	a2	a3	...



# SCADE V-Cycle

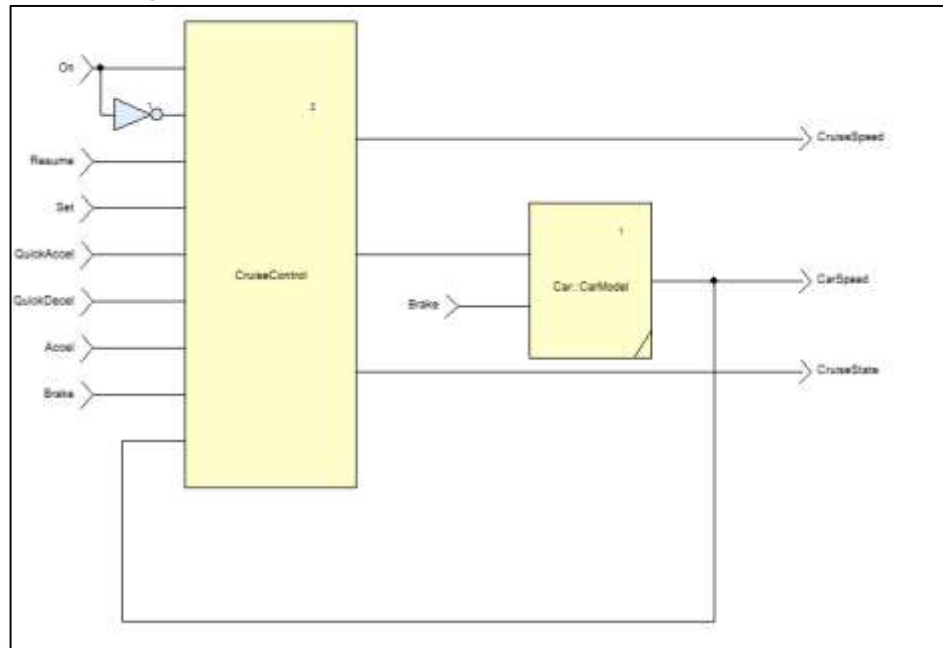


# Software Architecture

Once the software requirements phase is achieved, software engineers define the software architecture before implementing the design.

A software architecture in a SCADE Suite model is a network of top level user-defined operators drawn as boxes.

Software architecture  
can be done with  
SCADE Architect,  
bundled with SCADE  
Suite





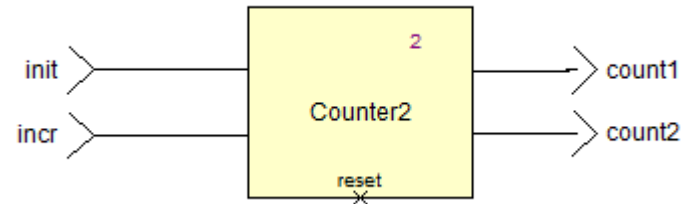
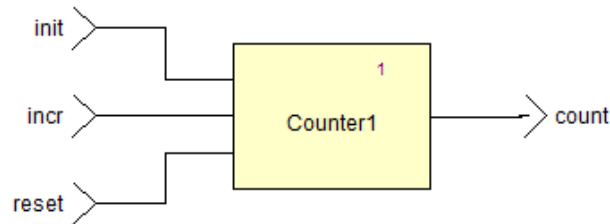
# Software Design

A SCADA Suite design of a software is composed of

- A network of user-defined operators and
- Predefined operators such as '+', 'and'

Operators may be purely combinatorial, or may have memory (like a counter or an integrator)

- Independent operators data are not ordered and compute in parallel
- Operators only communicate through the flows



# User Operator

## Set of equations:

- Graphical, State Machine or textual operator
- One definition for each output and each local variable
- Indifferent order in equations

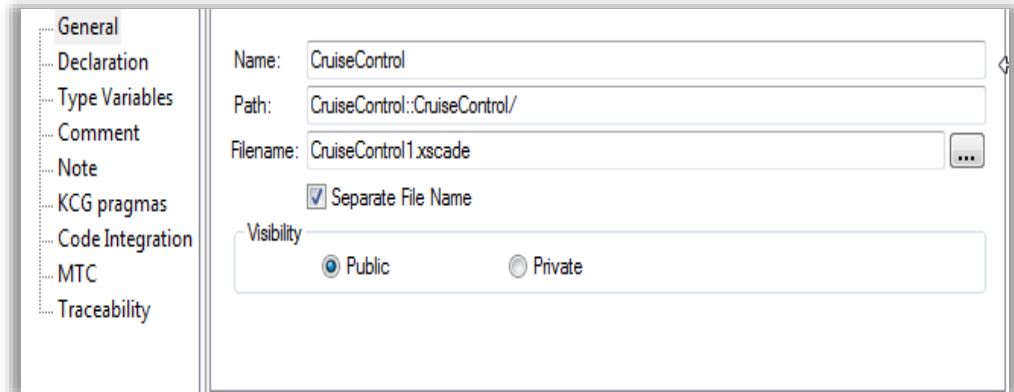
## Modularity / Reuse:

- An operator is a black box and is connected through its formal interface
- Each operator defined by the user can be reused

# Operator Creation

Click on:  (Shortcut: *Ctrl+Shift+N*)

Select the created operator  
in the Scade view



General

Declaration

Type Variables

Comment

Note

KCG pragmas

Code Integration

MTC

Traceability

Name: CruiseControl

Path: CruiseControl::CruiseControl/

Filename: CruiseControl1.xscade

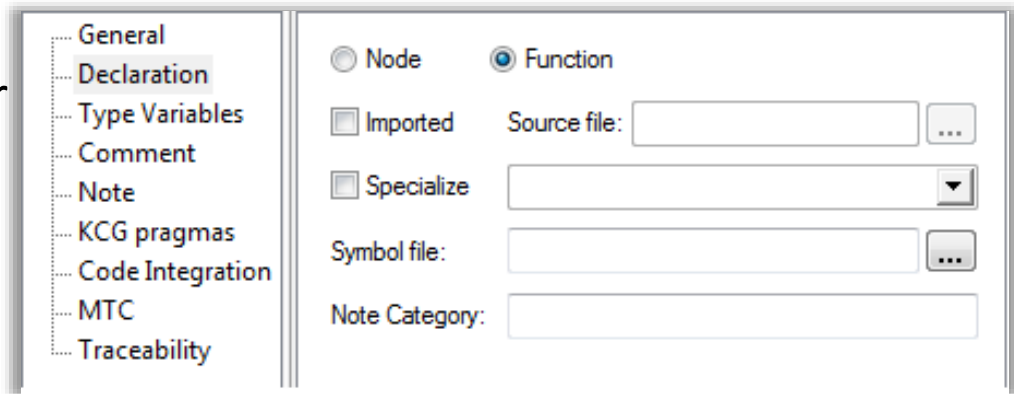
☒ Separate File Name

Visibility

☒ Public ☐ Private

Give a meaningful name to:

- The operator
- The file where the operator is saved if separated



General

Declaration

Type Variables

Comment

Note

KCG pragmas

Code Integration

MTC

Traceability

☐ Node ☒ Function

☐ Imported Source file:

☐ Specialize

Symbol file:

Note Category:

# Operator Creation

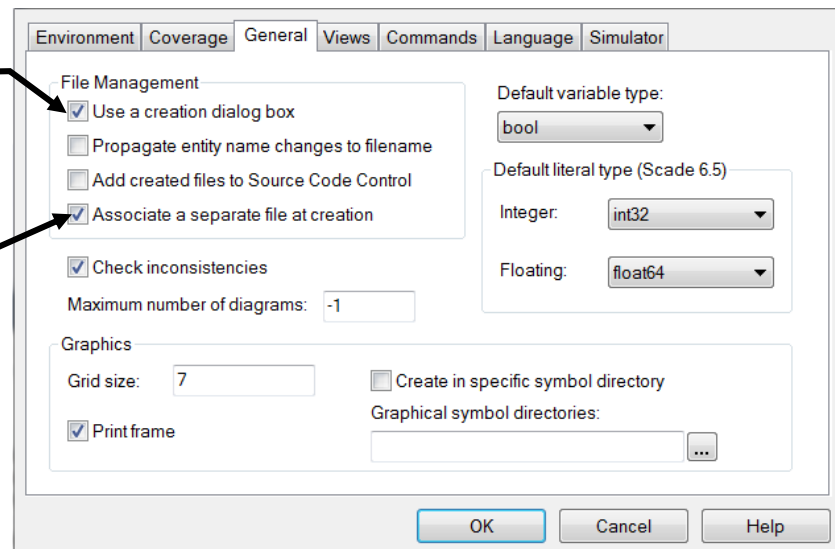
It is possible to set global options to:

- Directly set the name of the operator and its file using a creation dialog box
- Propagate changes on the operator name to the file name

Open *Tools > Options > General tab*

Select to have a dialog box entry to edit the operator's name

Select to propagate the operator's name to the underlying file name

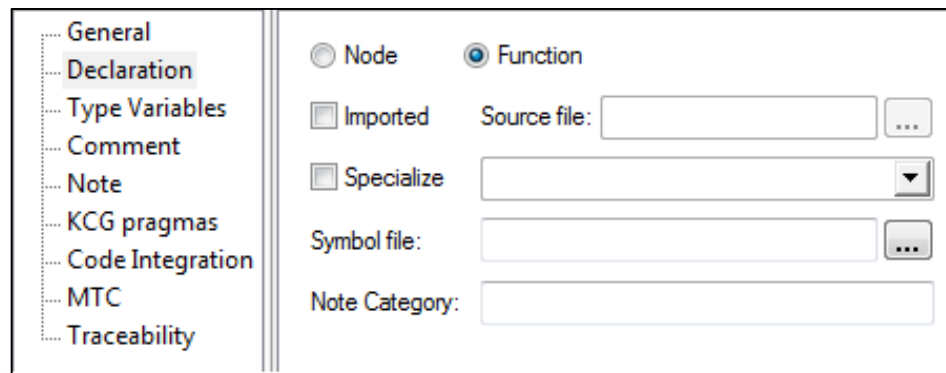


# Functions/Nodes

By default, an operator is declared as a node (operator with internal memory)

A function is a stateless operator:

- Contains no internal memory (State machine, last, pre, fby, etc. do have memory!)
- Enables function specific code generator optimizations
- Setting accessed via the property dialog:

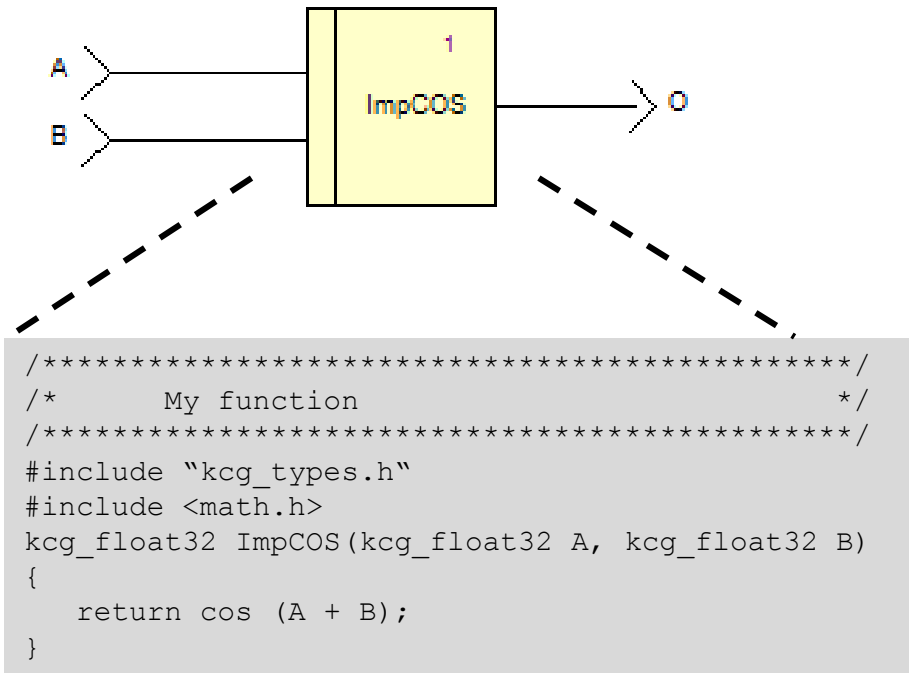


The screenshot shows a property dialog box with a left sidebar and a main content area. The sidebar contains the following items: General, Declaration (highlighted), Type Variables, Comment, Note, KCG pragmas, Code Integration, MTC, and Traceability. The main content area has two radio buttons at the top: 'Node' and 'Function', with 'Function' selected. Below these are four rows of settings: 1) 'Imported' checkbox (unchecked) and 'Source file:' text box with a browse button (...). 2) 'Specialize' checkbox (unchecked) and a dropdown menu. 3) 'Symbol file:' text box with a browse button (...). 4) 'Note Category:' text box.

# Imported Operators (C Code)

The aim of an imported operator is to use external code:

- The operator's name and interface are defined in SCADE Suite, but not the operator's body
- The implementation is provided through the target code

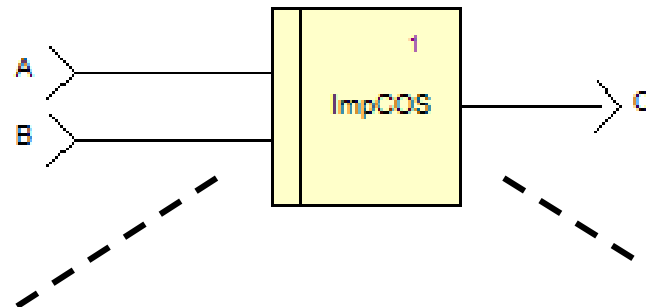


# Imported Operators (Ada)

Ada

The aim of an imported operator is to use external code:

- The operator's name and interface are defined in SCADE Suite, but not the operator's body
- The implementation is provided through the target code

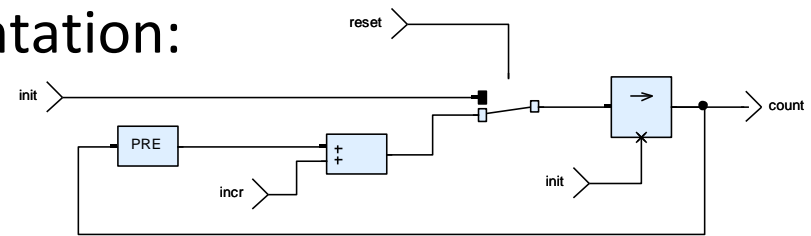


```
function ImpCOS(  
  A : in Kcg_Config.Kcg_Float32;  
  B : in Kcg_Config.Kcg_Float32) return  
Kcg_Config.Kcg_Float32  
is  
  O : Kcg_Config.Kcg_Float32;  
begin  
  O:= cos (A + B);  
  return O;  
end ImpCOS;
```

# Operator Interface

## Graphical SCADE Suite representation:

- An operator is a block diagram

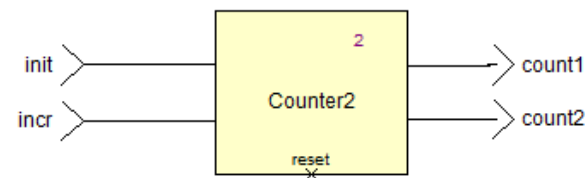
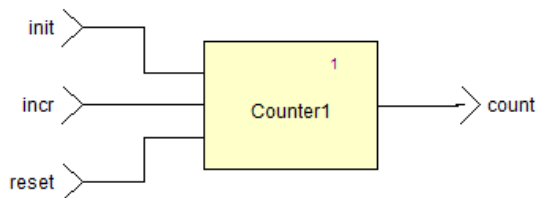


An operator is seen as a black box and is connected through its formal I/O interface

## Textual Scade representation:

```
node Counter1(init:int32; incr:int32; reset:bool)
returns (count:int32)
```

```
node Counter2(init:int32; incr:int32; reset:bool)
returns (count1:int32; count2:int32)
```



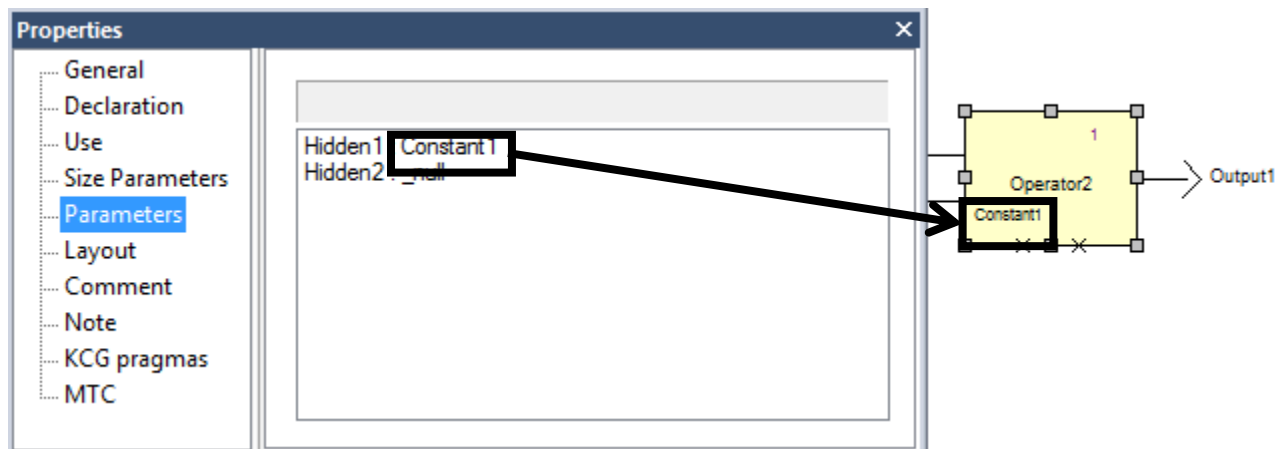


# Hidden Inputs

## Goals:

- Reduces the number of links and increases readability
- Allows to distinguish between parameters and data flows
- Makes a graphical difference when the operator is instantiated

Direct assignment by connection or by name:



From a code generation point of view, no difference between a hidden input and an input

# Operator Interface Creation

Input:



(Shortcut: Ctrl+Shift+I)

Output:



(Shortcut: Ctrl+Shift+O)

Hidden input:



(Shortcut: Ctrl+Shift+H)

- Select the created input/output
- Give a meaningful name
- Select the type of the input or output

General

Declaration

Clock

Comment

Note

ASAM MCD-2 MC

KCG Pragmas

Simplorer

Traceability

Type: bool

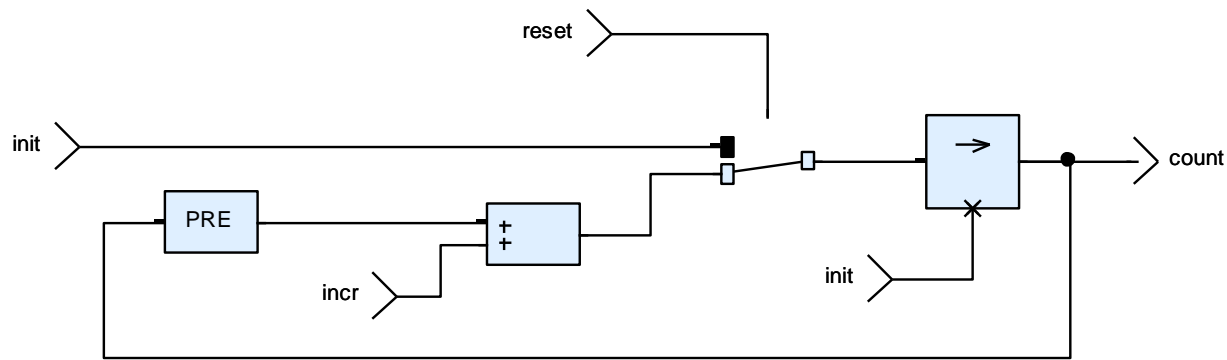
Last:

Default:

Kind: → Output

# Operator Body

A block diagram example:



A textual diagram example:

```
count = init -> if reset then init else pre(count) + incr;
```

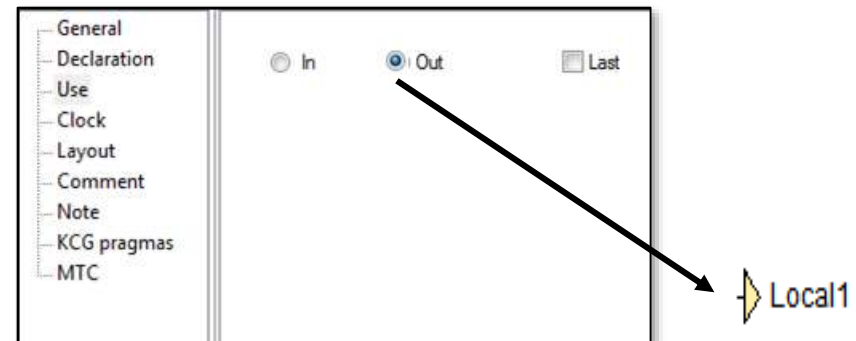
# Local Variables

Seen only by the operator/conditional block/state in which it is declared: 

- Must be defined only once
- Can be read several times

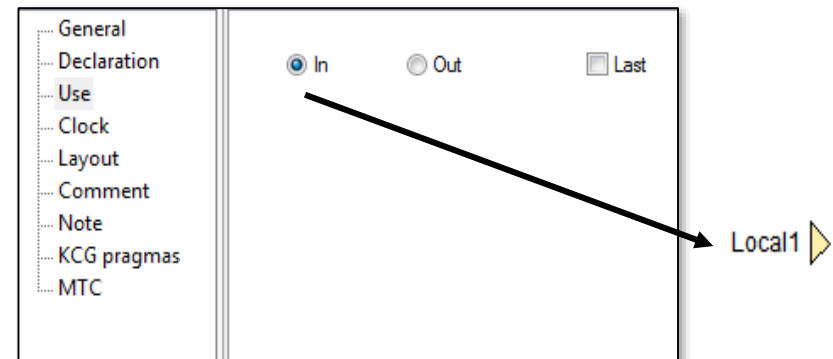
To produce it :

- Open its Properties window
- Click under Use tab on Out
- Or press on CTRL + Drag&Drop it from its definition




To use it :

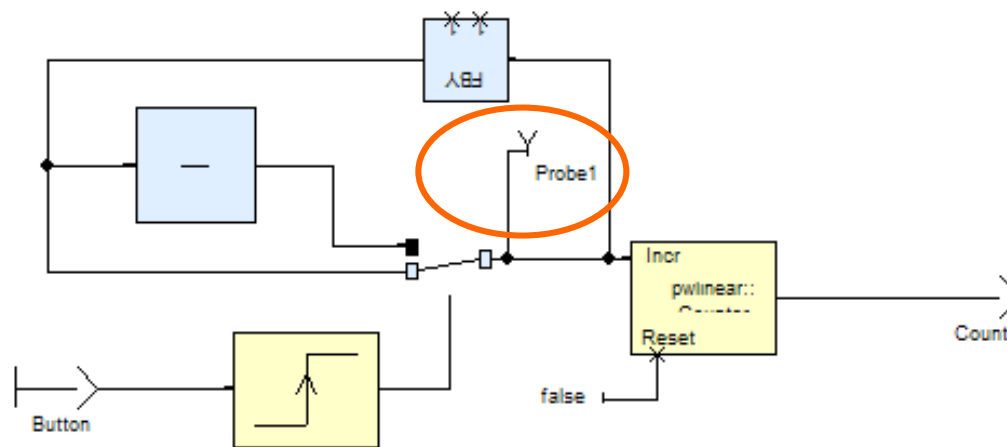
- Open its Properties window
- Click under Use tab on In



# Probe

Defined an observation point of internal variable  
Can be accessed during simulation or execution of the  
generated code  
Created by clicking on button 

Enhance observation by allowing the watch of any internal  
flow during simulation or execution of the generated code



# Sensor

Global variable accessible in read-only mode


Assumed to be provided by the environment

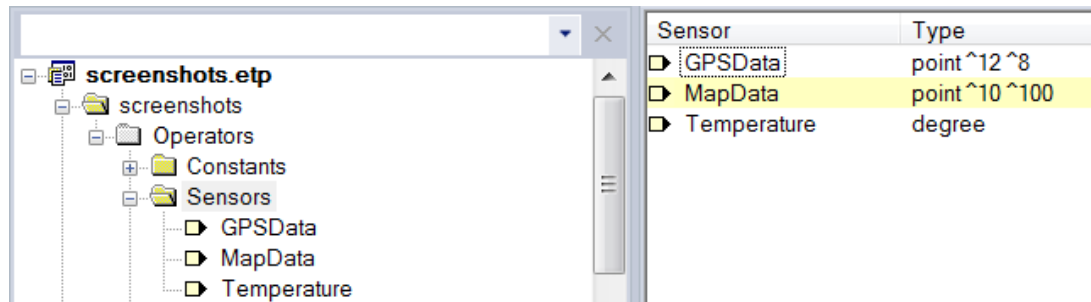
Remains constant during the whole execution cycle

Provides an alternative to the operator input interface, intended for data that must be accessed throughout the design

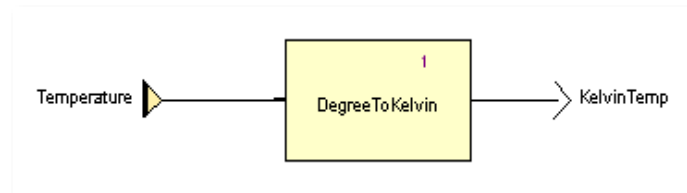
# Creating and Using Sensors

Select a package/project where the sensor must be created

- Create a sensor by clicking on: 
- Give it a name and select the type of the sensor



Use a sensor by drag and drop on a diagram



# Lab 2

Lab Support p.23-24

## Objective:

Define the architecture operators

## Requirements:

Time: 10 min

Create the 4 following operators (and packages)

Operators	Package	Comments
CruiseRegulation	CruiseControl	manages the vehicle speed when the Cruise Control is active
CruiseSpeedMgt	CruiseControl	manages the value of the cruise speed according to the driver commands
CruiseControl	CruiseControl	manages the cruise control behavior
SystemSimul	System	connects the other operators and the Car model



# Scade Predefined Operators

Basic operators of the Scade language

Can be used either in **textual or graphical** design

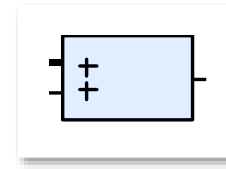
- Graphically, they have a blue fill color to quickly identify them

Fully documented in the **Scade Language Reference Manual**

Grouped by functionality:

- Mathematical
- Logical
- Bitwise
- Time
- Comparison
- Choice
- Structure/arrays
- Higher order: activation constructs

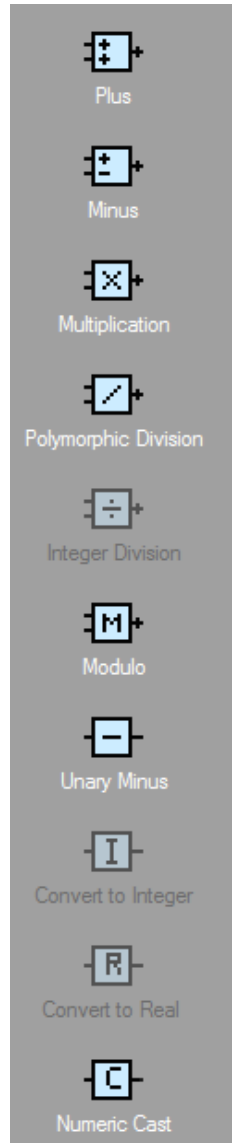
# Mathematical Operators



They apply to numeric types

- Type interface: 'N or 'N, 'N → 'N, where N is Numeric (float or integer)
- Example

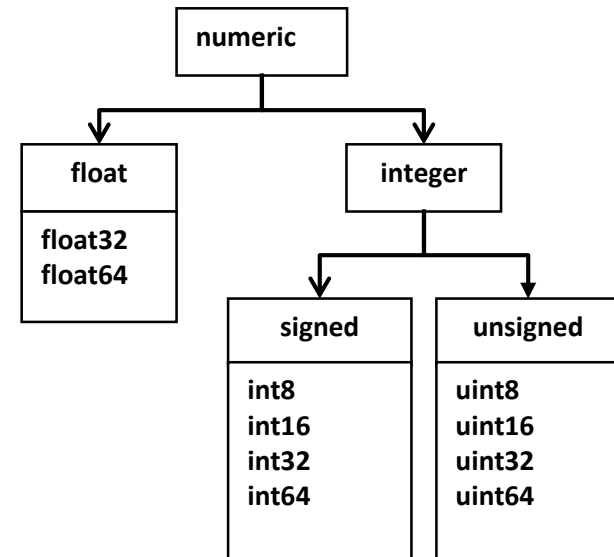
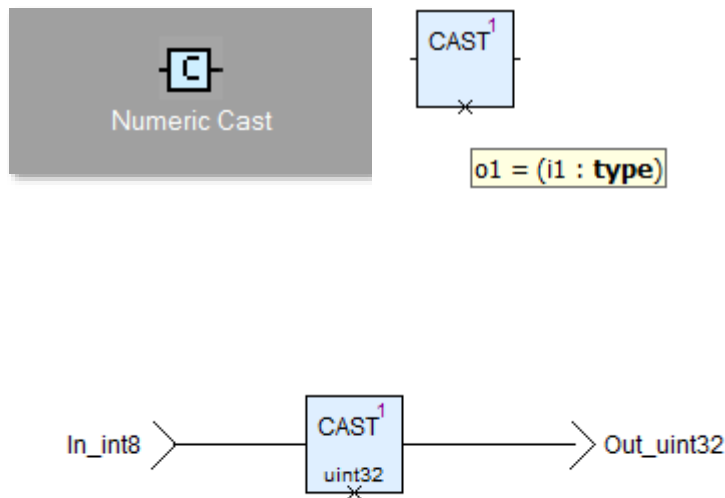
Cycle	1	2	3	4
a	a1	a2	a3	a4
b	b1	b2	b3	b4
c = a+b	a1+b1	a2+b2	a3+b3	a4+b4



# Cast

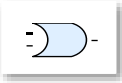
A **cast** operator allows the conversion of any data into an expression of the proper numeric type:

- $(a + b: \text{uint16})$ : conversion to uint16
- $(4: \text{int8})$ : specification of a literal



# Logical Operators

They apply to the bool type:

- Type interface: bool or bool, bool  $\rightarrow$  bool
- Note:  $\#(e1, \dots, en)$  is true if zero or one of the  $e_i$  is true
- Example 

Cycle	1	2	3	4
a	a1	a2	a3	a4
b	b1	b2	b3	b4
$C = a \text{ or } b$	$a1 \text{ or } b1$	$a2 \text{ or } b2$	$a3 \text{ or } b3$	$a4 \text{ or } b4$



And



Or



Exclusive Or




Sharp



Not

# Bitwise Operators

They apply to unsigned integer types

- Type interface:  
unsigned integer, unsigned integer → unsigned integer
- Example: 

Cycle	1	2	3	4
a	a1	a2	a3	a4
b	b1	b2	b3	b4
C = a & b	a1 & b1	a2 & b2	a3 & b3	a4 & b4



# Add New Input for Predefined Operators

Select a predefined operator in the diagram view:

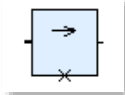
- From the Properties Use tab:
  - Add new input pins:
    - Use “Input number” selector
    - Enter the number of inputs directly in the field

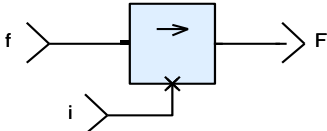
The screenshot shows the 'Use' tab in the Properties panel. On the left is a tree view with 'Use' selected. The main area contains two sections: 'Operator' and 'Higher Order'. The 'Operator' section has an 'Input number' field with the value '3', a small red 'X' icon to its right, and a 'Symmetric' checkbox. Below it is an 'Instance name' field with the value '2'. The 'Higher Order' section has a dropdown menu set to '<none>' and an empty 'Instance name' field. A black arrow points from the 'Input number' field to the 'Instance name' field.

- Remove input pins:

# Time Operators: Initialization ->


Initializes flows


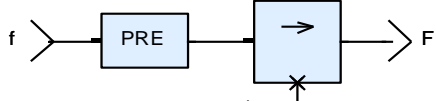
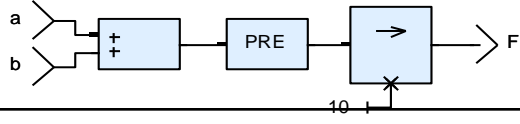
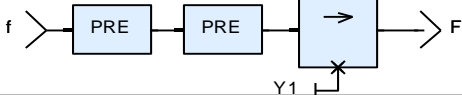
- Type interface: 'T, 'T  $\rightarrow$  'T, where 'T is any Type
- Example: 

Cycle	1	2	3	4
<b>i</b>	<b>i1</b>	<b>i2</b>	<b>i3</b>	<b>i4</b>
<b>f</b>	<b>f1</b>	<b>f2</b>	<b>f3</b>	<b>f4</b>
<b>i -&gt; f</b> 	<b>i1</b>	<b>f2</b>	<b>f3</b>	<b>f4</b>

# Time Operators: pre

Produces the previous value of its input

- Type interface: 'T  $\rightarrow$  'T, where 'T is any Type
- At the first cycle, the previous value is not defined
- The Initialization operator ensures the complete definition
- Examples: 

f		f1	f2	f3	f4
<b>pre (f)</b>		<b>nil</b>	<b>f1</b>	<b>f2</b>	<b>f3</b>
<b>1 -&gt; pre (f)</b>		<b>1</b>	<b>f1</b>	<b>f2</b>	<b>f3</b>
<b>10 -&gt; pre (a+b)</b>		<b>10</b>	<b>a1+b1</b>	<b>a2+b2</b>	<b>a3+b3</b>
<b>y -&gt; pre (pre (f))</b>		<b>y1</b>	<b>nil</b>	<b>f1</b>	<b>f2</b>



# Exercise 1: Rising Edge

## Objective

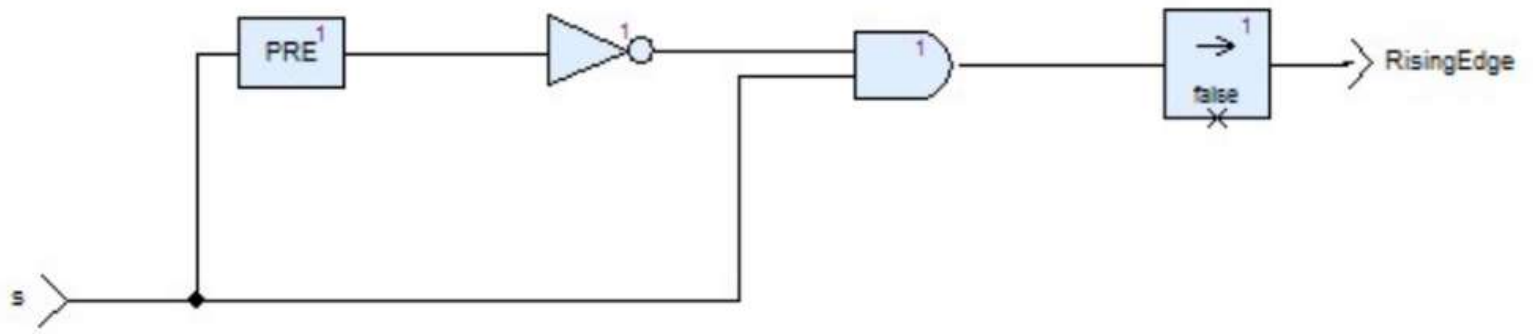
Implement a rising edge operator using `pre/init` primitives

## Requirements

The `RisingEdge` operator take boolean input `s` and output a boolean `R` according to:

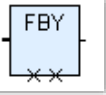
```
if ( $s_t$  and  $s_{t-1}$ )  
     $R_t = \text{true}$   
else  
     $R_t = \text{false}$ 
```

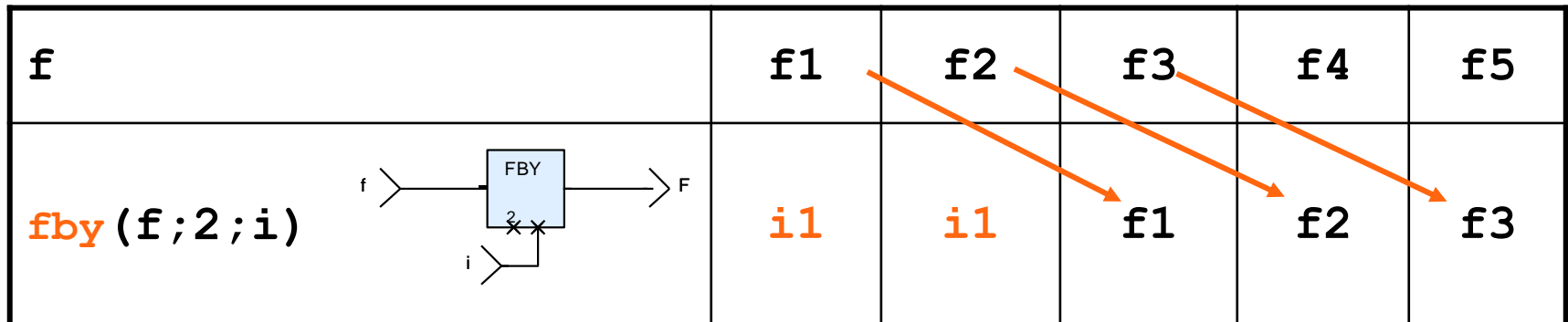
# Exercise 1: Solution



# Time Operators: fby

Delays its inputs by a certain number of cycles while providing a default value:

- Type interface: 'T, int, 'T → 'T, where T is any type
- Introduces a delay
- The number of past cycles must be a strictly positive integer value
- Example: 



- $\text{fby}(f;n;i)$  is equivalent to :  
 $i \rightarrow \text{pre}(i \rightarrow \text{pre}(\dots \rightarrow \text{pre}(f)))$

## Exercise 2: Rising Edge

### Objective

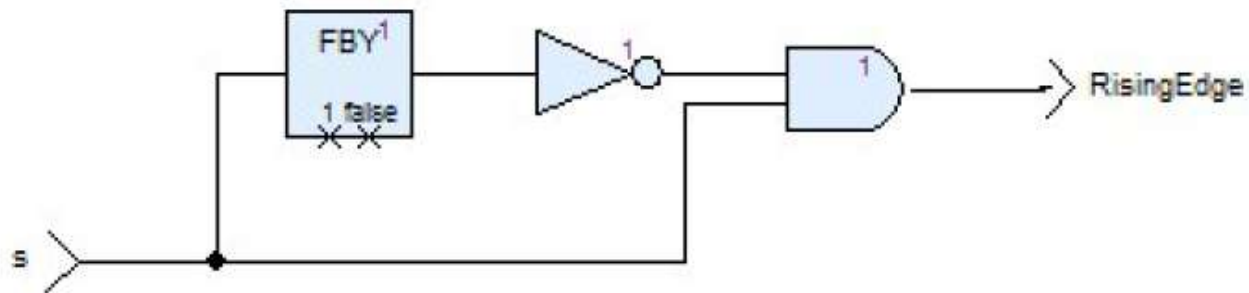
Implement a rising edge operator using `fb` primitives

### Requirements

The `RisingEdge` operator take boolean input `s` and output a boolean `R` according to:

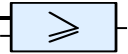
```
if ( $s_t$  and  $s_{t-1}$ )  
     $R_t = \text{true}$   
else  
     $R_t = \text{false}$ 
```

## Exercise 2: Solution

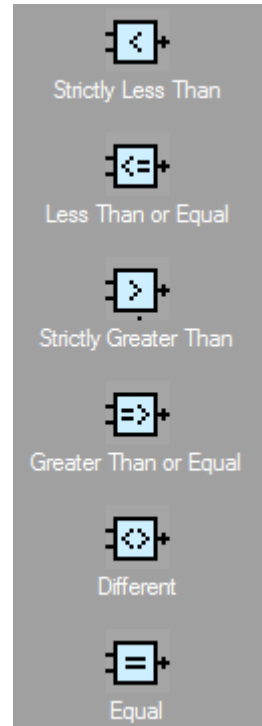


# Comparison Operators

They apply to integer or real types

- Type interface: 'N, 'N → bool, where 'N is a numeric type
- Example 

Cycle	1	2	3	4
a	-145	-1	154	-1456
b	-1	-2	154	1785
c = a >= b	false	true	true	false



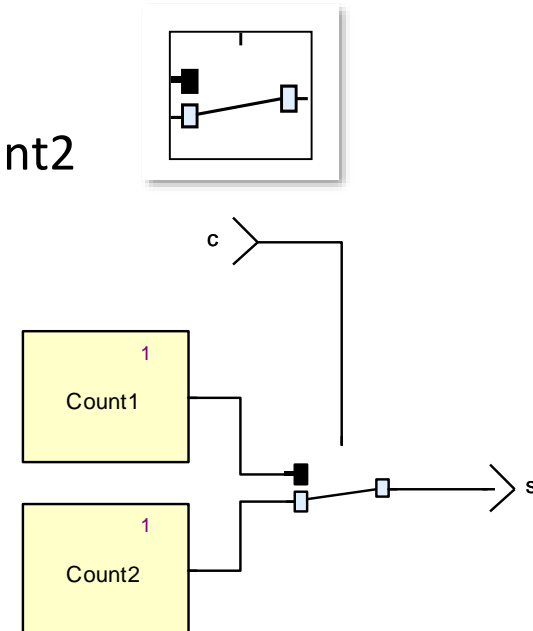
# Choice Operators: if-then-else (Reminder)

Receives two flows and outputs one of them depending on a Boolean condition

Both “then” and “else” expressions are always evaluated independently of the condition value.

## Example:

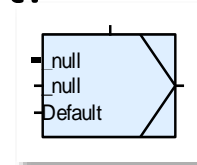
- Consider two counter operators count1 and count2
- $s = \text{if } c \text{ then Count1}() \text{ else Count2}();$ 
  - Each counter is incremented at each cycle regardless of  $c$
  - This behaviour complies with the intuition behind the reading of the graphical representation



# Choice Operators: case

Receives  $2N+1$  inputs and produces one output:

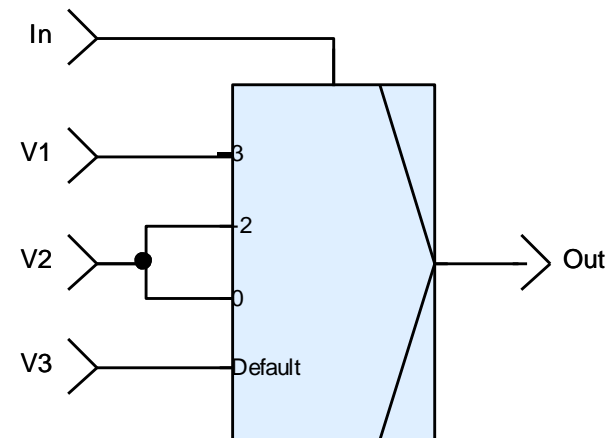
- $N$  inputs  $V_i$ , candidate value for the output
- $N$  constant inputs, one per  $V_i$  plus a default case
- 1 input to be tested against the constants



All the inputs are computed before the choice

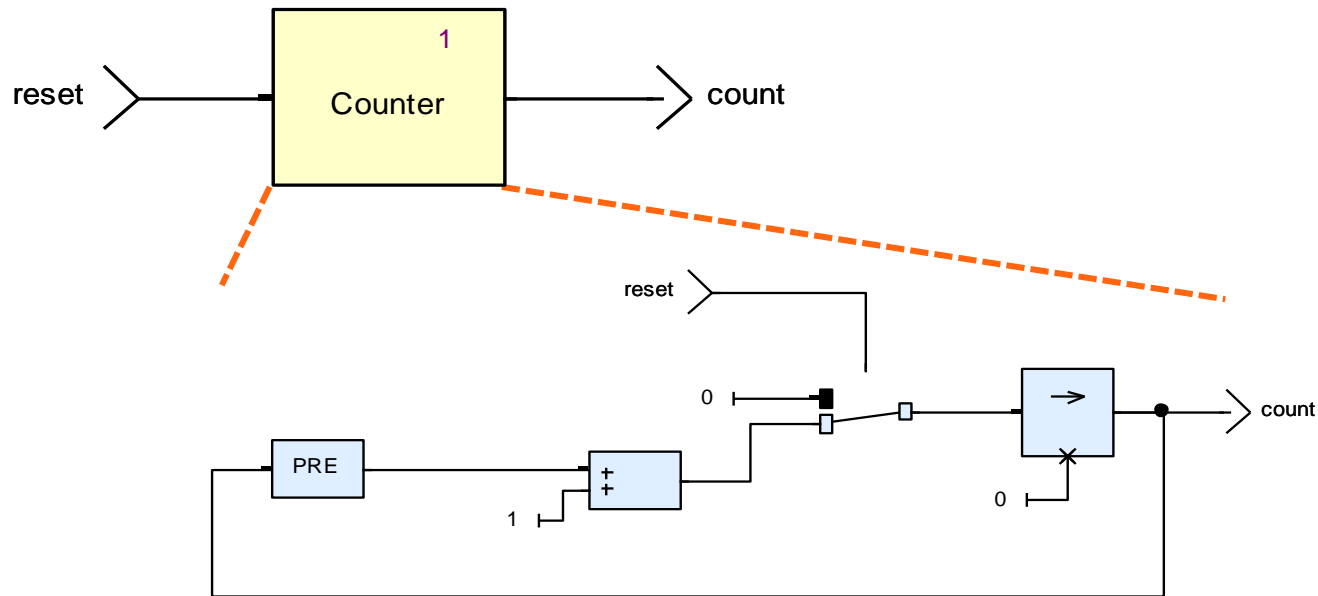
## Example:

- Out = case In of  
    3 : (V1)  
    -2 : (V2)  
    0 : (V2)  
    default : (V3);
- Equivalent to:  
    Out = if In=3 then V1  
        else (if In=-2 or 0 then V2  
              else V3)





# Choice Operators Example: Simple Counter



```
node Counter (reset: bool) returns (count: int32)
count = 0 -> if reset then 0 else PRE(count) + 1;
```

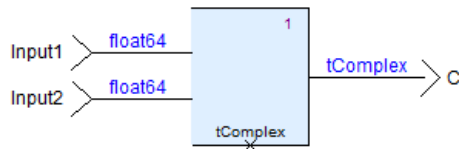
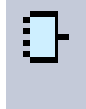
Cycle	1	2	3	4	5
count (reset =false)	0	1	2	3	4

# Structured Type Operators

## Operators on structures:

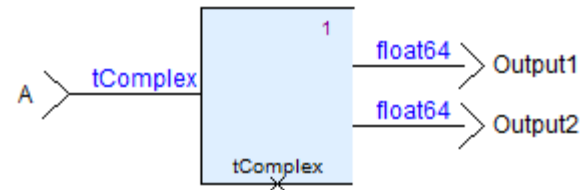
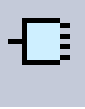
### Make

- CTRL+Drag&Drop on type
- Compose on output
- Properties/Parameters on Make

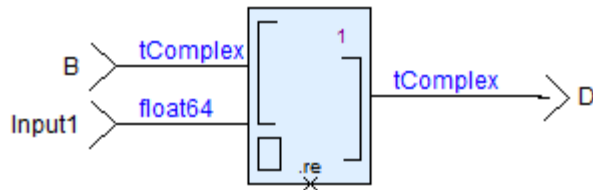
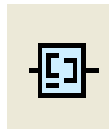


### Flatten

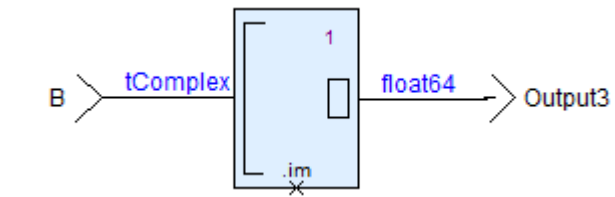
- Drag&Drop on type
- Decompose on input(s)
- Properties/Parameters on Flatten



### New Assign of a Structured Element



### New Projection (select a field)



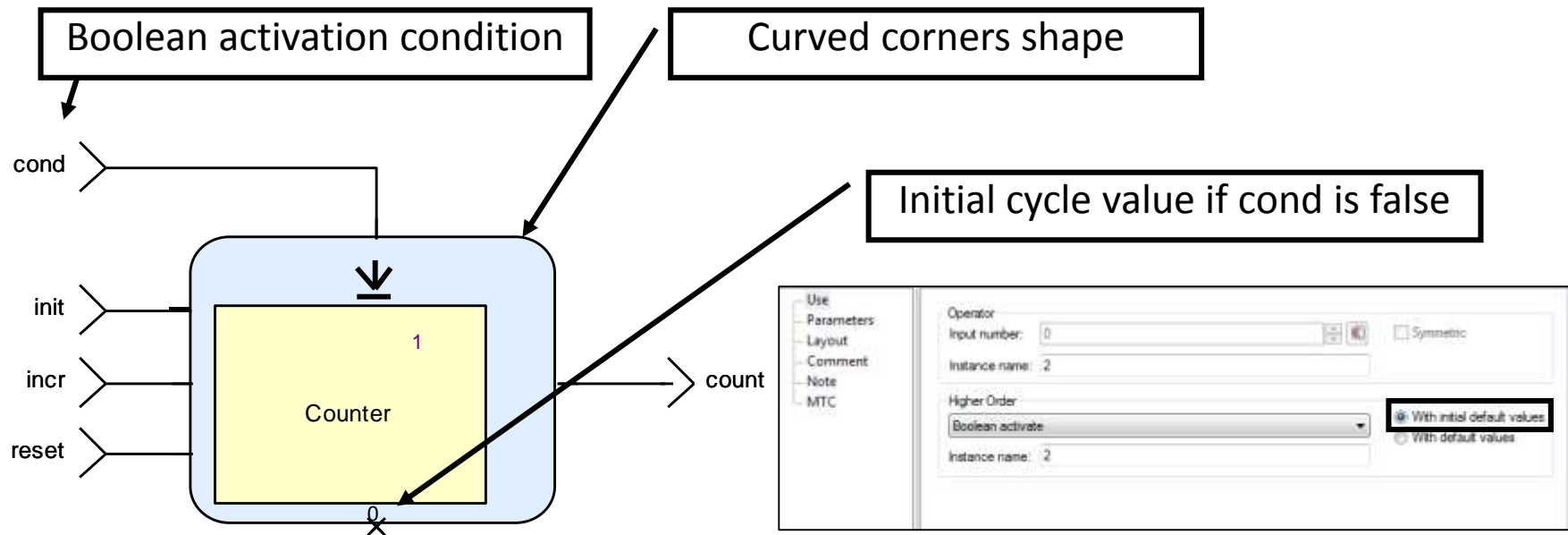
# Higher order: Activation Constructs

SCADE Suite provides constructs to control the activation of operators:

- Boolean Activate
- Restart

# Boolean Activate Operators: Memorized Outputs

Operator outputs are memorized when the activation condition is false



```
count =  
(activate Counter every cond initial default (0)) (reset, init, incr);
```

# Exercise 3: Memorized Output

## Objective

Implement boolean activate higher order

## Requirements

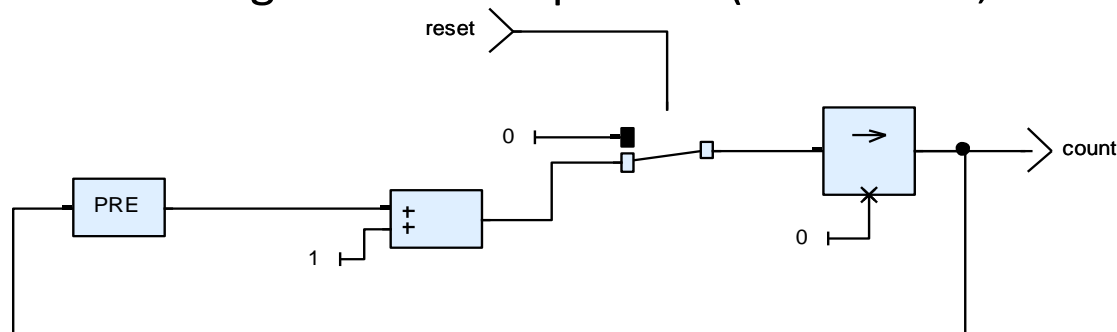
Create an operator which

```

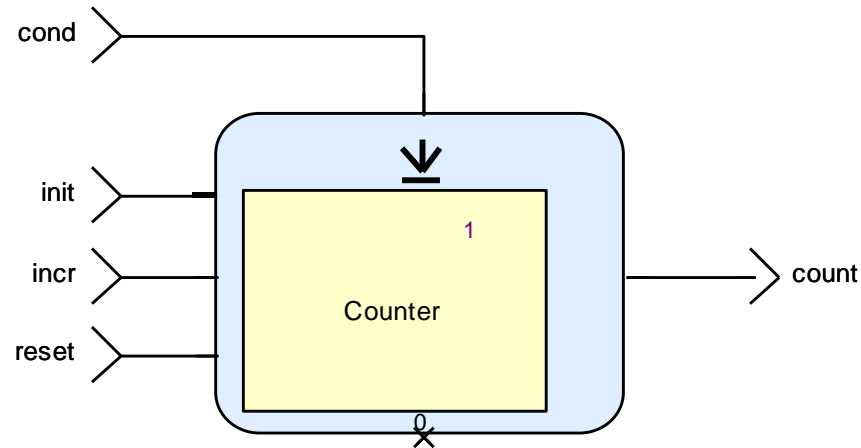
if enable = true
     $count_t = count_{t-1} + 1$ 
Else
     $count_t = count_{t-1}$ 
  
```

Name	Kind	Type
enable	input	bool
count	output	uint16

**Tip:** create the following Counter operator (reset: bool, count: uint16)



# Exercise 3: Solution

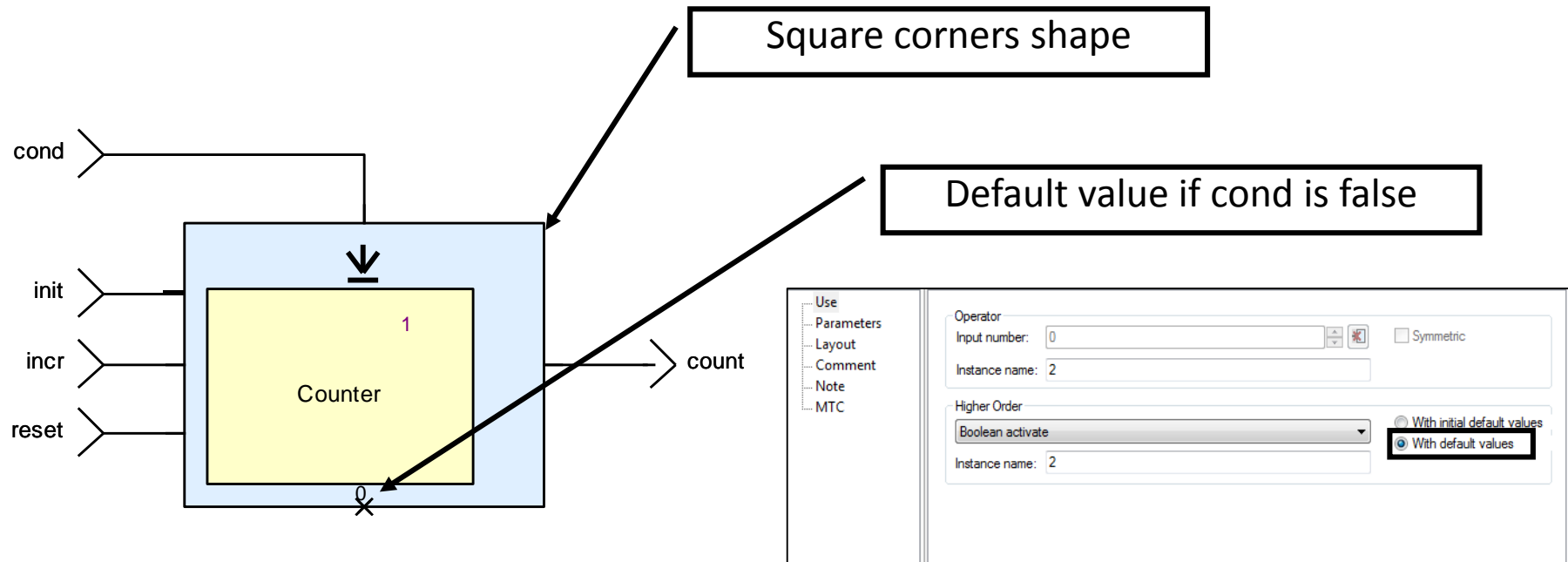


cond	true	false	true	false	false
Activate Counter	count1=Counter()	count1	count3=Counter()	count3	count3

cond	false	false	true	false	false
Activate Counter	0	0	count3=Counter()	count3	count3

# Boolean Activate Operators: Default Outputs

Operator outputs have a default value when the activation condition is false



```
count = (activate Counter every cond default (0)) (reset, init, incr);
```

# Exercise 4: Default Output

## Objective

Implement boolean activate higher order

## Requirements

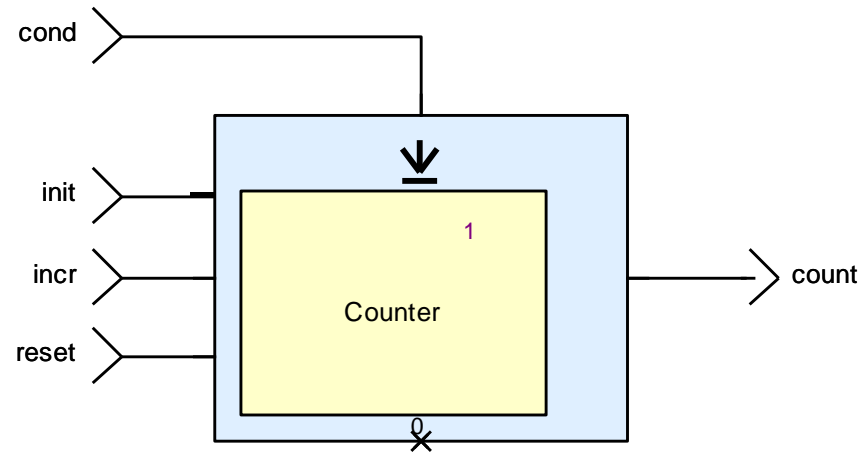
Create an operator which

```
if enable = true
     $count_t = count_{t-1} + 1$ 
Else
     $count_t = 0$ 
```

Name	Kind	Type
enable	input	bool
count	output	uint16



# Exercise 4: Solution

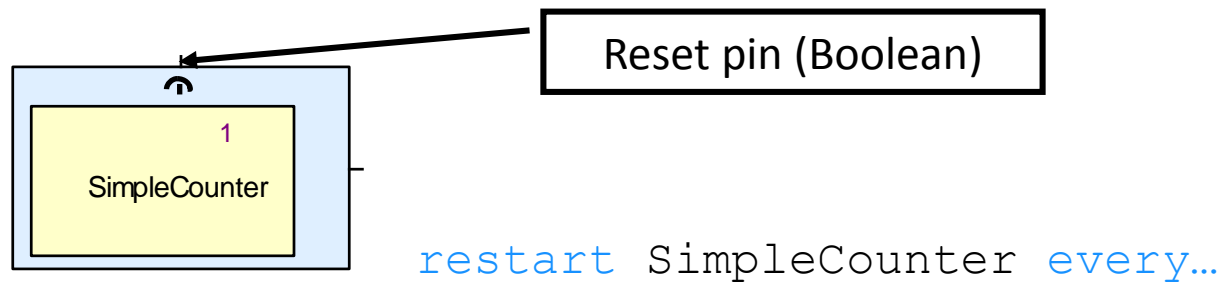


cond	true	false	true	false	false
activate Counter	count1=Counter()	0	count3=Counter()	0	0

cond	false	false	true	false	false
activate Counter	0	0	count3=Counter()	0	0

# Restart Operator

The restart operator applies on operators



When activated, resets the operator to its initial cycle

- Applies on the whole sub-hierarchy of the operator

The effect is the resetting of the memories of the operator, as Init, pre/Fby, State machines, ...

# Lab 3 (1/4)

Lab Support p.26-47

## Objective:

Design the `CruiseRegulation` operator

## Requirements:

Time: 30 min

Implement the `CruiseRegulation` function

*(requirements CC\_HLR\_CDC\_01 to 05)*

The `CruiseRegulation` operator manages automatically the vehicle speed when the Cruise Control is on: the regulation is reset when the Cruise Control is on, and frozen when the throttle output is saturated

## Lab 3 (2/4)

Create the SaturateThrottle operator (*requirement CC\_HLR\_CDC\_05*)

Name	Kind	Type
throttleIn	input	CarType::tPercent
throttleOut	output	CarType::tPercent
saturate	output	bool

```
A = if (throttleIn < ZERO_PERCENT)
    then ZERO_PERCENT
    else throttleIn;
B = if (throttleIn > REGUL_THROTTLE_MAX)
    then REGUL_THROTTLE_MAX
    else A;
throttleOut = B;
```

# Lab 3 (3/4)

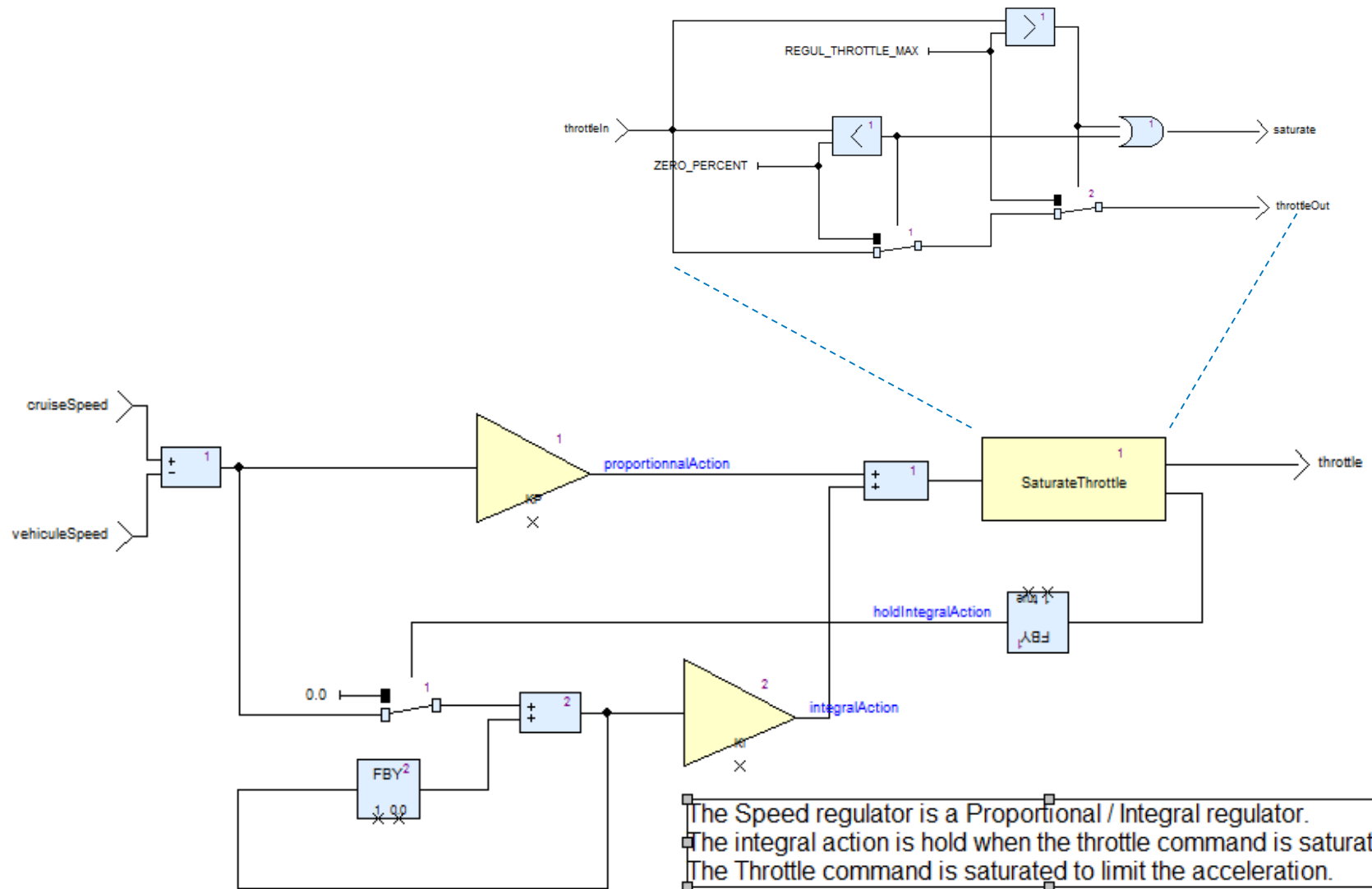
Create the CruiseRegulation operator (*requirement CC\_HLR\_CDC\_01 to 05*)

Name	Kind	Type
cruiseSpeed	input	CarType::tSpeed
vehicleSpeed	input	CarType::tSpeed
throttle	output	CarType::tPercent

- $$\text{Throttle} = K_p \cdot (\text{CruiseSpeed}_t - \text{VehicleSpeed}_t) + \sum_0^t (K_i \cdot (\text{CruiseSpeed}_t - \text{VehicleSpeed}_t))$$
- Saturate the throttle with SaturateThrottle operator
- $$\text{Saturate} = (\text{ThrottleIn} > \text{REGUL\_THROTTLE\_MAX}) \text{ or } (\text{ThrottleIn} < \text{ZERO\_PERCENT})$$
- $$\text{if saturate then Integral} = 0.0$$

$$\text{else Integral} = \text{cruiseSpeed}_t - \text{vehicleSpeed}_t$$
- $$\text{IntegralAction} = \sum_0^t K_i (\text{CruiseSpeed}_t - \text{VehicleSpeed}_t)$$

# Lab 3 (4/4): CruiseRegulation



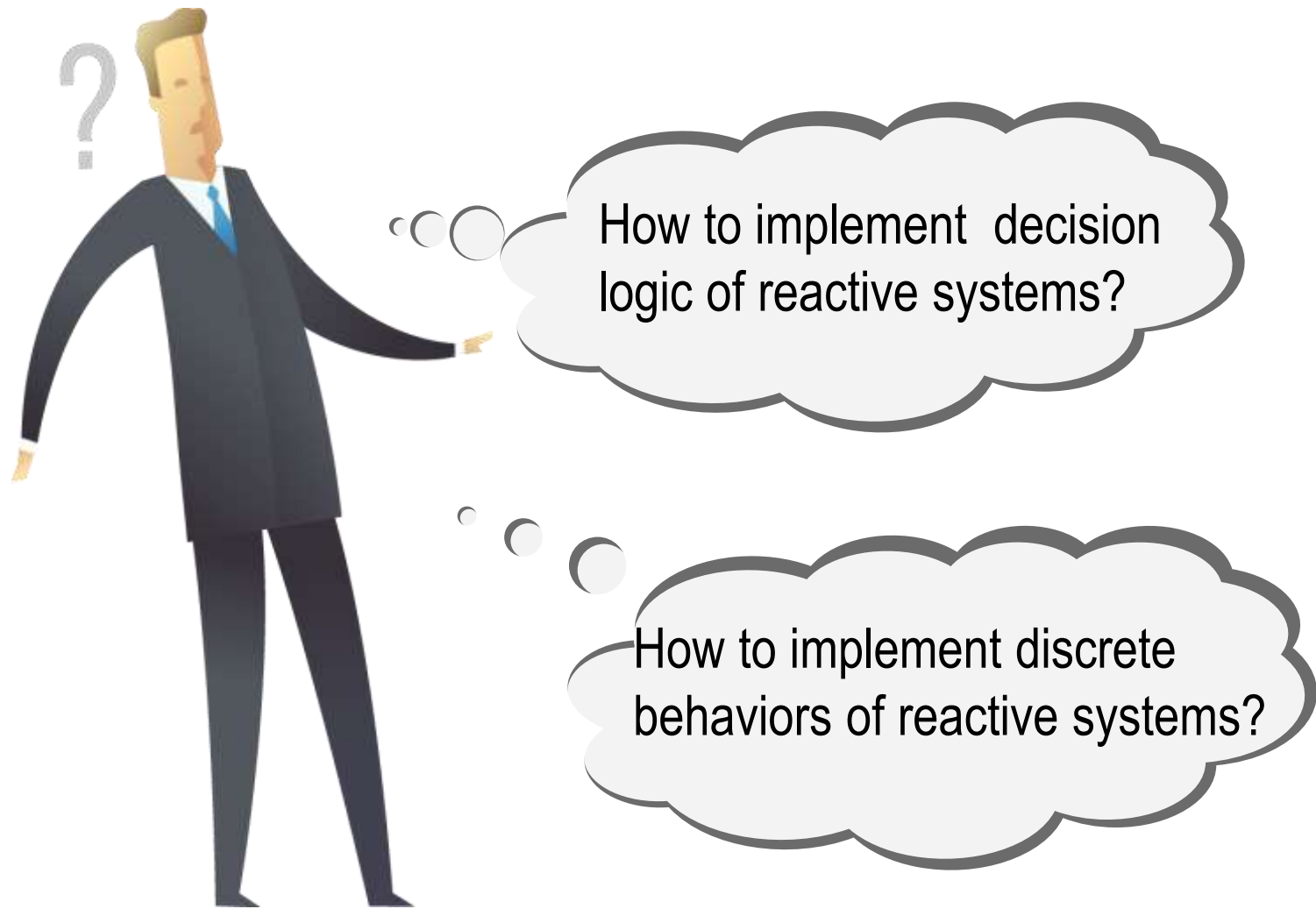
# AGENDA

Type declaration

Constant declaration

Operator and interface declaration

**Control flow declaration**



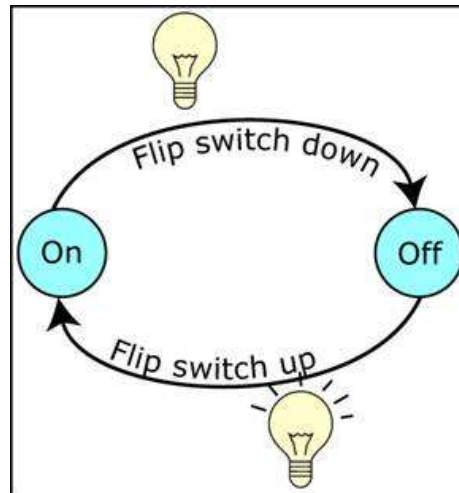


# Control Flows

Control flows are flows modelling discrete behaviors or decision logic of reactive systems

In SCADE Suite, find the following control flow constructs:

- Conditional blocks: If Block, When Block
- State machines:



# When Do You Use a Conditional Block?

A Conditional Block implements control flows to perform a test and then take action based on the test result

Because it allows a program to take a different path depending on some condition(s)

A Conditional Block has no memory so

- The generated code does not contain memories
- Use it when the conditions do not depend on the past and do not depend on the previous mode

# Applications for Conditional Blocks

## Flight-control, cockpit displays

- Modes, and user interaction control (buttons, navigation controls, etc.)



## Train speed control, interlocking, traffic lights, cruise control

- Modes, and user interaction control



## Vending machines, Automatic Teller Machine, elevators, etc.

- Modes and user interaction control



# Applications for Conditional Blocks

Used to implement the cruise speed management:

- CruiseSpeedMgt operator manages the value of the cruise speed according to the driver commands
- Conditional Block is used to fix, increase, reduce speed when the associated button is pressed

## Requirements:

CC\_HLR\_CSM\_01 to CC\_HLR\_CSM\_05

- Set or QuickAccel or QuickDecel buttons pressed when the CruiseControl is enabled

# When Do You Use a State Machine?

A State Machine implements control flows to perform a predetermined sequence of actions depending on a sequence of events or several conditions

Because

- it is in only one state at a time among a finite number of states
- it can change from one state to another when initiated by a triggering event or condition

A State Machine has memory so

- The generated code contains always memories
- Use it when the events or conditions depend on the past or the previous mode

# Applications for State Machines

## Flight-control, cockpit displays

- Display modes, user interaction control



## Train speed control, interlocking, traffic lights

- Modes, current states and transition logic



## Vending machines, Automatic Teller Machine, elevators, cruise control, etc.

- Modes, user interactions
- Hierarchy for the organization modes
- Concurrent events and behavior
- Synchronization of events
- Preemptive events



# Applications for State Machines

Used to implement the Cruise Control behavior:

- A State Machine is implemented to define the different states, ON, OFF, STDBY, etc.

## Requirements:

CC\_HLR\_CCB\_01

- When the driver starts the car, the CruiseControl shall be off
- The CruiseState output shall be set to OFF

CC\_HLR\_CCB\_02

- The CruiseControl shall be set on when the driver presses the on button

CC\_HLR\_CCB\_05

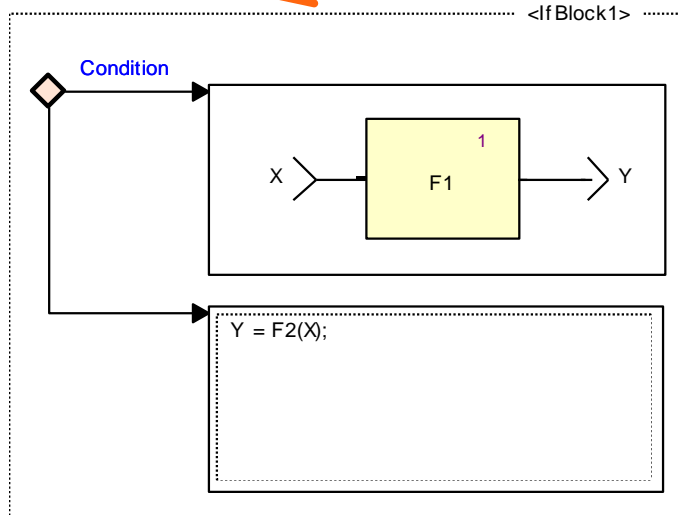
- The CruiseControl system shall be automatically disabled when the accelerator pedal is pressed, or when the car speed is outside the speed limit
- The CruiseState output shall be set to STDBY

# Conditional Blocks

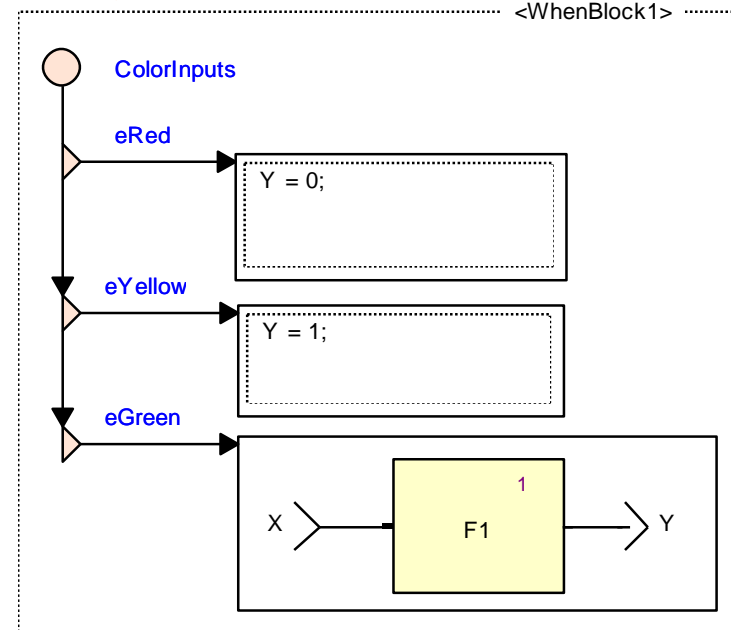
If and When Blocks are control flow constructs

Their activation is a function of conditions computable in the current cycle

## Boolean



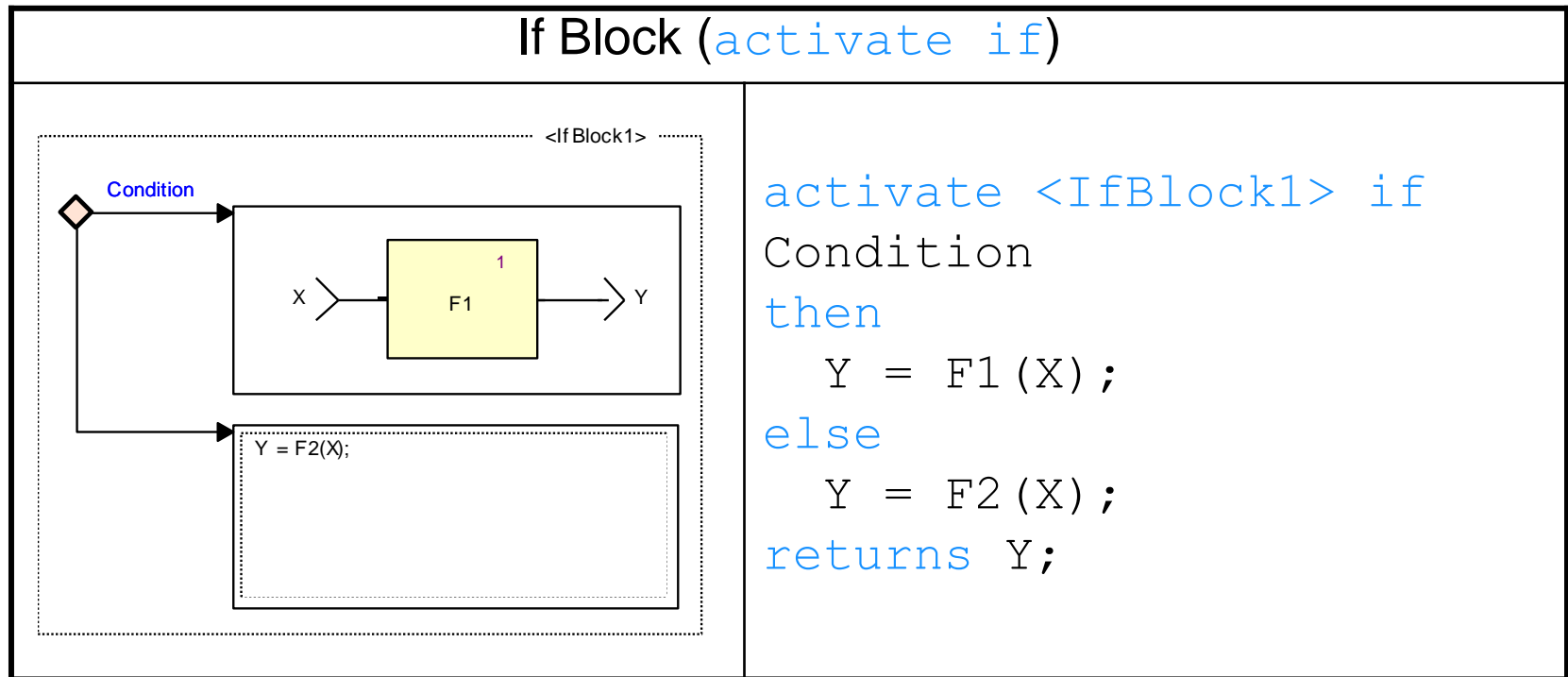
## Enumeration





# Conditional Blocks: If Block

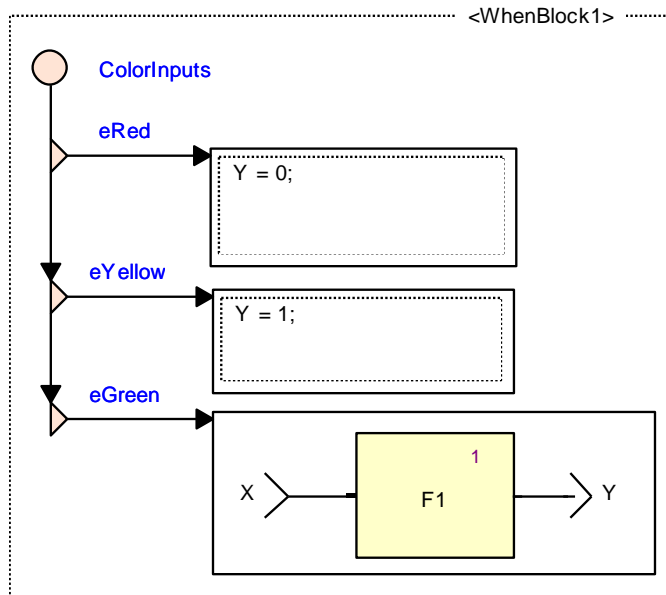
The control blocks implement any SCADE Suite equations, graphical or textual, including calls of user operators



# Conditional Blocks: When Block

Relies on an enumeration type

## When Block (activate when match)



```
activate <WhenBlock1> when  
ColorInputs match  
| eRed : Y = 0;  
| eYellow: Y = 1;  
| eGreen : Y = F(X) ;  
returns Y;
```

# Conditional Block: Data Flow Inside Block

Equations are calculated only when the block is active like an imperative-If and different from the If-Scade

Each declared flow (local or output) must have exactly one definition for each cycle where its block is active

What happens when a definition is missing in a given block?

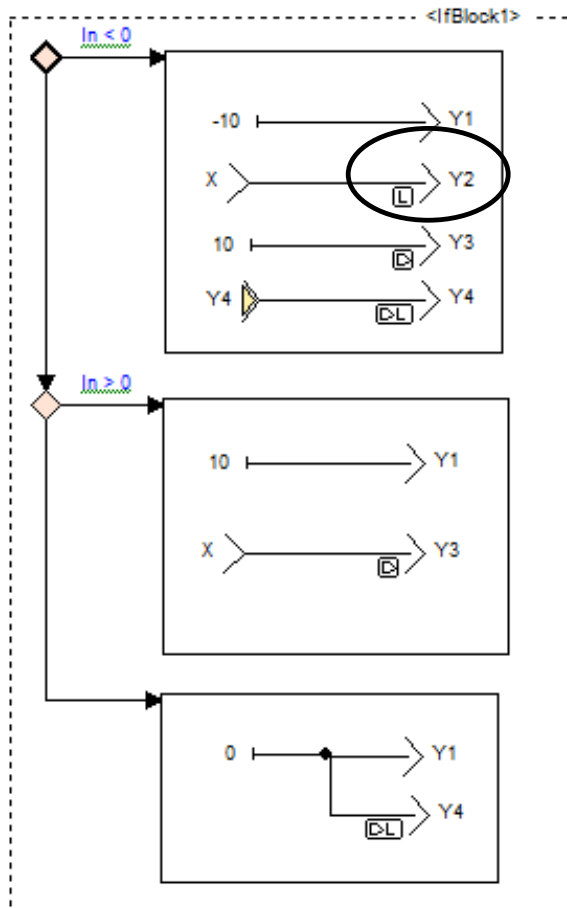
- Users can set either a **default value** or **maintain the last value of the flow**

The screenshot shows the configuration window for a block in ANSYS. On the left is a sidebar with a tree view containing the following items: General, Declaration (highlighted), Use, Clock, Layout, Comment, Note, KCG pragmas, and MTC. The main area on the right contains the following fields:

- Type: A dropdown menu with 'tSpeed' selected.
- Last: A text field containing 'Speed', followed by a dropdown arrow and a button with three dots.
- Default: A text field containing '0', followed by a dropdown arrow and a button with three dots.
- Kind: A dropdown menu with '→ Output' selected.

# Conditional Block: Data Flow Inside Block

Maintain the latest value: Y2



Y2 is produced  
with value X

*"maintain the latest value"*

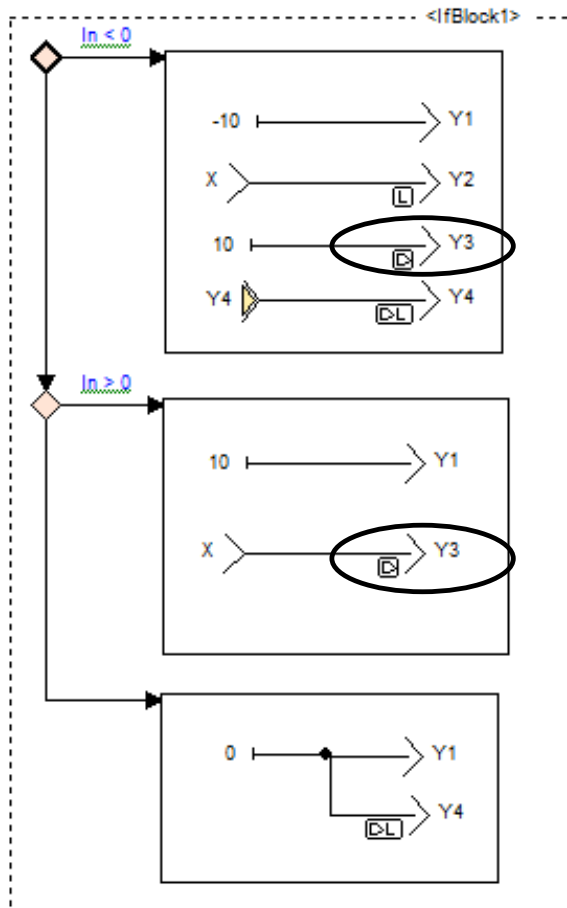
Y2 is produced  
with last value of Y2, or -5 at the  
first cycle

*"maintain the latest value"*

Y2 is produced  
with last value of Y2, or -5 at the  
first cycle

# Conditional Block: Data Flow Inside Block

Produce a default value: Y3



Y3 is produced  
with 10

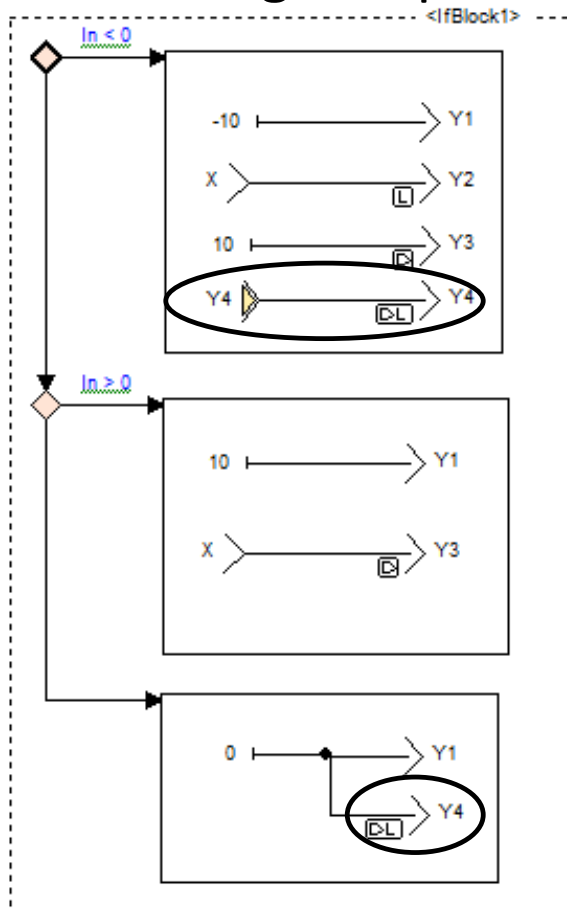
Y3 is produced  
with X

*"Produce a default value"*

Y3 is produced  
with 3

# Conditional Block: Data Flow Inside Block

If both Last and Default are defined, the Default value has the highest priority: Y4



Which value is produced for Y4?  
(see next slides)

*"Produce a default value"*

Y4 is produced  
with 3

Y4 is produced  
with 0

# Last Primitive

The last primitive is applied to an input/output/variable/signal to memorize and return its previous cycle value.

A flow has a single definition and must always be defined hence the need for last (and default) value to cover this rule.

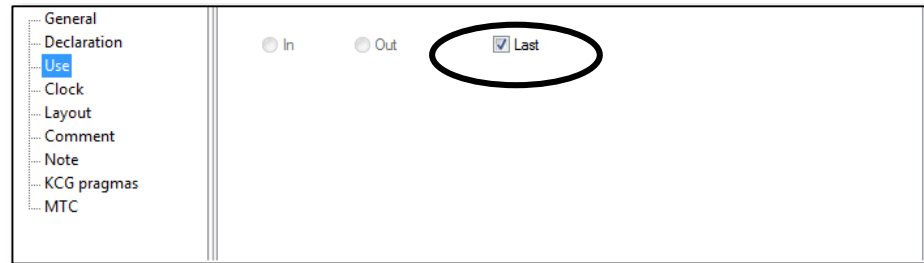
Differences Between last and pre:

- **pre** applies to an expression, **last** to an identifier (variables, inputs, outputs, signals)
- Same behavior outside control blocks
- Within a control block, pre provides the value at the previous cycle where the block was active, last provides the value at the previous cycle whether the block was active or not.

# Create Last Primitive

Drag & drop the variable and click Last in the “Use” tab

last 'AltitudeTarget

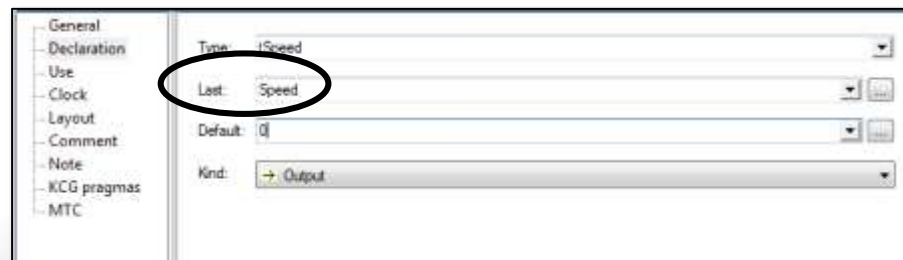


Use a textual expression of the form: last'VarName (Note the single quote in this expression)

last 'AltitudeTarget

A “last” variable must be defined on the first cycle.

Specify this initial value in the “Last” text box in the variable declaration tab





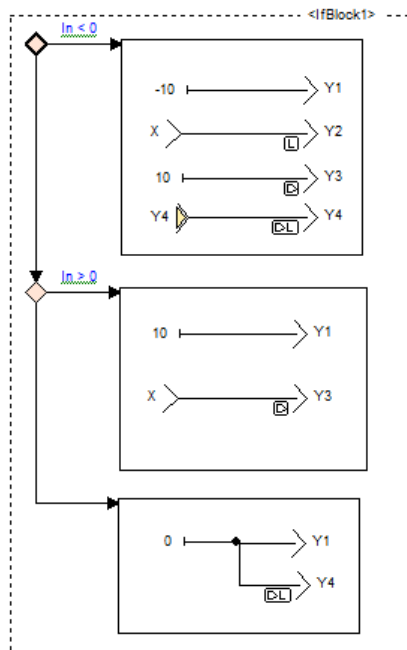
# Exercise 5: last primitive

## Objective

Observe the `last` primitive behavior

## Requirements

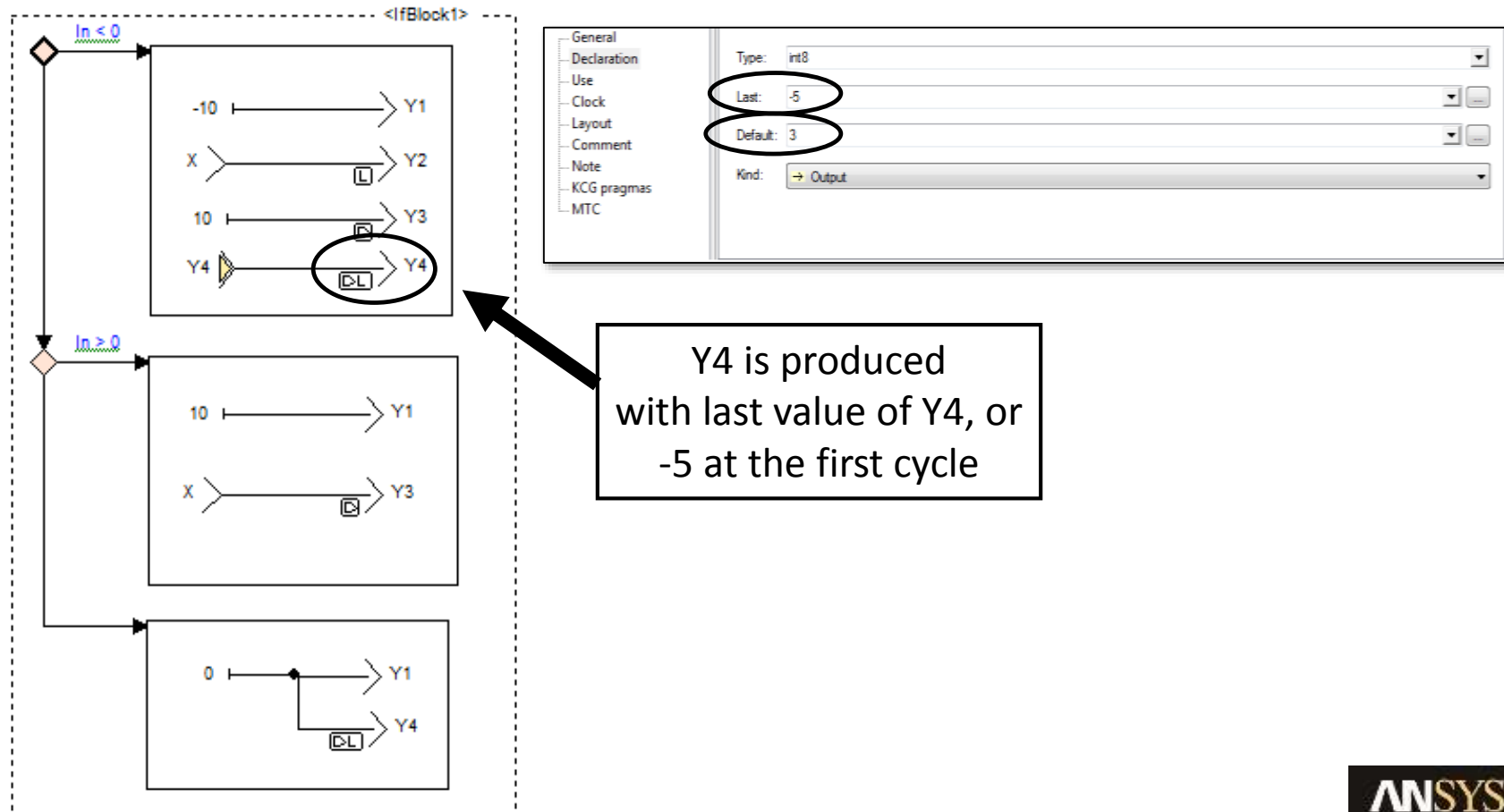
Create the following operator and simulate to observe the behavior



Name	Kind	Type	default/last
X	input	int8	-
Y1	output	int8	-
Y2	output	int8	default: -5
Y3	output	int8	last: 3
Y4	output	int8	default: -5 last: 3

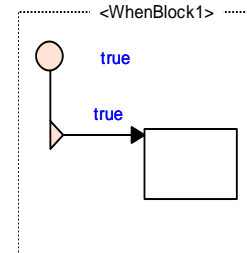
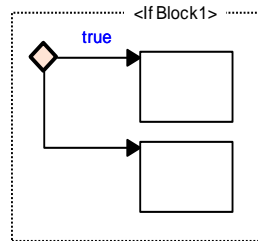
## Exercise 5: Solution



If both Last and Default are defined, the Default value has the highest priority: Y4

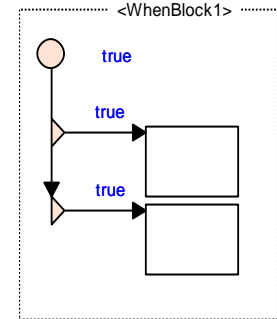
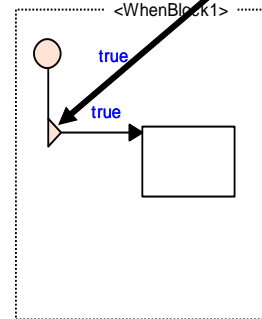
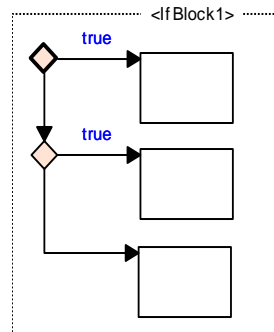
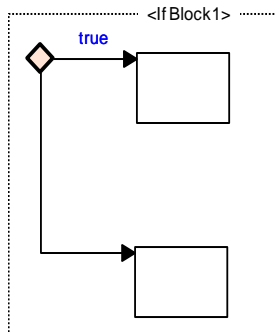


# Conditional Block: Creation

Click on  or  and draw a rectangle in the diagram to add an If or When Block.

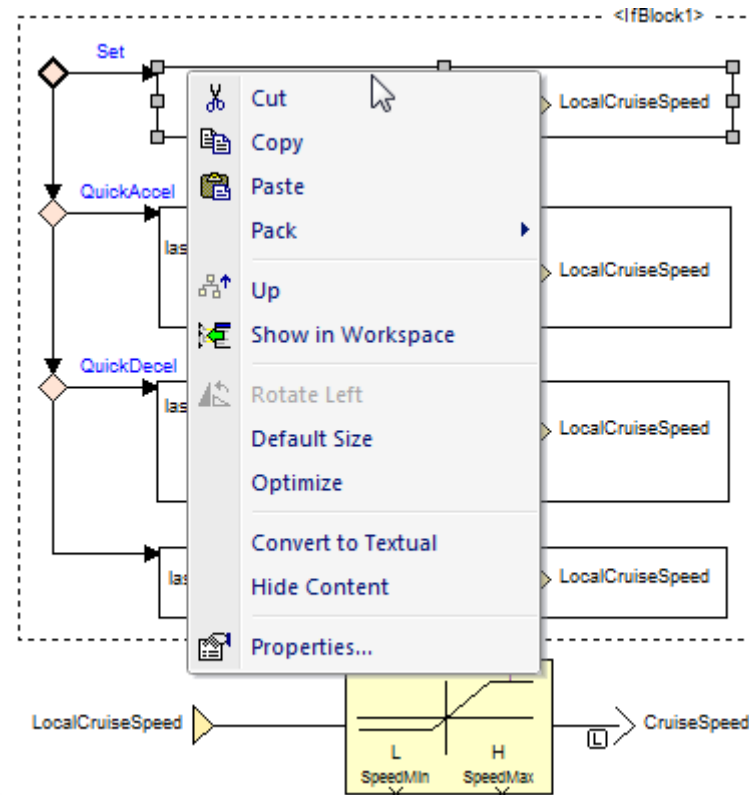


It is possible to add a branch by clicking on  (If) or  (When) and clicking on an else, or case, to add the new one.



# Block Activation: Textual Equations

Right click on the control block to enter textual equations and select "Convert to Textual"



# Lab 4 (1/3)

Use Lab Support p.49-63

## Objective:

Implement the `CruiseSpeedMgt` operator (**use of a conditional block**)

## Requirements:

Time: 20 min

The `CruiseSpeedMgt` operator manages the value of the cruise speed according to the driver's commands (Set or QuickAcccel or QuickDecel buttons pressed) when the Cruise Control is enabled  
(*requirements CC\_HLR\_CSM\_01 to 05*)

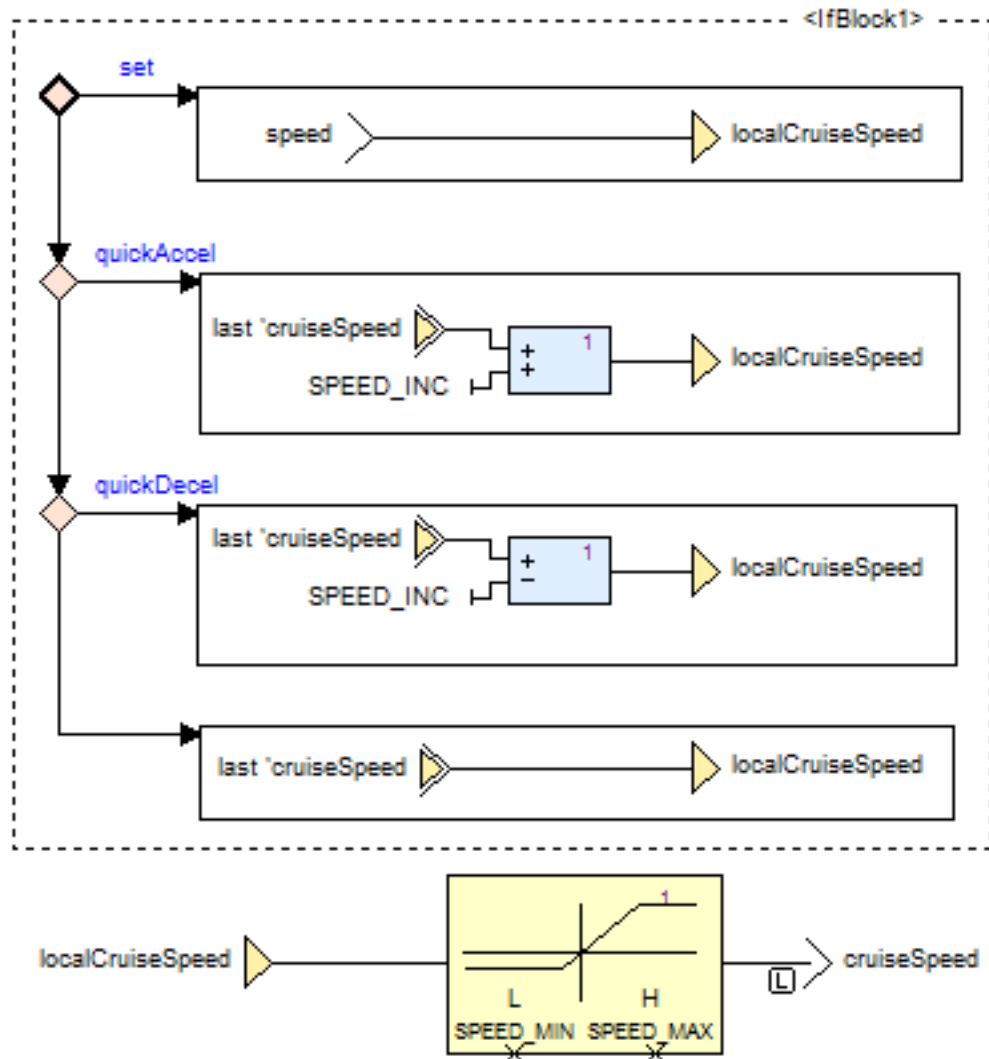
## Lab 4 (2/3)

In package `CruiseControl`, create `CruiseSpeedMgt` operator

Name	Kind	Type
Set	input	bool
QuickAccel	input	bool
QuickDecel	input	bool
Speed	input	<code>CarType::tSpeed</code>
CruiseSpeed	output	<code>CarType::tSpeed</code>

- Set the cruise speed to the current speed when the Set Button is pressed
- Increase the cruise speed of `SpeedInc` when the QuickAccel button is pressed
- Decrease the cruise speed of `SpeedInc` when the QuickDecel button is pressed
- Maintain the cruise speed to the last value (the default value of Last is Speed) when no button is pressed
- The cruise speed is well maintained between `SpeedMin` and `SpeedMax` km/h (use `pwlinear::LimiterUnSymmetrical` library operator)

# Lab 4 (3/3): CruiseSpeedMgt

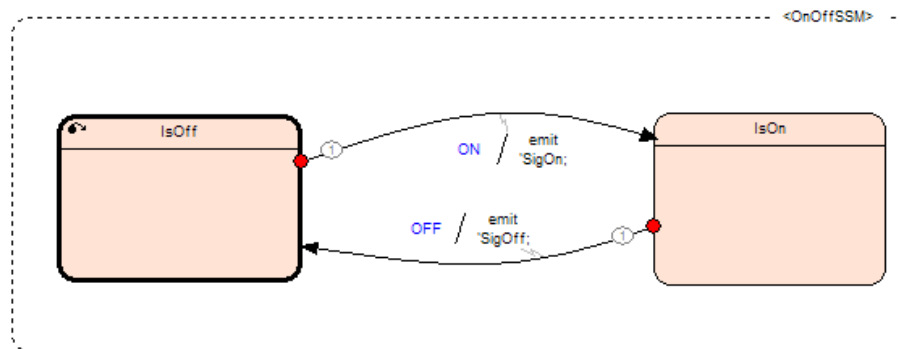


# State Machines

State machines are models describing the behavior of a software

Composed of a finite number of states and transitions

Considered as Control flow constructs where their activation is a function of the previous state

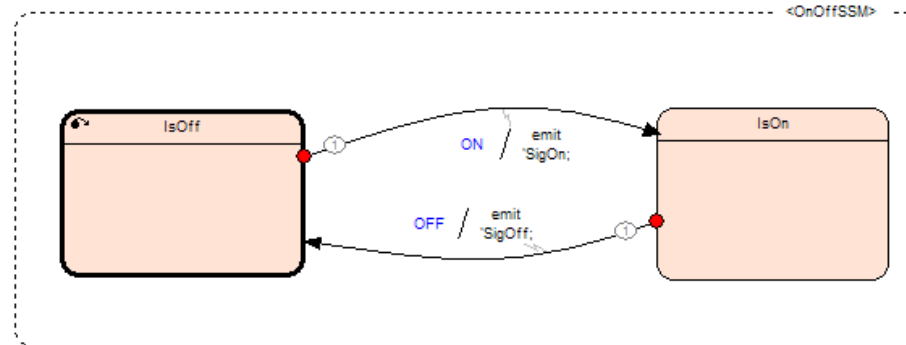




# State Machines

States are represented by boxes

Transitions are represented by arcs



States and transitions have labels for names and behaviors

# State Machines

A SCADE State machine has following properties:

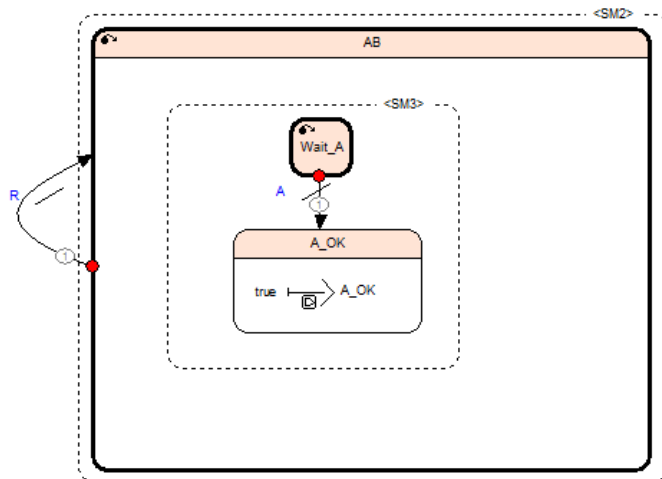
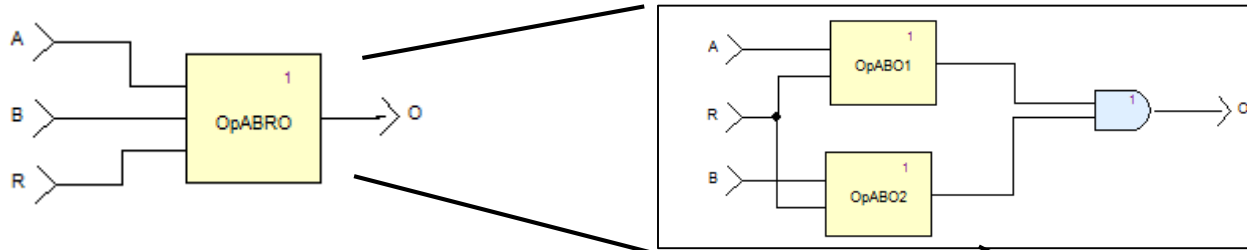
- Hierarchy for state architecture
- Parallelism for capturing the concurrency
- Synchronization of parallel state machines
- Preemption (transition) for handling errors or exceptions
- Smooth and consistent mix with data-flow design

=> Factorization of the model architecture: write things once and better readability

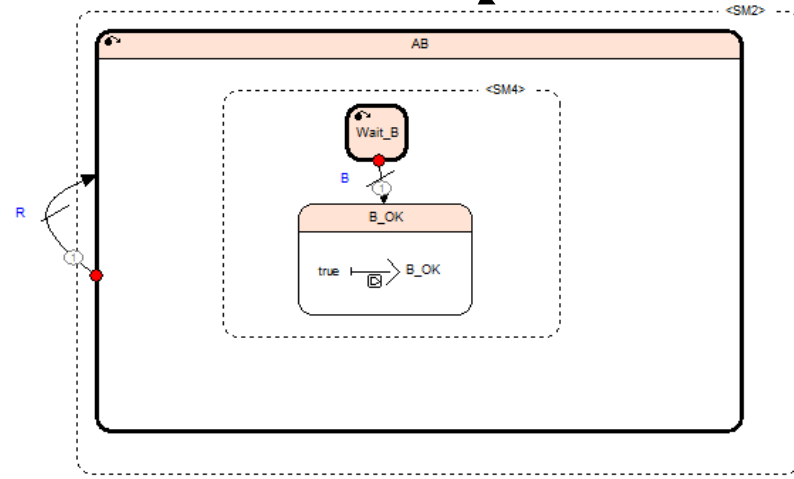
# Design without Parallelism

O is set to true when A and B are set to true and associated states are reached

AB state is reset when R is set to true



OpABO1 node



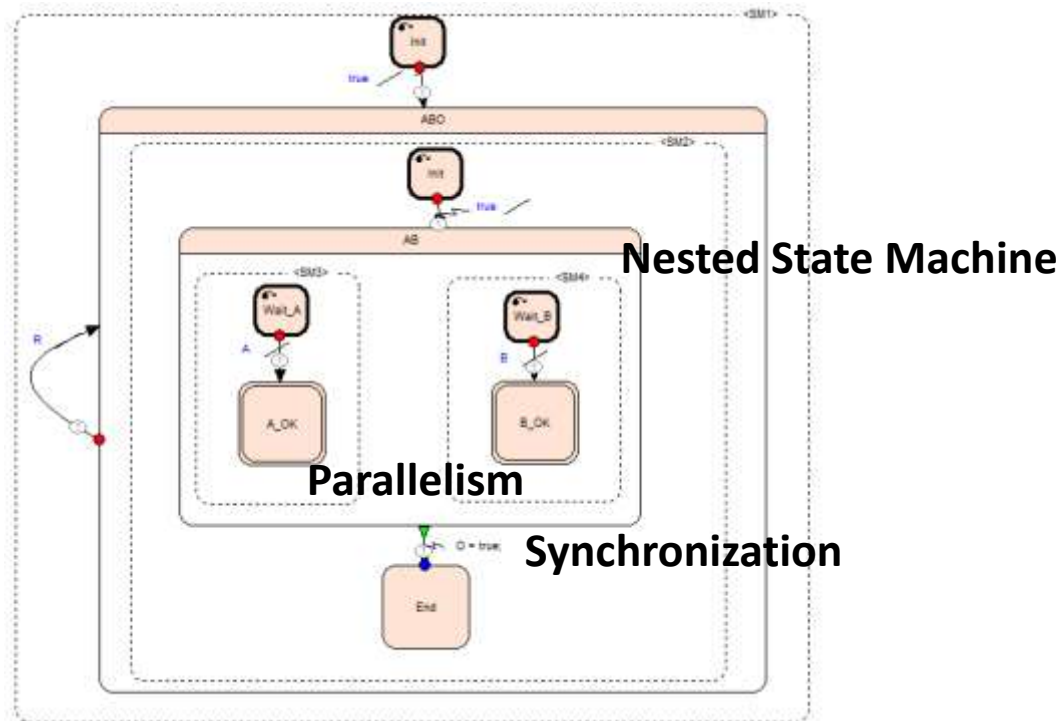
OpABO2 node

# Design with Parallelism

O is set to true when A and B are set to true and associated states are reached

- Synchronization of parallel States Machines

ABO state is reset when R is set to true



# State Machine Definition

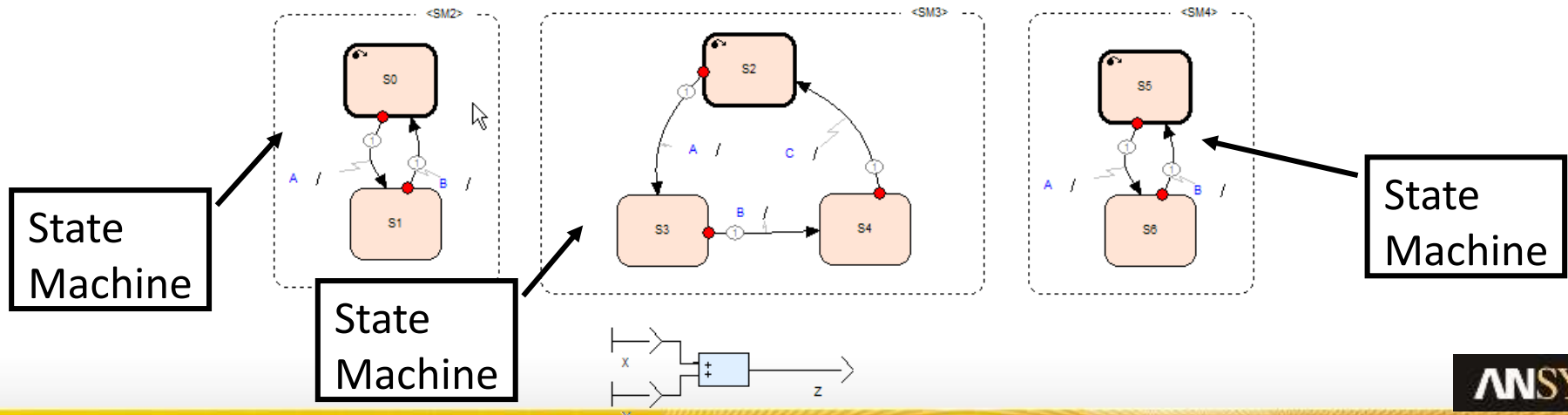
Created within a SCADE Suite operator:

- Only contains states and transitions
- Can own several parallel State Machines, each containing one or more states

Has a unique initial state:

- When another state is set as initial, any existing initial state is reset

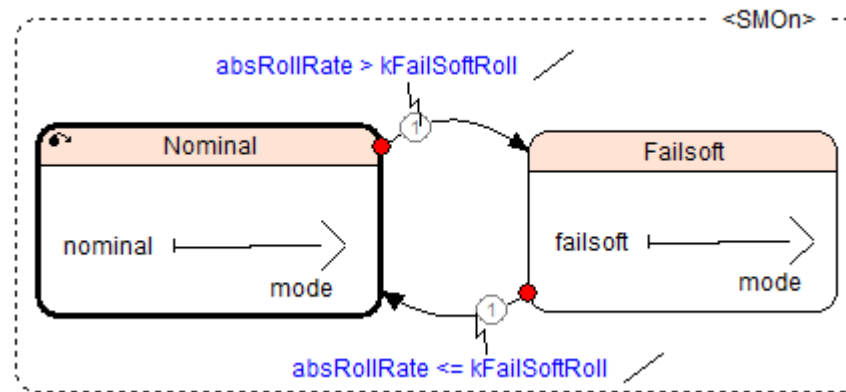
Surrounded by dashed lines to delimit it from the other State Machines or data-flow behaving in parallel



# State Machine States

## States:

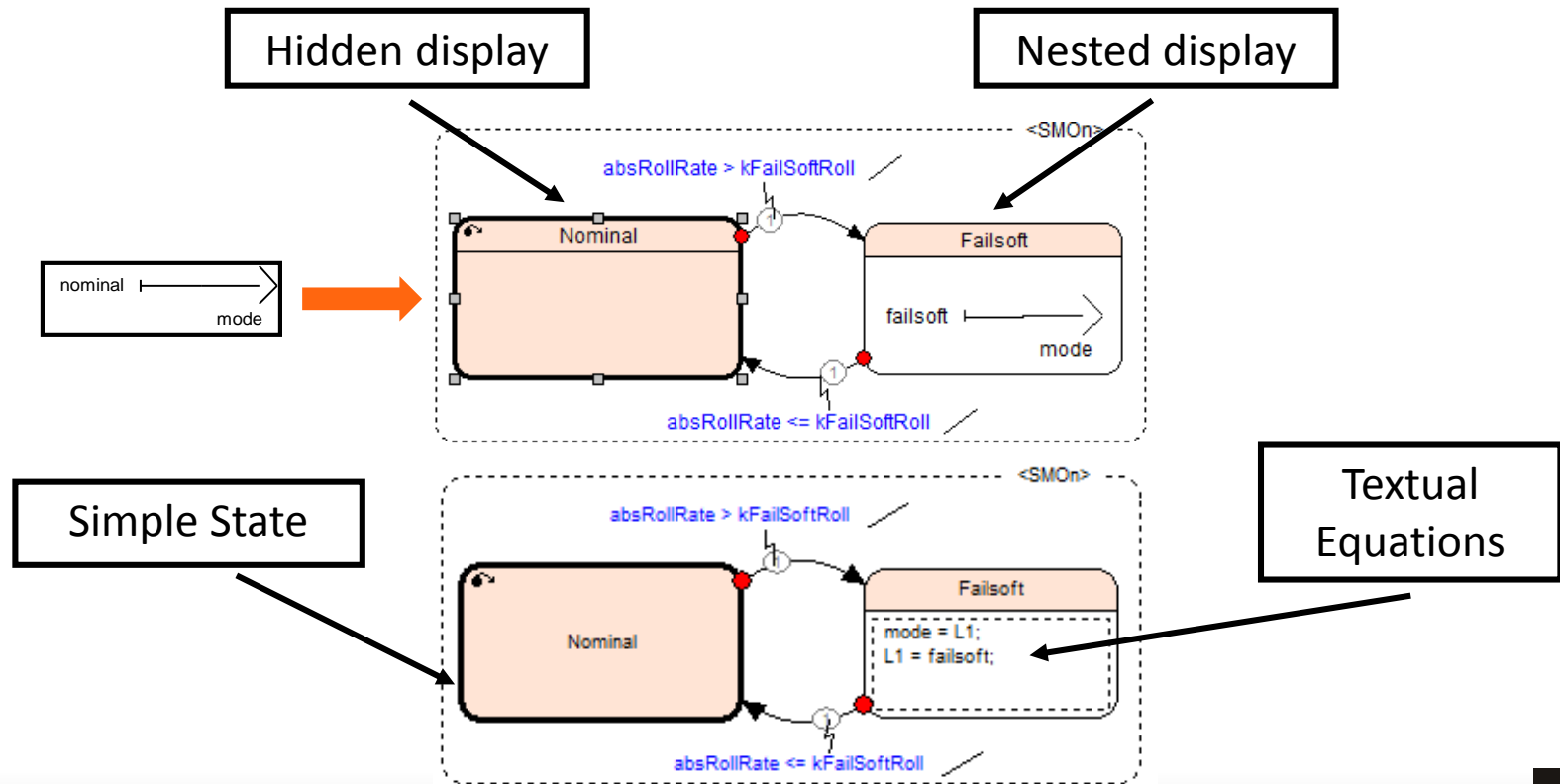
- Graphical rectangle that must be named
- At each cycle, a state is either active or not
- Only one state can be active/executed in a State Machine during a cycle



# State Machine States

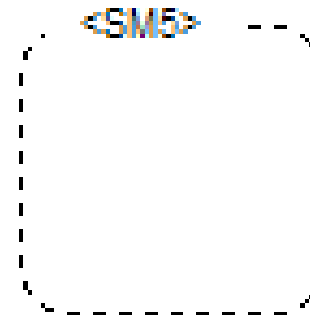
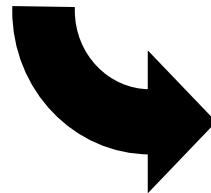
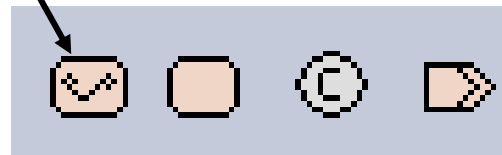
States may contain nested State Machine and graphical or textual data flows

Easy way to model control structures and state hierarchy



# Creating a State Machine

Click on the “New State Machine” to create an empty State Machine





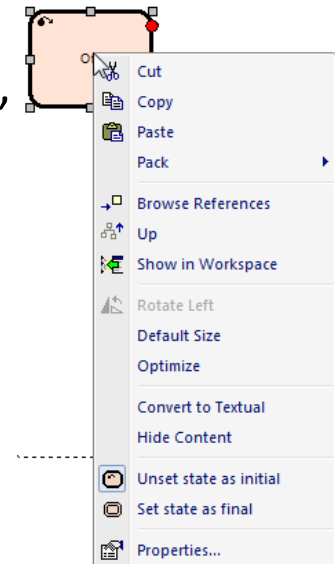
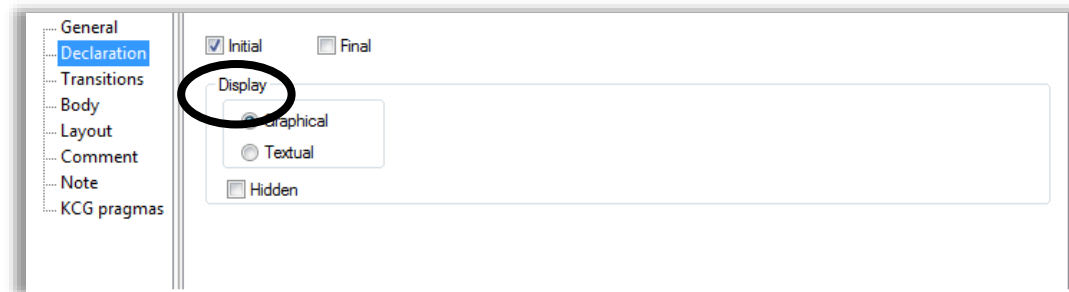
# Initial State

Graphically depicted with a bolder outline

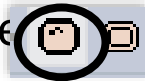
The initial state is active whenever the State Machine is started

How to set a state as initial?

- From its contextual menu “Set/Unset state as initial”
- From its Properties window “Initial” checkbox



- From the State Machine toolbar dedicated button

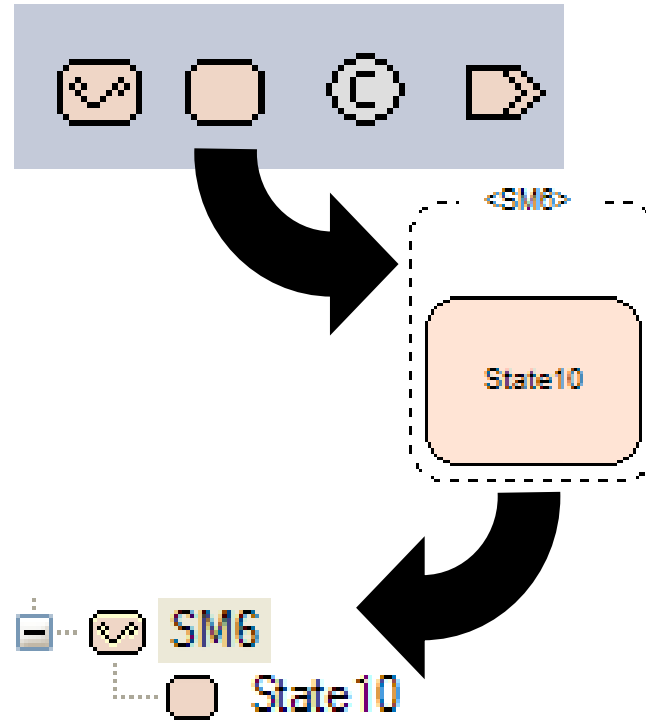


When a state is added to a State Machine, with no initial state, the state becomes initial

# State Editing: Creating States

## Creating states:

- Click the New State button
  - Click into state machine
  - Click & drag to custom its size



States are visible in the Scade View under their State Machine entries

# State Editing: State Properties

General

Declaration

Transitions

Body

Layout

Comment

Note

KCG pragmas

Name: Enabled

Path: CruiseControl::CruiseControl/SM1:Enabled:

Filename: CruiseControl1.xscade

☐ Separate File Name

Visibility

☒ Public ☐ Private

Name: a SCADE Suite identifier: Enabled ...

General

Declaration

Transitions

Body

Layout

Comment

Note

KCG pragmas

☒ Initial ☐ Final

Display

☒ Graphical ☐ Textual ☐ Hidden

Mark state as initial or final, graphical, textual and toggle hidden

General

Declaration

Transitions

Body

Layout

Comment

Note

KCG pragmas

Target States: ↑ ↓

Condition:

StandBy - (1)

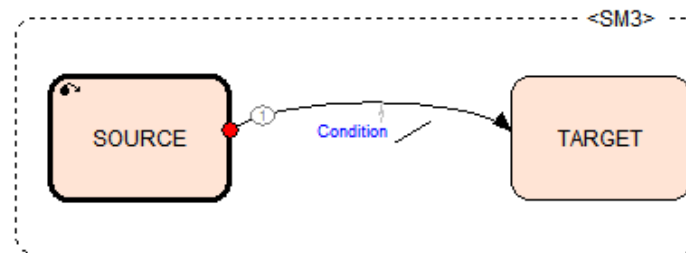
StdbyCondition

Outgoing transitions with priority and respective textual Scade code

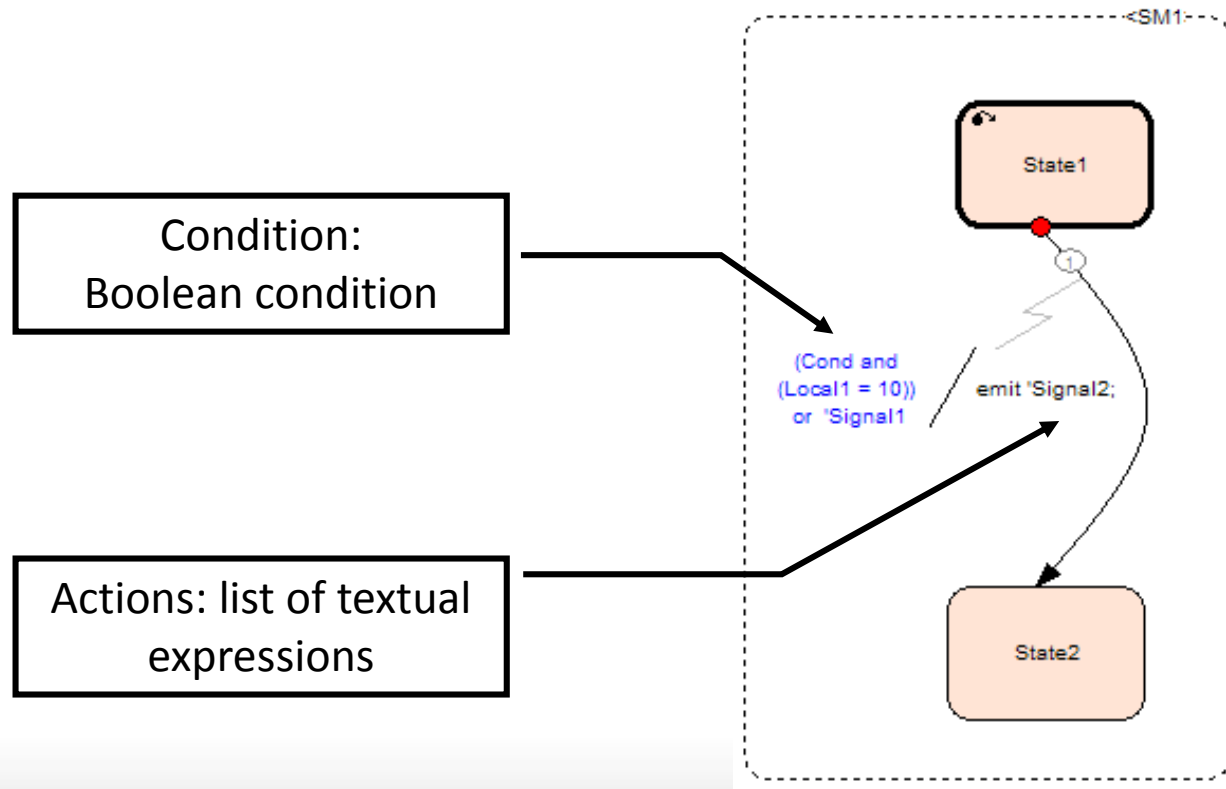
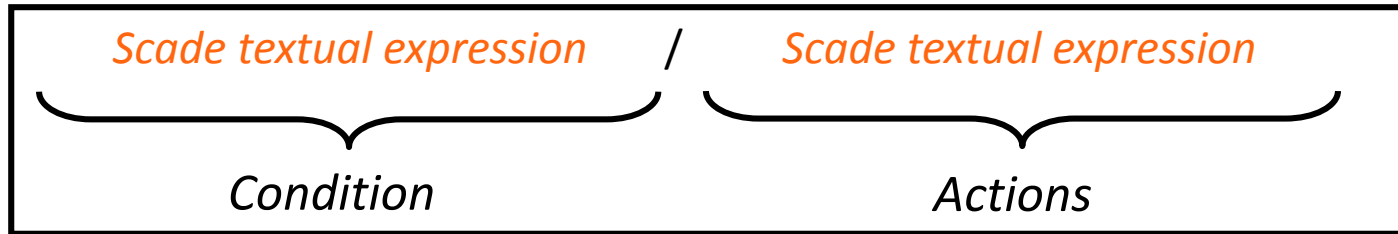
# State Machine Transitions

Have one source state and one target state

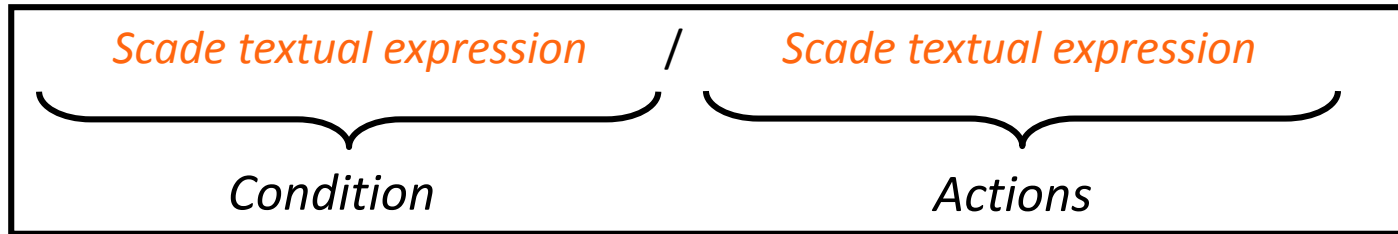
Label of the form **Condition / Actions**



# Transition Labeling



# Transition Labeling



## Condition:

- Boolean Scade expression
- Useful operators: times and last

```
2 times On
3 times (Local1 > 18.0)
5 times 'Signal1
last 'On
last 'Local1 > 18.0
last 'Signal1
2 times (last 'Local1 > 18.0)
```

# Transition Labeling



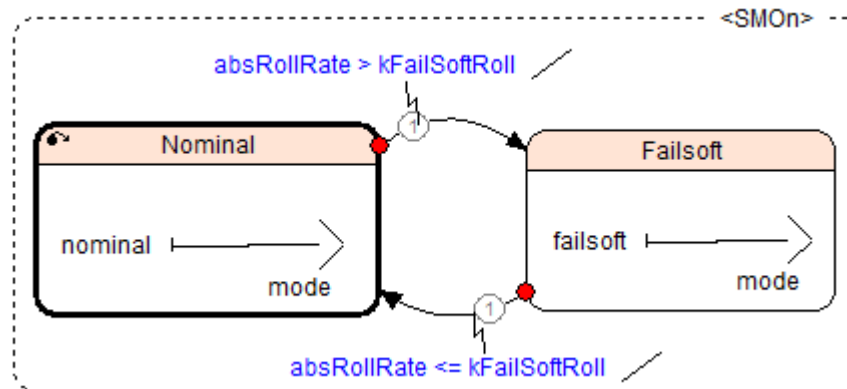
## Actions:

List of signal emissions and variable definitions, which can contain any Scade expression:

```
emit 'Signal1;  
Local1 = 15.0;  
StopEngine = true;  
MaxValue = Max(Value1, Value2);
```

# Transition Firing

A transition can be fired if its trigger condition is true and if its source state was active at the previous cycle:



- If  $\text{absRollRate} > \text{kFailSoftRoll}$  is true and if the Nominal state was active at the previous cycle, this state is preempted and state Failsoft is activated

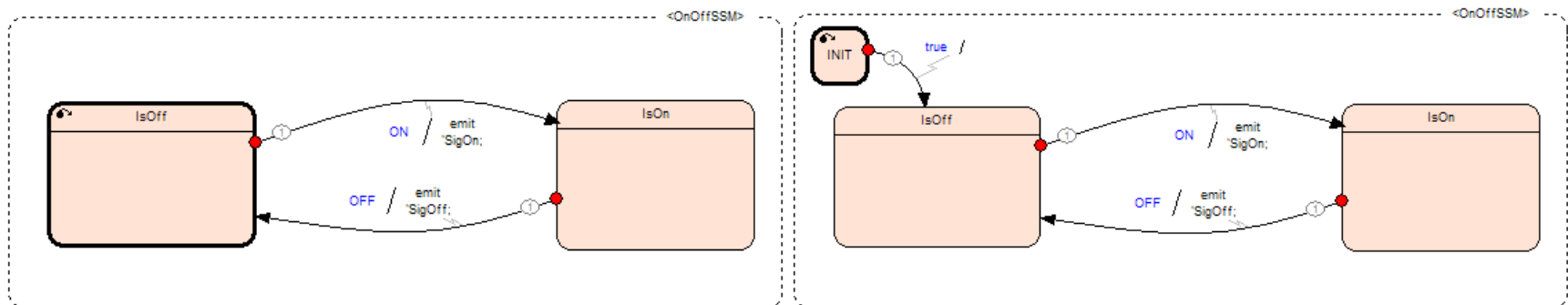
At most one transition is fired in a cycle (per State Machine).



# Transitions and Initial State

At the initial cycle, the transitions of the initial state are immediately fire-able.

At the other cycles, the initial state behaves like any other state.

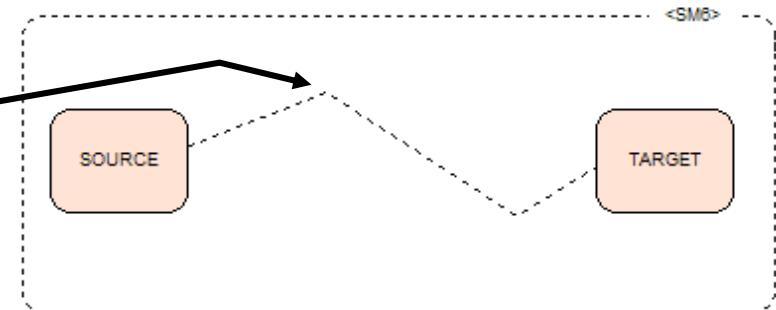


# Transition Editing: Creating a Transition

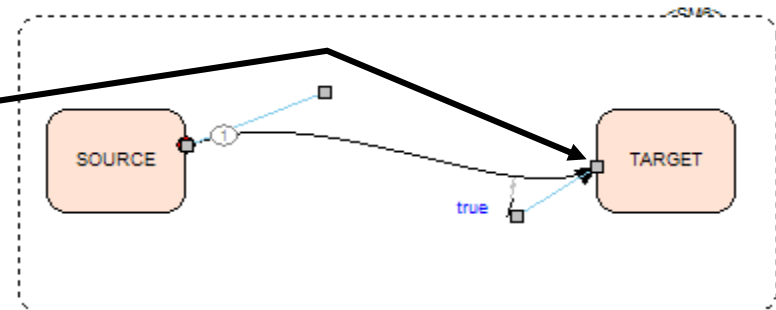
1. Position the mouse cursor over the edge of the source state and click the left mouse button



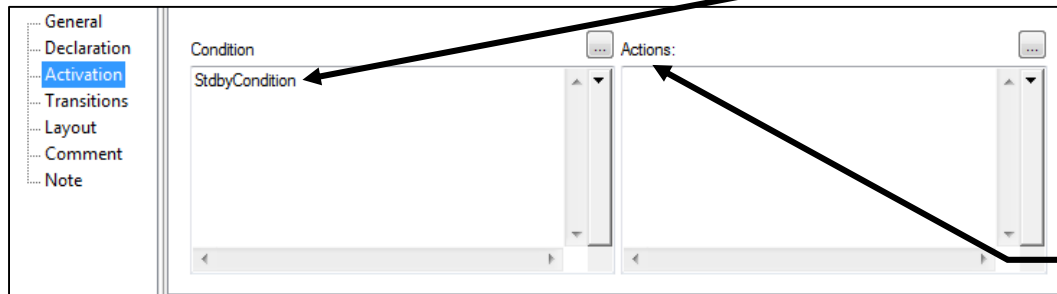
2. Click as many times as needed for a path to the border of the target state



3. Click on the border of the target state to create the final transition

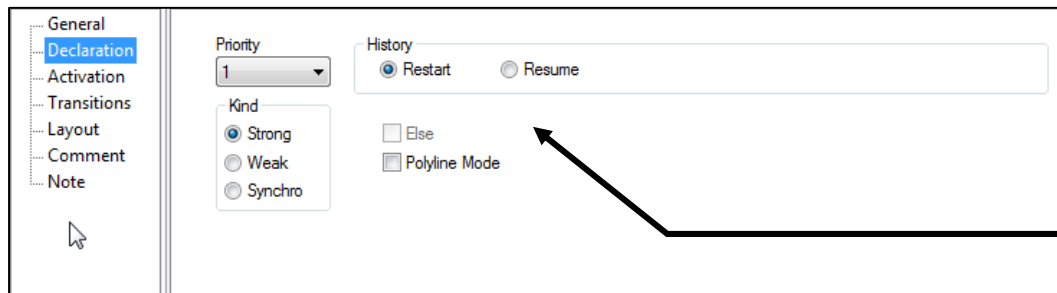


# Transition Editing: Transition Properties



Condition: Scade expression that triggers the transition when true

Actions: Scade expression that describes actions to be taken when the transition is triggered

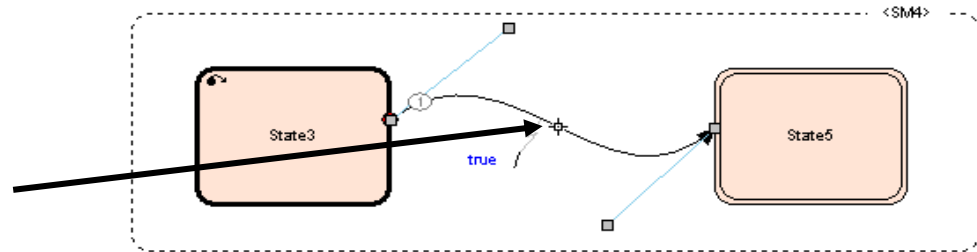


Priority: ensures determinism  
Kind: preemption type  
History: resume states  
Polyline: straight lines

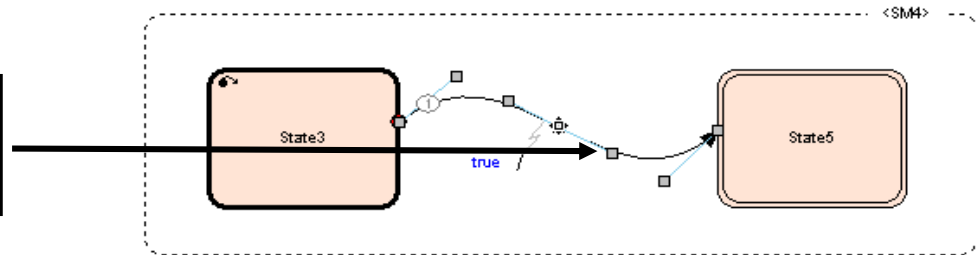
# Transition Editing: Modifying Layout

## Additional control vectors

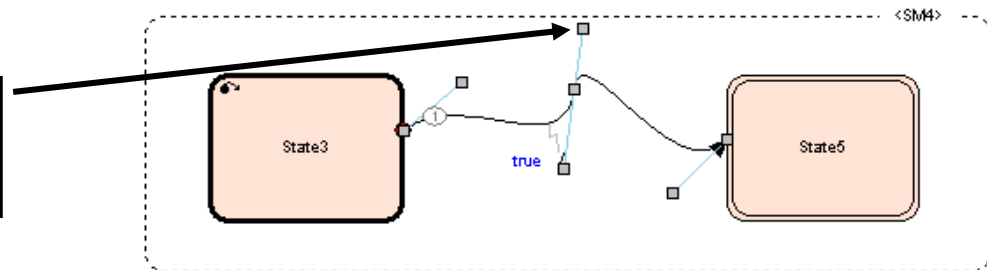
1. Position mouse cursor over a selected transition and holding down CTRL key



2. Keep CTRL pressed and click to add control vector



3. Move control points to modify transition

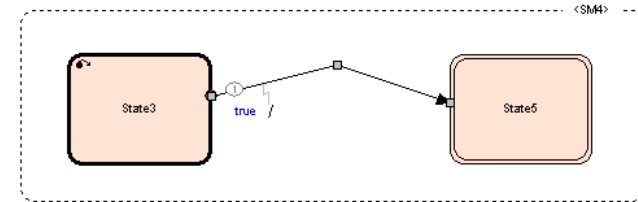
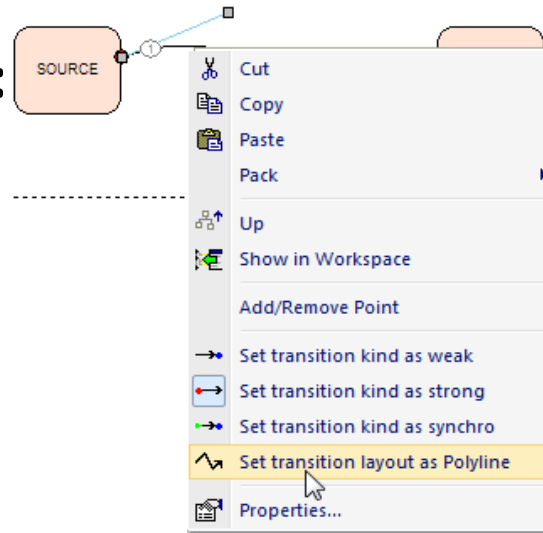


# Transition Editing: Modifying Layout

Two layout types:

- Bezier
- Polyline

Contextual access:



Menu access:

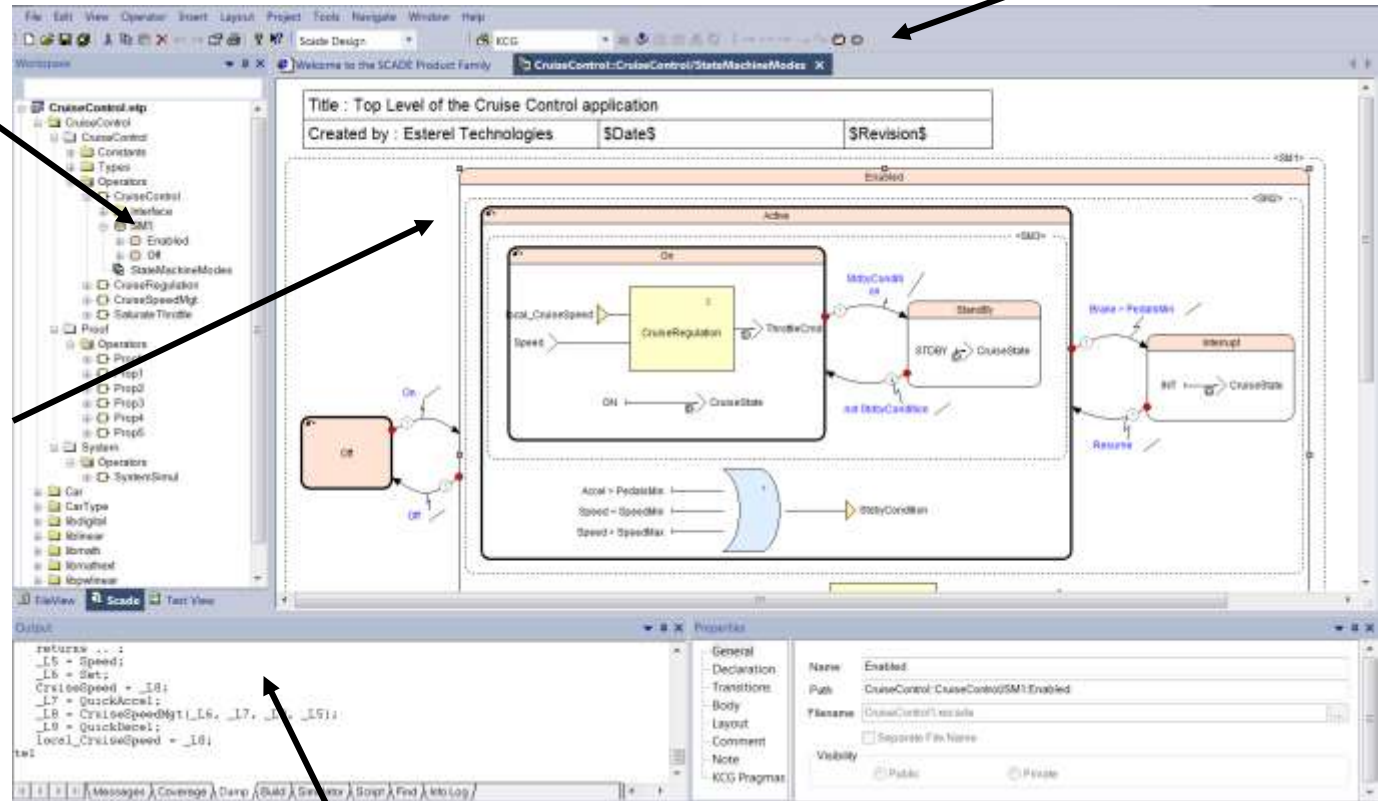


# State Machine in Scade View

State Machine Toolbar

State Machine in  
Scade View

State Machine in  
diagrams



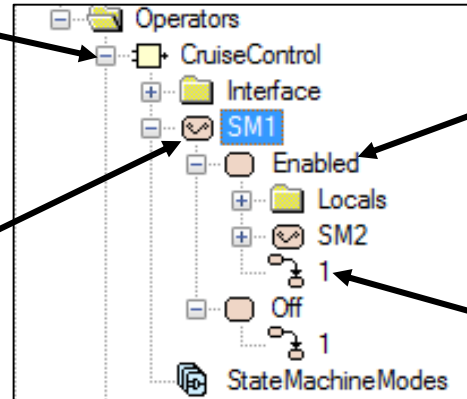
Dump tab:  
State Machine textual representation

# State Machine in Scade View

Node containing a State Machine

SM Name:

- Must be unique in the node
- Can be set and modified in the State Machine Properties Window



State Machine State

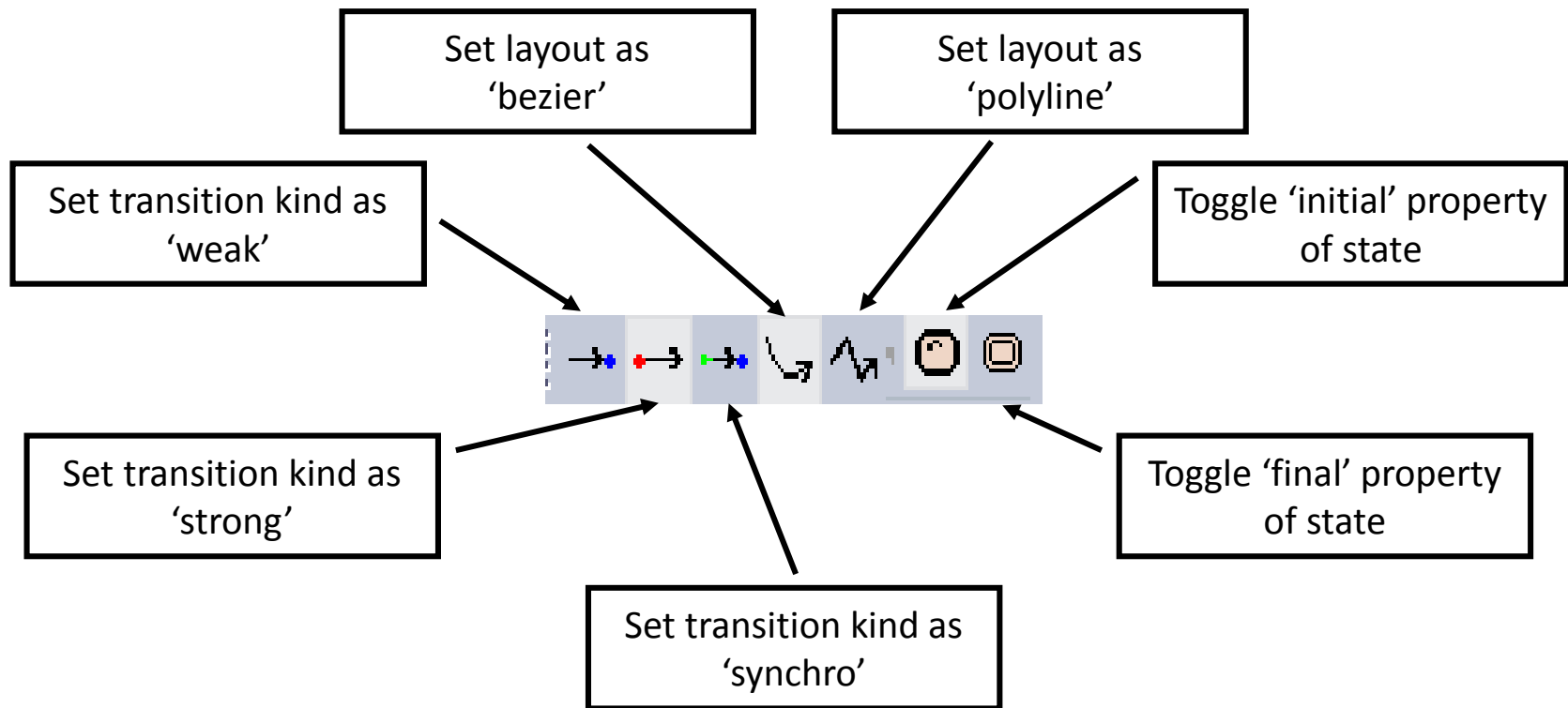
State Machine Transition

The State Machine Properties Window for 'SM1' is shown. It has a left sidebar with tabs: 'General' (selected), 'Comment', 'Note', 'KCG pragmas', and 'Traceability'. The main area contains the following fields:

- Name: SM1
- Path: CruiseControl::CruiseControl/SM1:
- Filename: CruiseControl1.xscade
- ☐ Separate File Name
- Visibility: ☐ Public ☐ Private

# State Machine Toolbar

Shortcuts to some properties and settings:





# Exercise 6

## Objective

Create a State Machine

## Requirement

Create the `OnOff` operator in a new project

**Project Name:** `SSM_Exercises`

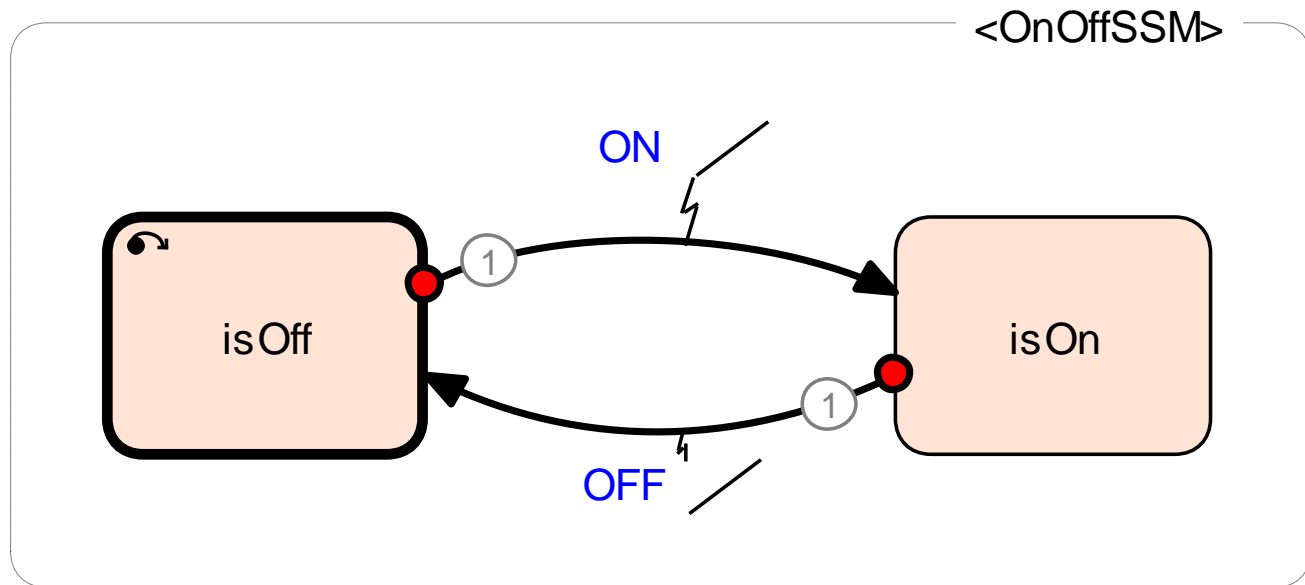
**Operator Name:** `OnOff`

Name	Kind	Type
ON	input	bool
OFF	input	bool

Initial state: button is in `isOFF` state

Behavior: if `ON` is true and `OnOff` is in `isOFF` state, `OnOff` goes to `isON` state. If `OFF` is true and `OnOff` is in `isON` state, `OnOff` goes to `isOFF` state

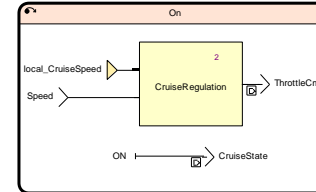
## Exercise 6: Solution



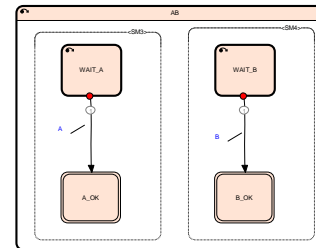
# State Machine State Activities

A state, "similar" to a SCADE operator, can contain:

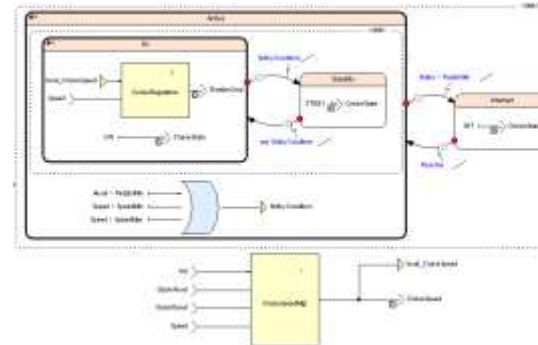
- Graphical data flow equations



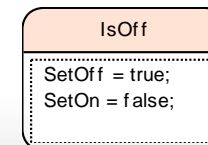
- Others State Machines



- And both



Can also be converted to a textual state



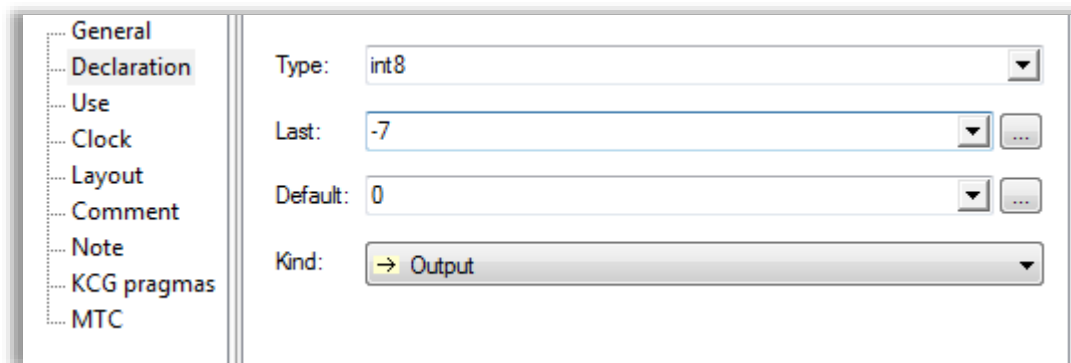
# Data Flow Inside State

Equations are calculated only when the state is active.

Each declared flow (local or output) must have exactly one definition for each cycle where its scope is active.

What happens when a definition is missing in a given state?

- Users can set either a **default value** or **maintain the last value of the flow**



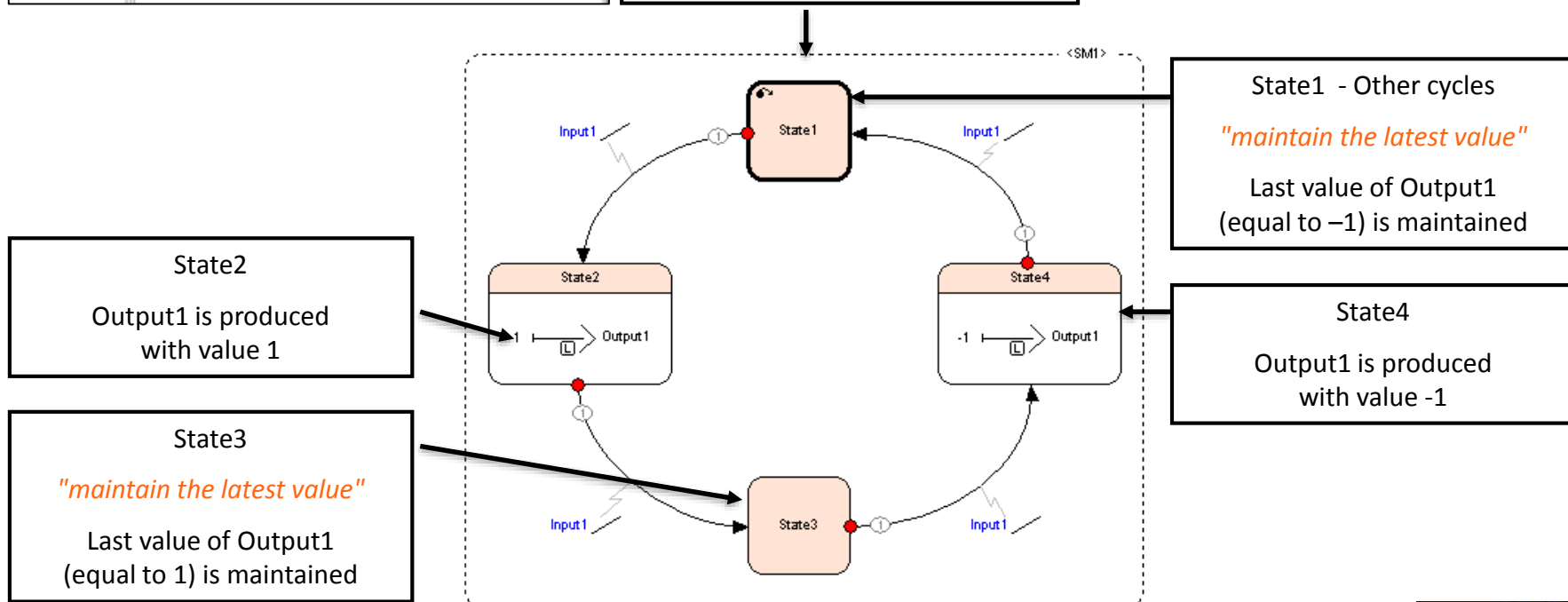
The screenshot shows the 'Declaration' tab of a dialog box in ANSYS StateFlow. The left sidebar contains a tree view with the following items: General, Declaration (selected), Use, Clock, Layout, Comment, Note, KCG pragmas, and MTC. The main area of the dialog has four fields: 'Type' is a dropdown menu set to 'int8'; 'Last' is a text field containing '-7' with a dropdown arrow and an ellipsis button; 'Default' is a text field containing '0' with a dropdown arrow and an ellipsis button; and 'Kind' is a dropdown menu set to '→ Output'.

# Data Flow Inside State

Maintain the latest value

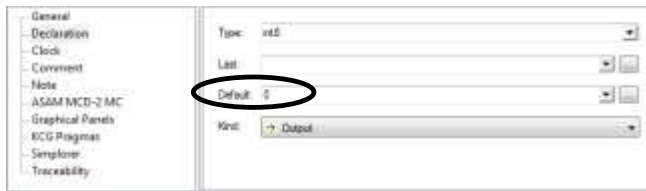


Initial State1 - First cycle  
*"maintain the latest value"*  
Output1 not yet produced.  
then set to 0 (initial last value)



# Data Flow Inside State

## Produce a default value



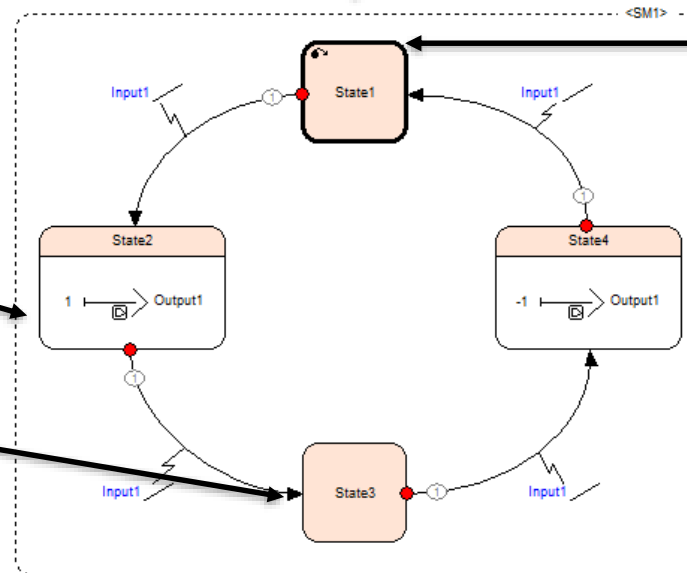
Initial State1 - First cycle

*"Produce a default value"*

Output1 is produced  
with default value 0

State2  
Output1 is produced  
with value 1

State3  
*"Produce a default value"*  
Output1 is produced  
with default value 0



State1 - Other cycles

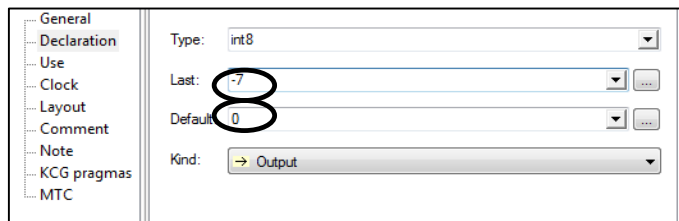
*"Produce a default value"*

Output1 is produced  
with default value 0

State4  
Output1 is produced  
with value -1

# Data Flow Inside State

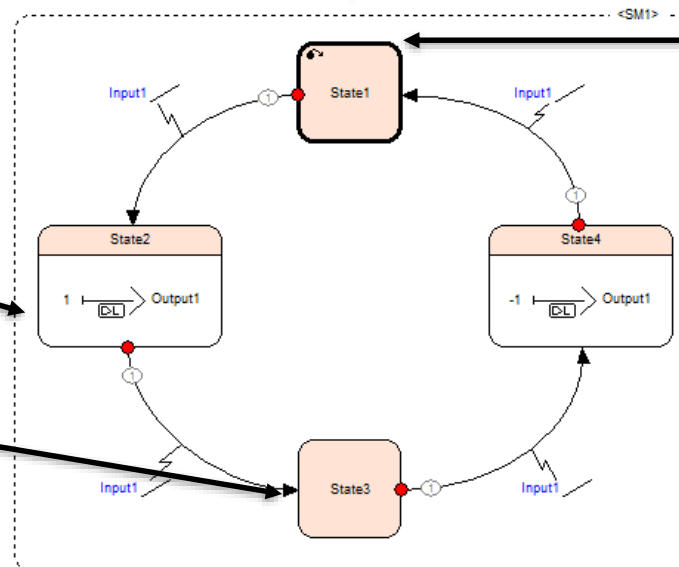
If both Last and Default are defined, Default has a higher priority than Last one



Initial State1 - First cycle  
*"Produce a default value"*  
Output1 is produced with default value 0

State2  
Output1 is produced with value 1

State3  
*"Produce a default value"*  
Output1 is produced with default value 0

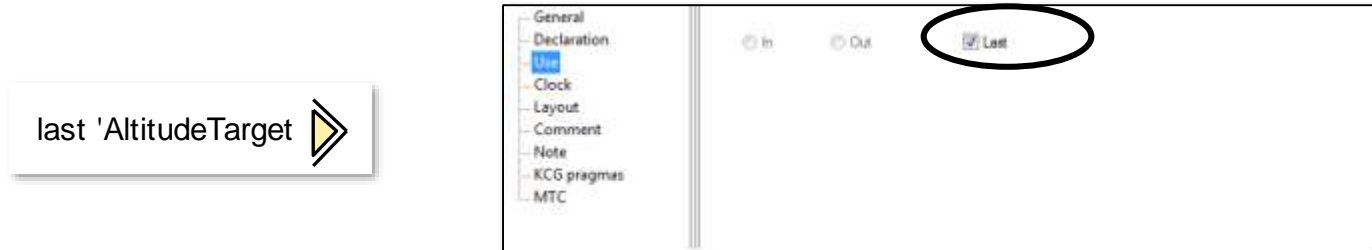


State1 - Other cycles  
*"Produce a default value"*  
Output1 is produced with default value 0

State4  
Output1 is produced with value -1

# Create a Last Primitive

Drag & drop the variable and click Last in the “Use” tab



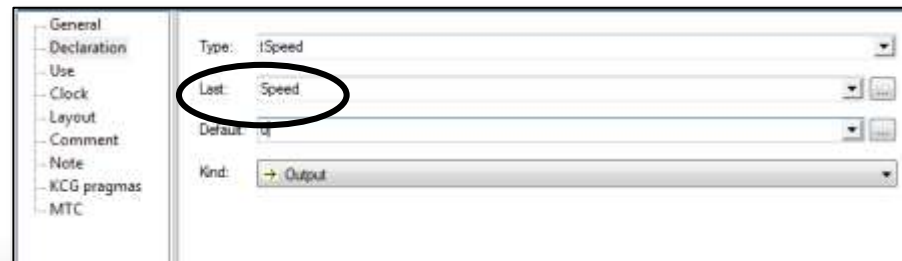
Use a textual expression of the form: last'VarName (Note the single quote in this expression)



last 'AltitudeTarget

A “last” variable must be defined on the first cycle.

Specify this initial value in the “Last” text box in the variable declaration tab

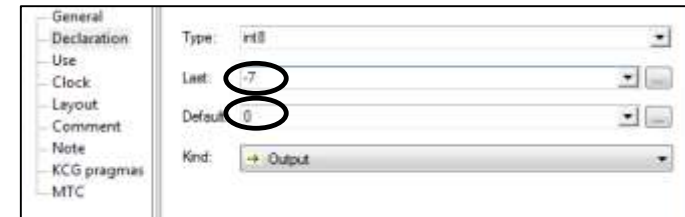




# Data Flow Inside State

## Objective:

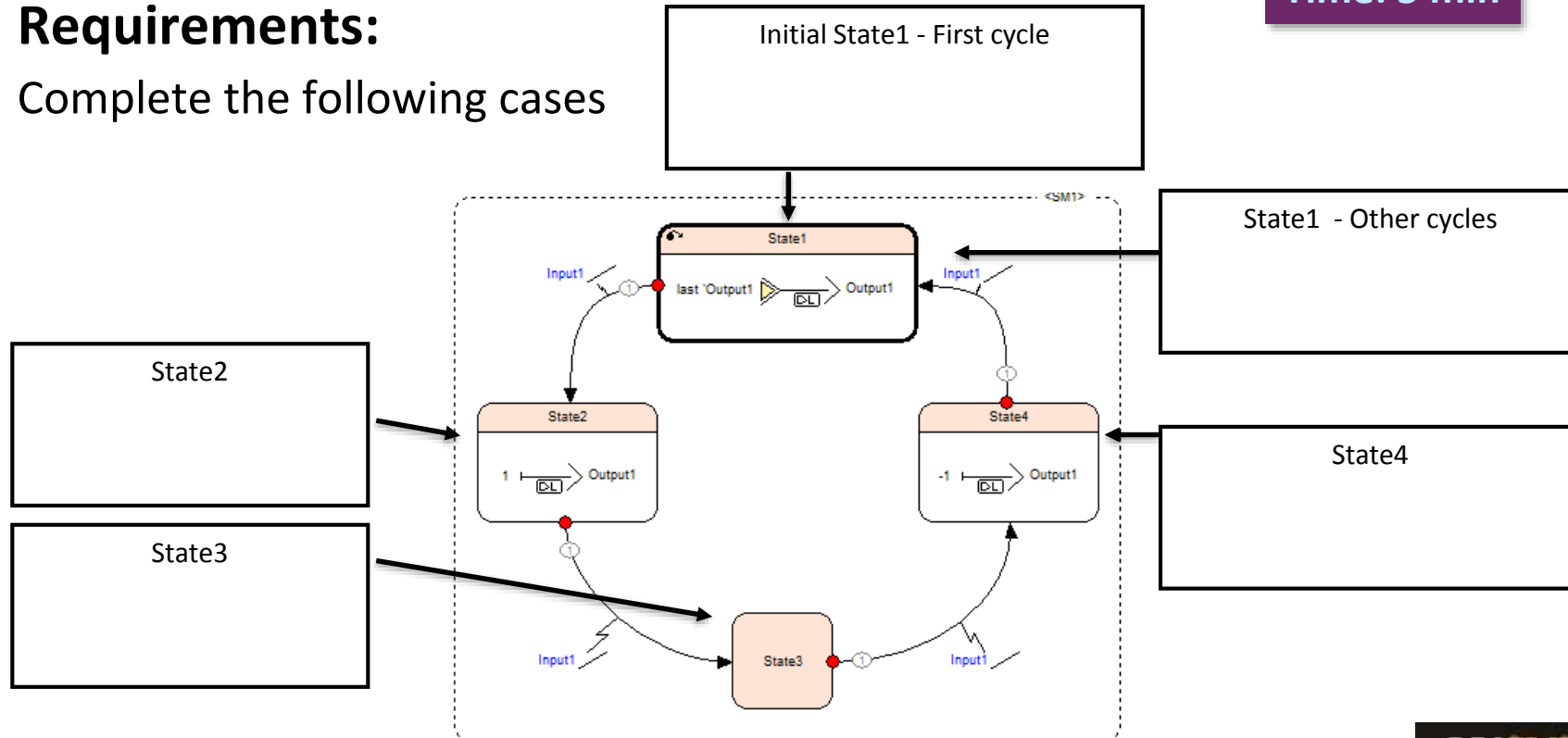
Define the produced value for Output1



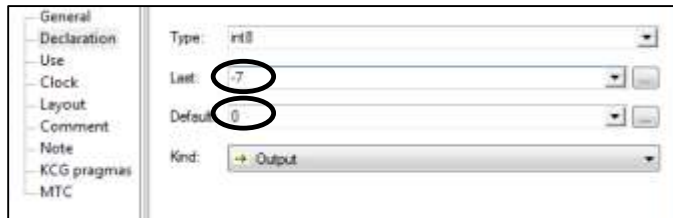
## Requirements:

Complete the following cases

Time: 5 min



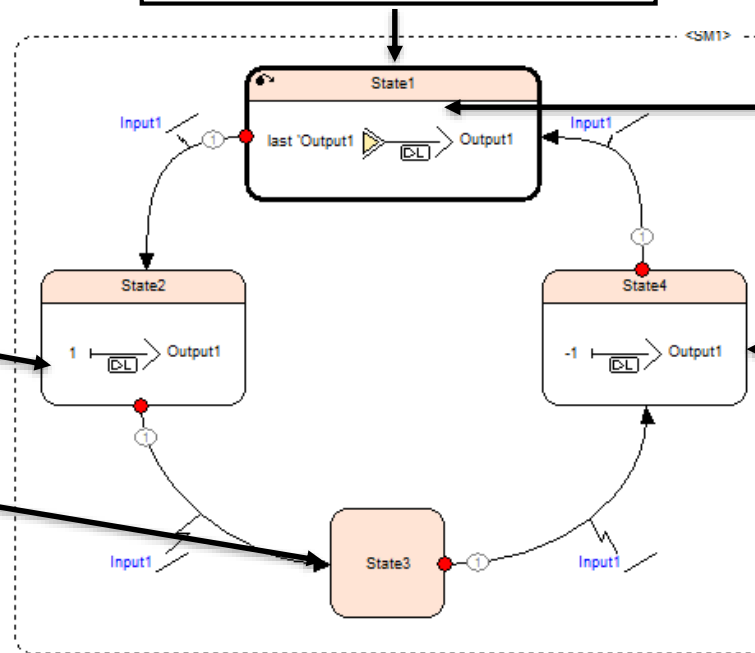
# Solution



Initial State1 - First cycle

*"maintain the latest value"*

Output1 not yet produced.  
then set to -7 (initial last value)



State1 - Other cycles

*"maintain the latest value"*

Last value of Output1  
(equal to -1) is maintained

State2

Output1 is produced  
with value 1

State3

*""Produce a default value"*

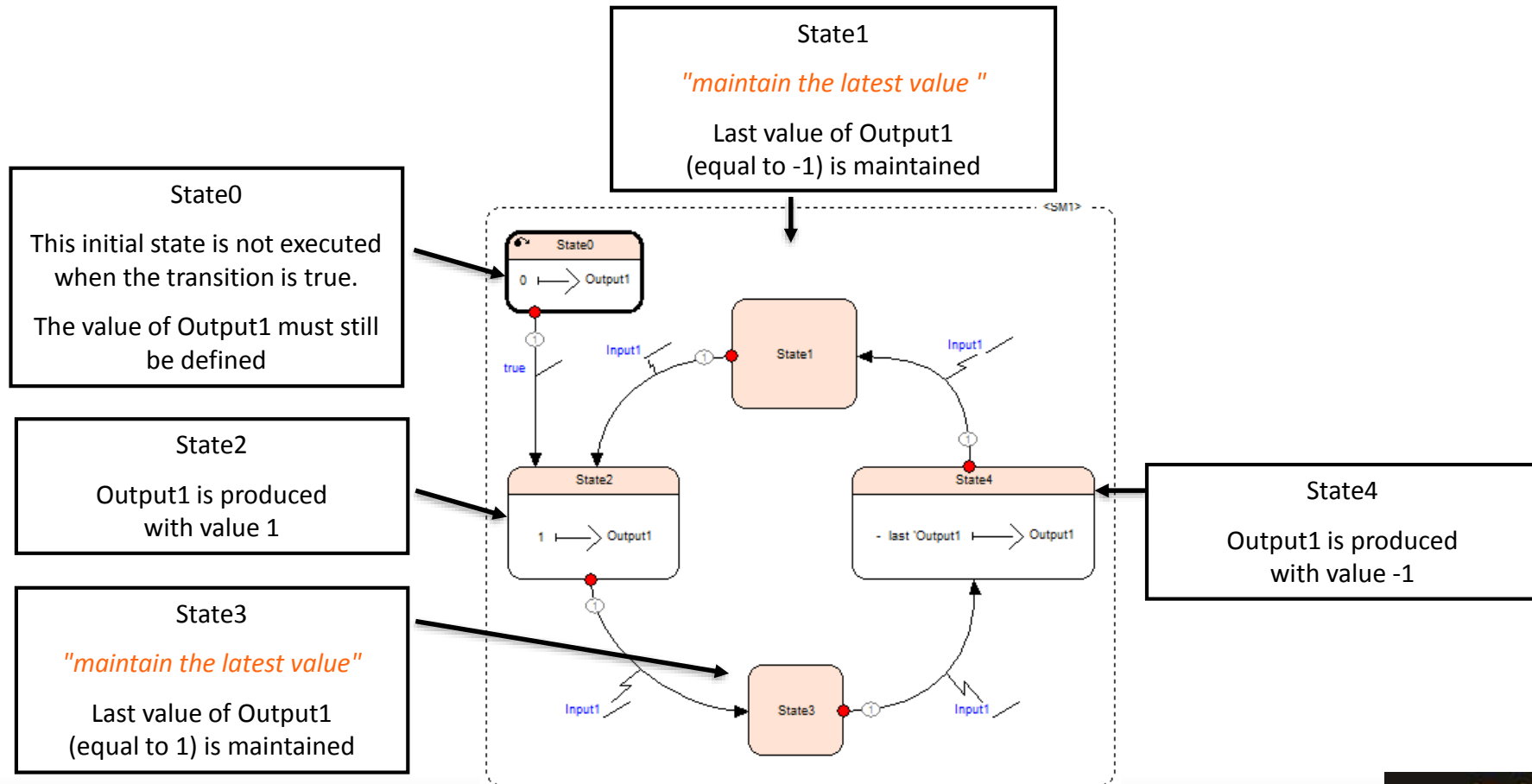
Output1 is produced  
with default value 0

State4

Output1 is produced  
with value -1

# Data Flow Inside State

If the variable is produced in the initial state but no default and last value is defined: the last behavior is kept, the variable must have a definition in every state accessible in the first cycle



# Exercises

OPTIONAL

# Exercise 7

## Objective

Observe the behavior of State Machine

## Requirement

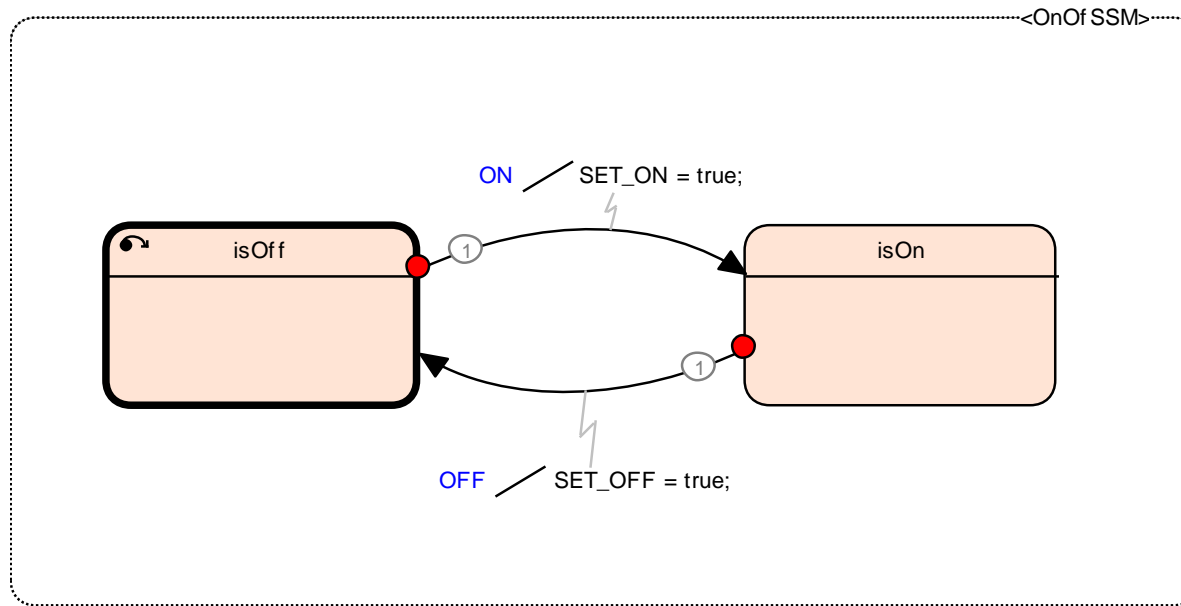
Add 2 output to the OnOff operator

Name	Kind	Type
ON	input	bool
OFF	input	bool
SET_ON	output	bool
SET_OFF	output	bool

The outputs SET\_ON and SET\_OFF shall be true only when transition is fired (no data flow in state)

***Use actions on transition and default values***

# Exercise 7: Solution



- Upon transition firing, SET\_ON and SET\_OFF are defined as true
- When no definition is activated the outputs' **default value** is false

# Exercise 8

## Objective

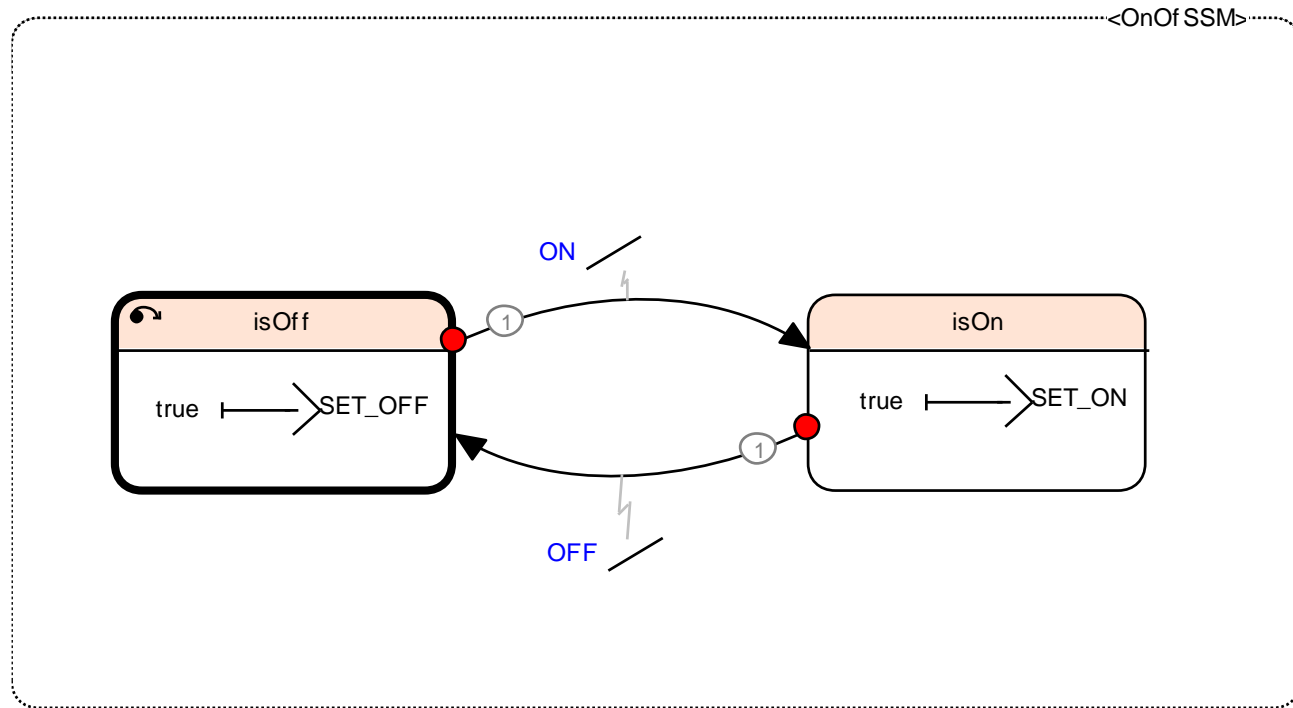
Observe the behavior of State Machine

## Requirement

- The SET\_ON output shall be true at all times when state ON is active
- The SET\_OFF output shall be true at all times when state OFF is active

***Use dataflow inside states and default values***

## Exercise 8: Solution



- Upon transition firing, states **IsOn** resp. **IsOff** are activated
- The operator outputs **SET\_ON** and **SET\_OFF** are continuously updated
- When no definition is activated the outputs' default value is false



# Exercise 9

## Objective

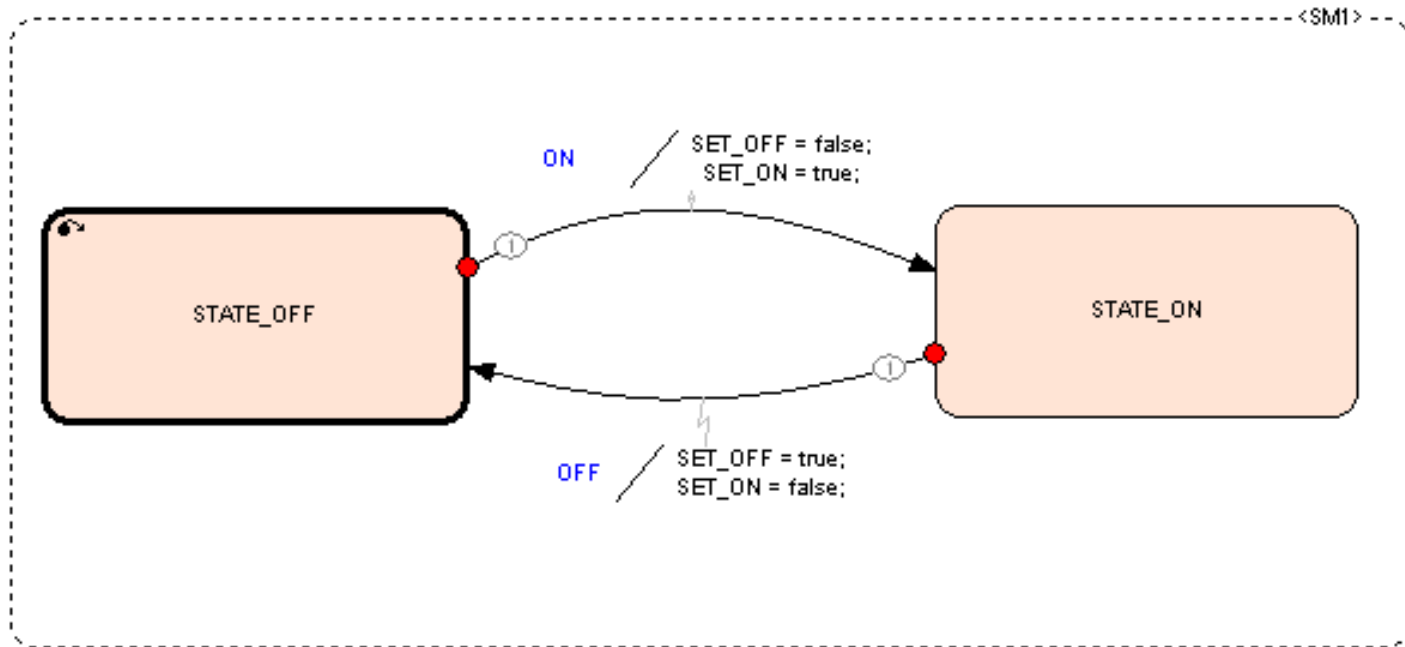
Observe the behavior of State Machine

## Requirement

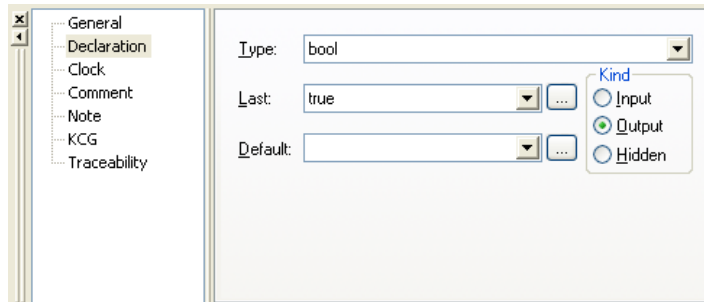
- The SET\_ON output shall be true at all times when isON state is active
- The SET\_OFF output shall be true at all times when isOFF state is active
- No data flow in state

***Use transition actions and last values***

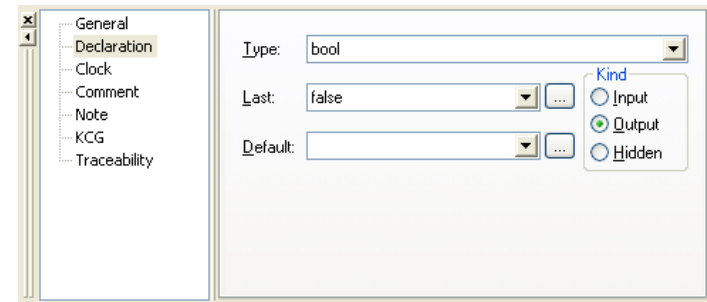
# Exercise 9: Solution



SET\_OFF properties window



SET\_ON properties window



# Exercise 10

## Objective

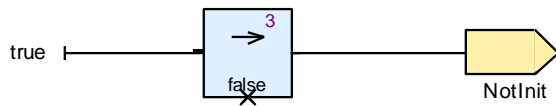
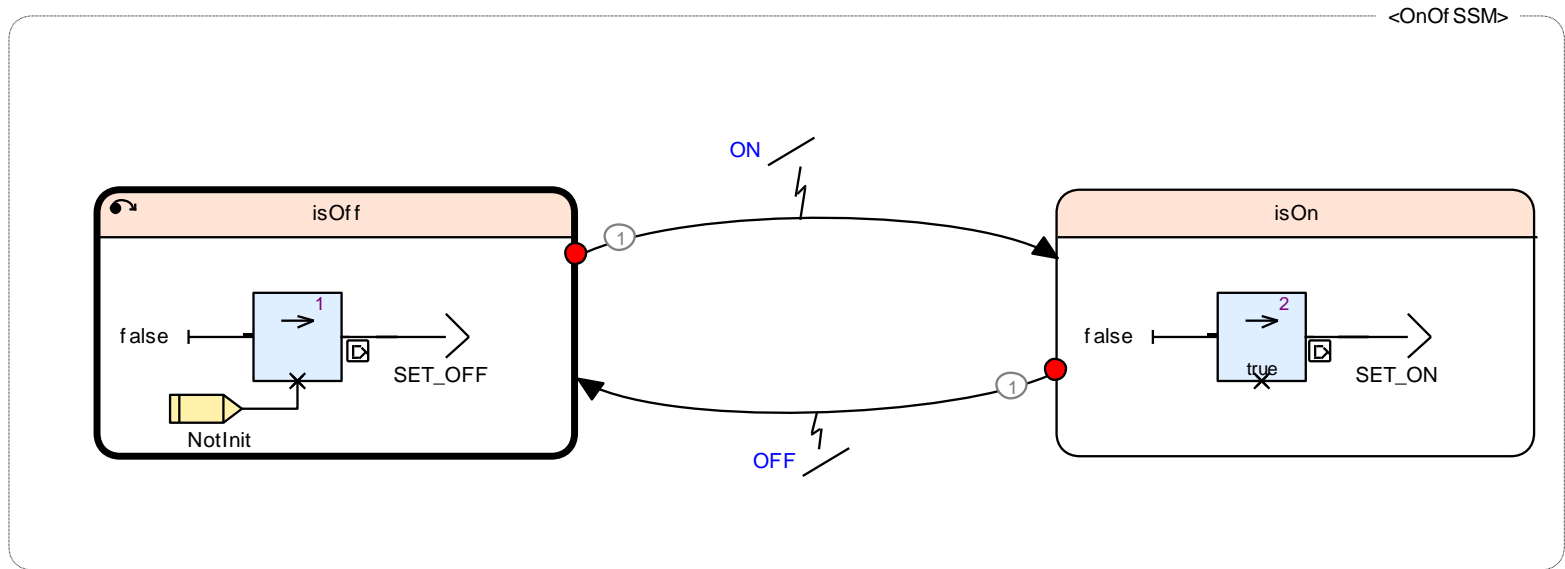
Observe the behavior of State Machine

## Requirement

SET\_ON and SET\_OFF true only when transition is fired

***Use dataflow inside state and the init operator***

# Exercise 10: Solution

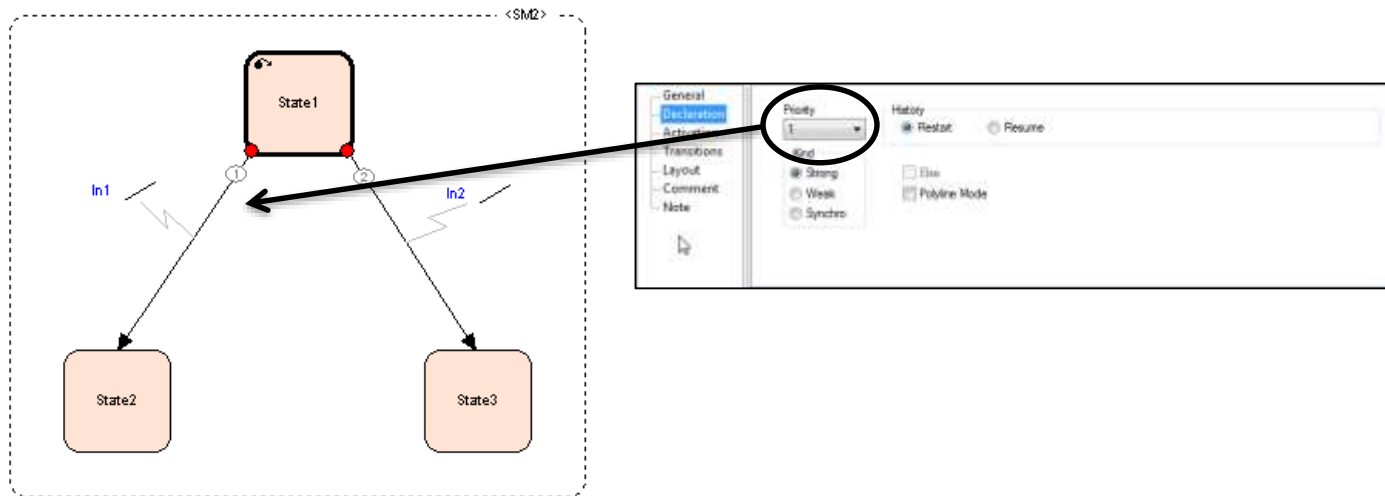


When no definition is activated the outputs' **default value** is false

# Transition Priorities

When conditions of several transitions starting from the same active state are true, only the one with the highest priority (1 is the highest) is fired

- The priority is mandatory to ensure determinism
- A default value is automatically added by the SCADE Suite Editor, but it can be manually changed by the designer



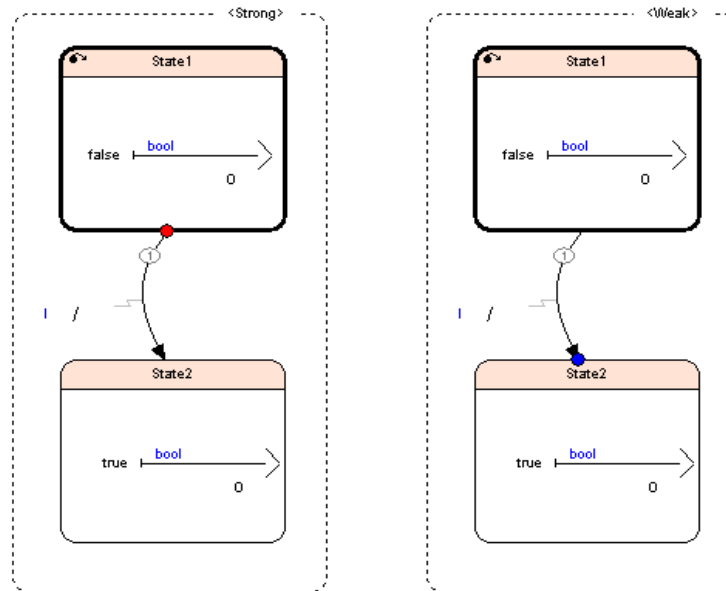
# Transition Preemption Types

## Three types of transition

- Strong: preempt the source state and set the target state as active, beginning its internal activity
- Weak delayed: let the source state finish its work during the cycle the transition is triggered. Target state activity is started at the next cycle
- Synchro: triggered on special condition (detailed after). Target state activity is started at next cycle (like weak)



# Strong and Weak Preemption



I	false	true	-
STRONG	O = false	O = true	O = true
WEAK DELAYED	O = false	O = false	O = true

# Exercise 11

## Objective

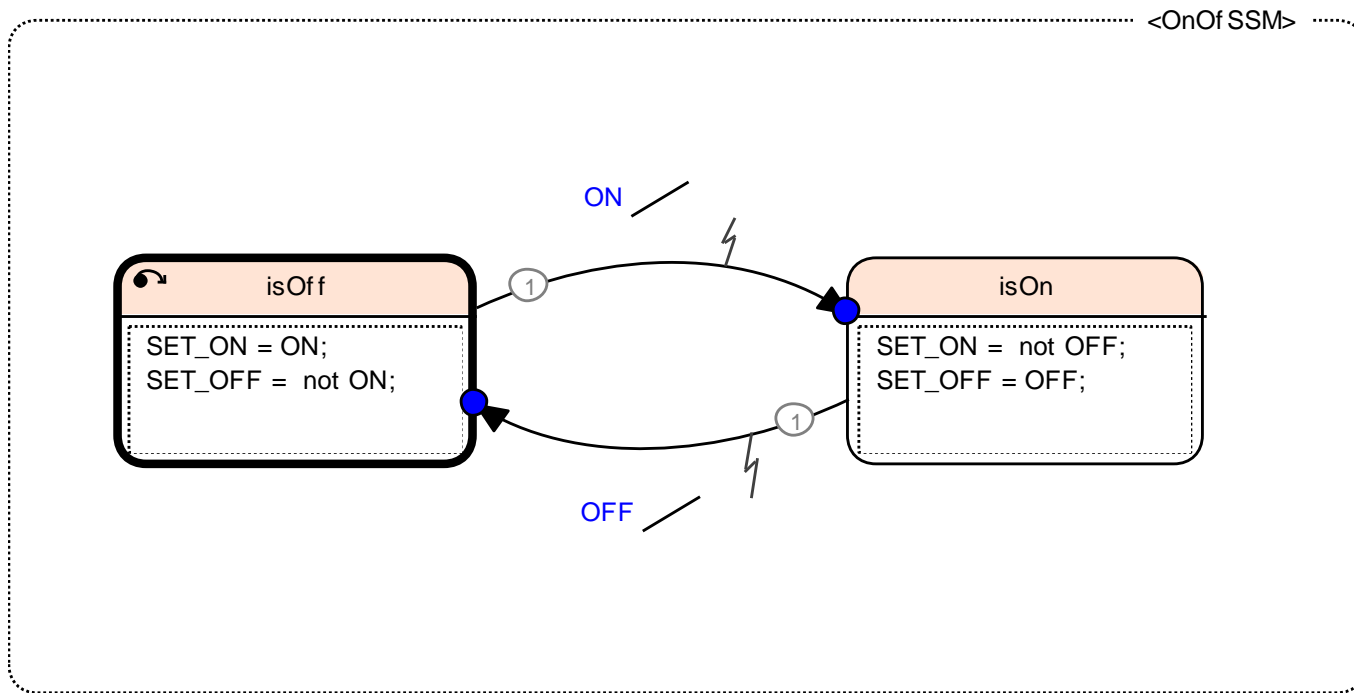
Using weak preemption transitions

## Requirement

- The outputs `SET_ON` shall be true immediately when the input `ON` is true, and stay true until the input `OFF` is true
- The outputs `SET_OFF` shall be true immediately when the input `OFF` is true, and stay true until the input `ON` is true



# Exercise 11: Solution



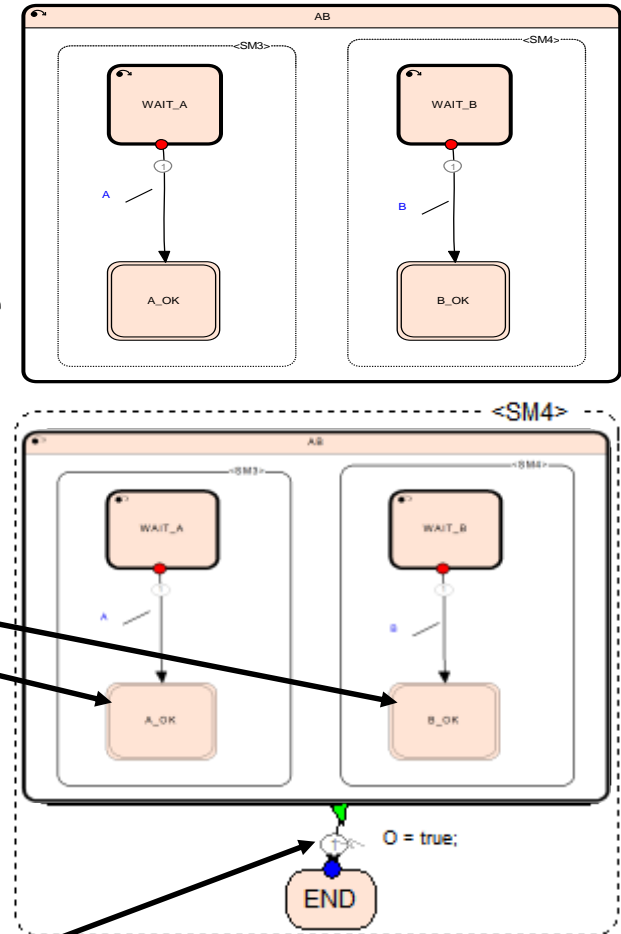
# Synchro Transition

Synchro transitions start with a green triangle

A synchro transition is fireable when all the final states in the source state are activated

Transitions can be defined out a final state

An effect action can be specified

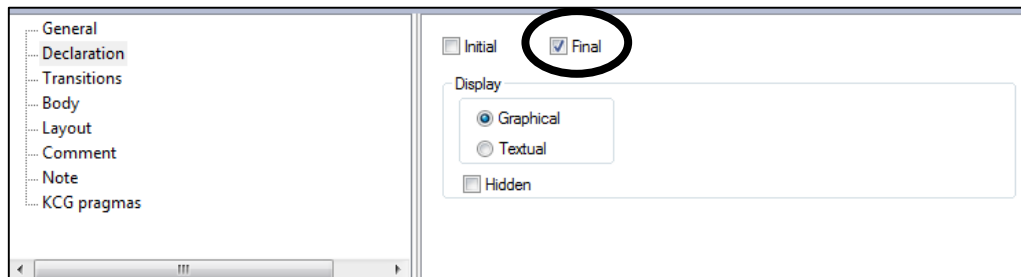
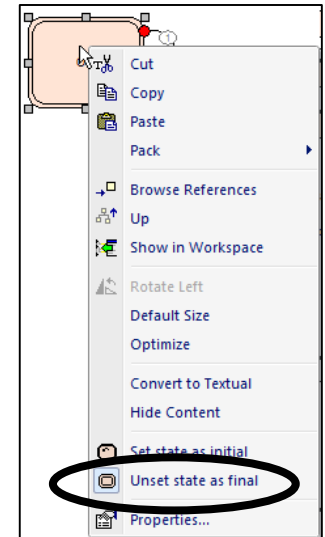


# Final State

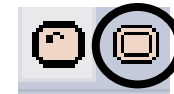
Graphically depicted with a doubled outline  
Used to mark a state for synchronization

How to set a state as final ?

- From its contextual menu: “Set/Unset state as final”
- From is Properties window: “Final” checkbox



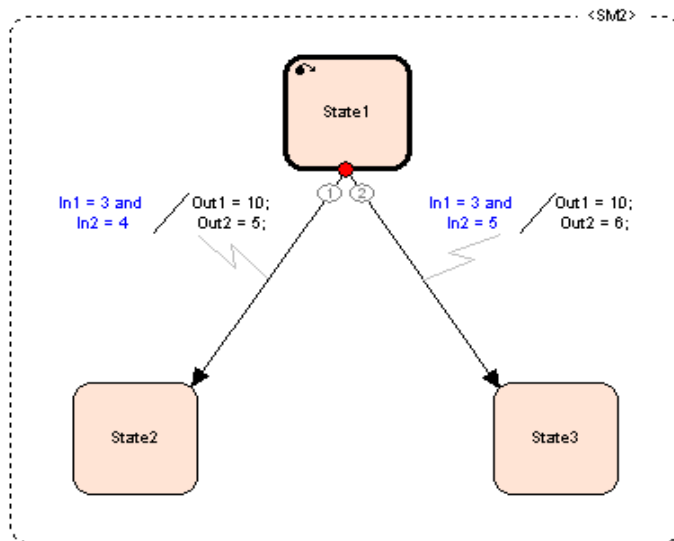
- From the State Machine toolbar dedicated button



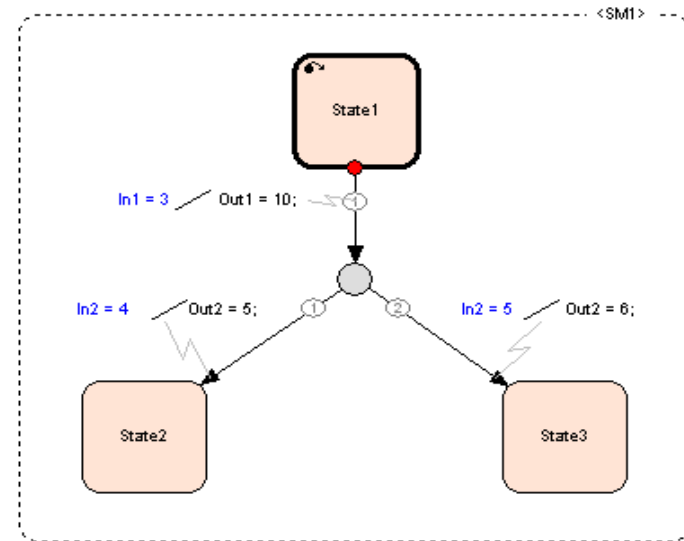
# Fork Transition

Fork transitions merge transitions with a common part in the condition:

- Common actions can also be merged on the common branch
- Actions are executed only on the fired transition
- The transition is only fired if a new destination state can be reached
- The else branch is not mandatory

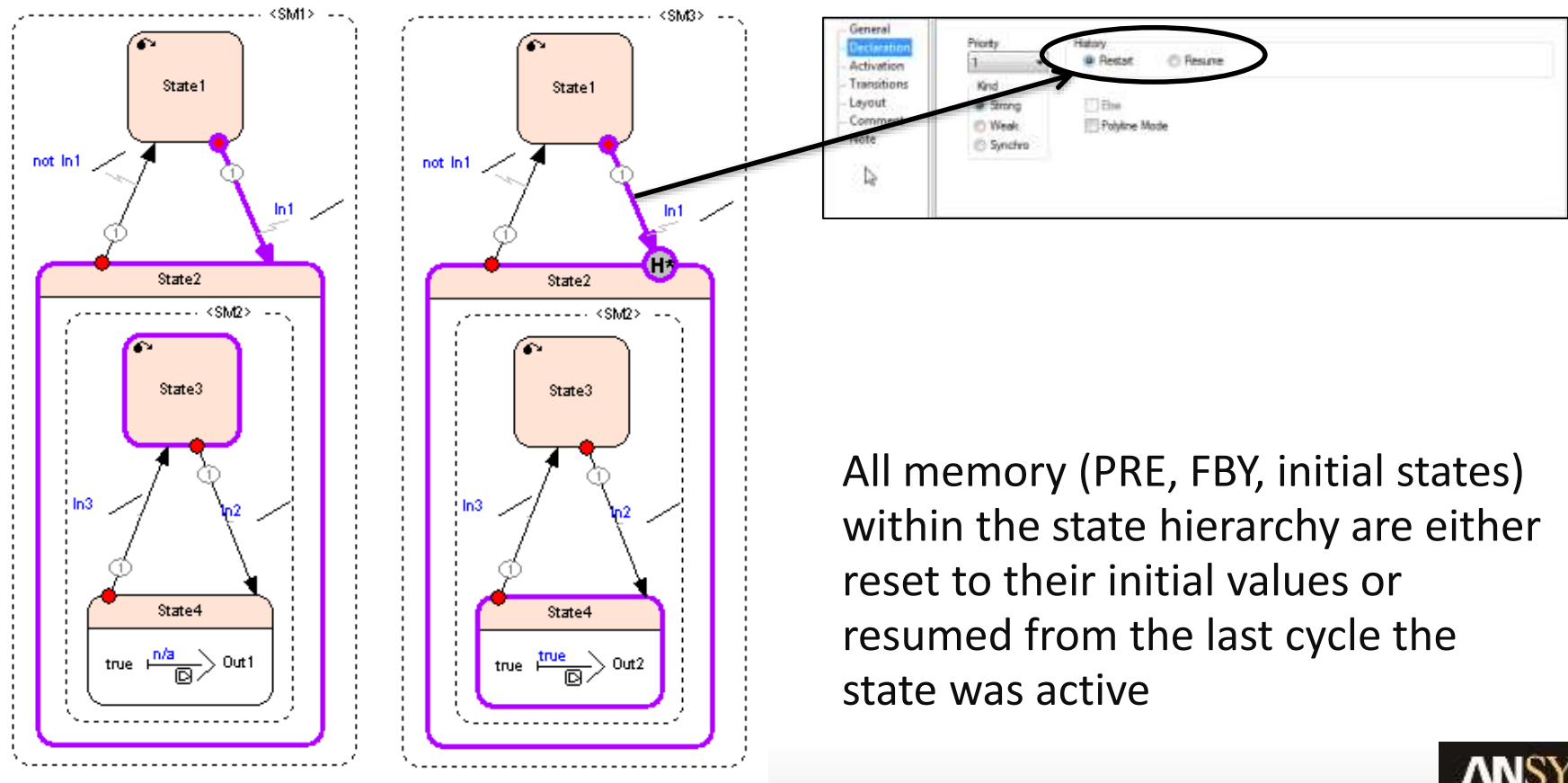


=



# History Transition

By default, a transition restarts its target state  
But it is possible to resume the target state

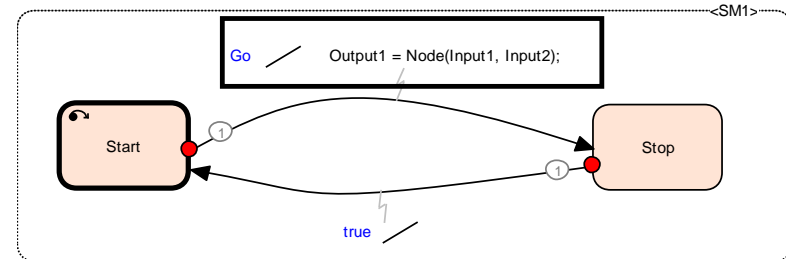
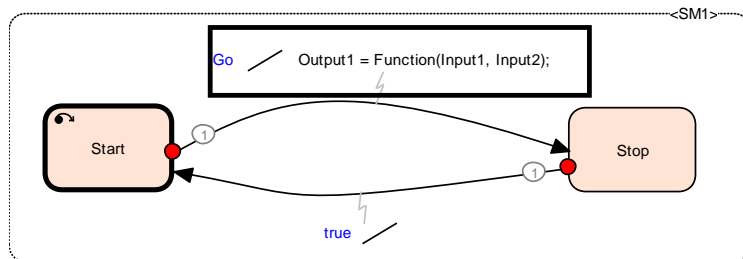


All memory (PRE, FBY, initial states) within the state hierarchy are either reset to their initial values or resumed from the last cycle the state was active

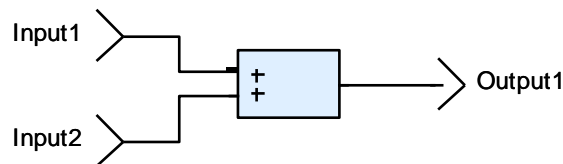
# Calling Operators in Actions

Functions and nodes are accepted as actions in transitions. Memories of nodes are restarted each time the transition is triggered.

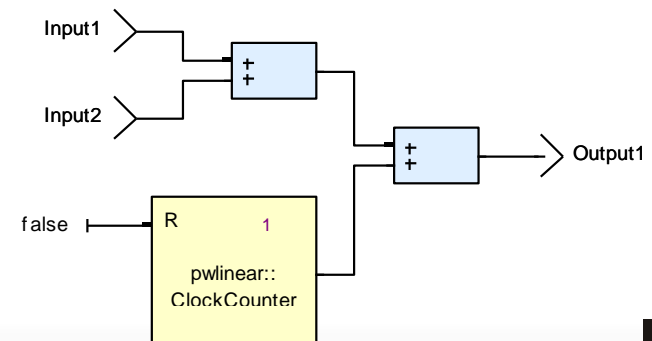
These following operators are equivalent:



## Function



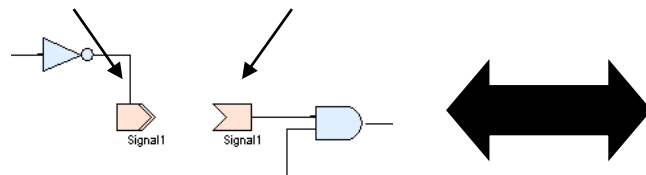
## Node



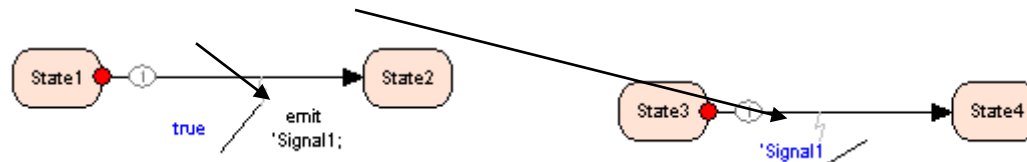
# Signals

## Signal usage type:

- Emit or Present



- Emit or Present



## Create signals:

- Using the New Signal button
- Using the contextual menu Insert > Signal in Scade view



# Signal Emission

Emitted signals are visible inside the context (operator, branch, state) in which the signal is defined.

A signal can be emitted several times in a given cycle.

In a given cycle, if a signal is emitted at least once then it is present in that cycle.



# Communication and Scope

Inputs and Outputs are visible throughout the operator they are defined in.

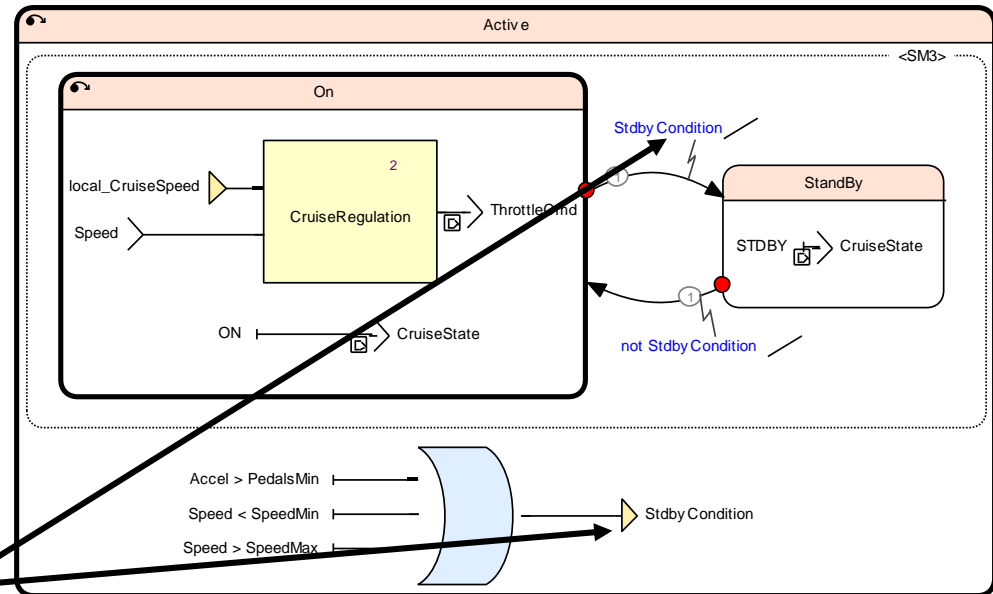
Signals and Locals are visible throughout the scope they are defined in:

- Operator

- State



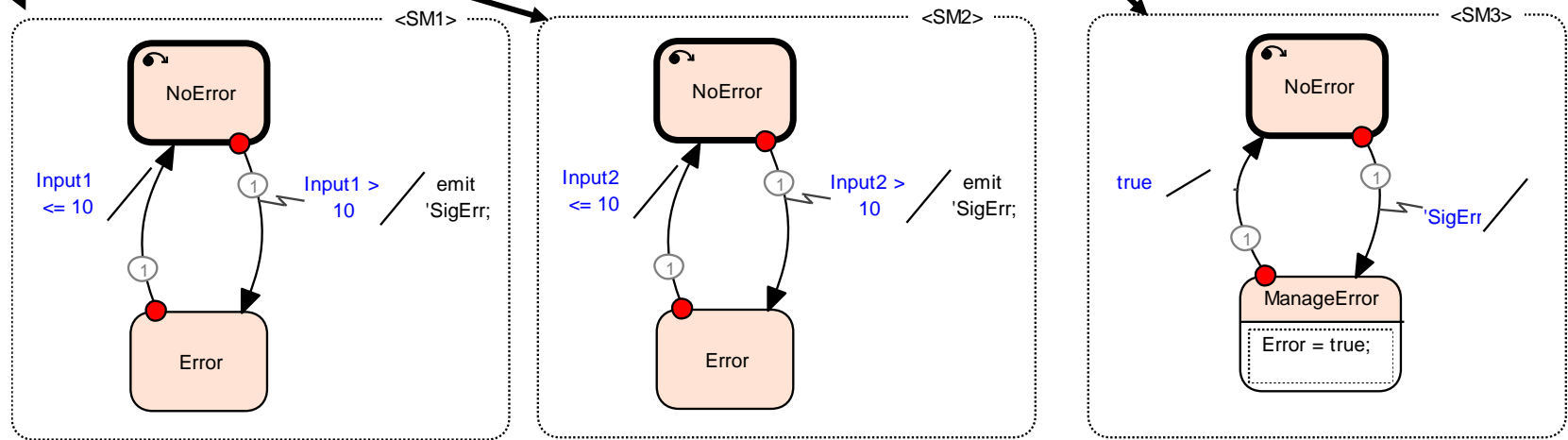
StdbbyCondition is only defined in the Active state and its hierarchy



# Communication and Scope

Signals are useful to catch a specific situation in several SMs:

- A signal is emitted in several parallel SMs when conditions are met
- A parallel State Machine waits for the presence of the signal to respond to the event



# State Machine: Causality Loops

Causality loops are:

- Cyclic dependencies of flow computation
- Or a mix of State/Transition execution and flow computation

Causality loops are detected by KCG and so by the SCADE Suite Checker

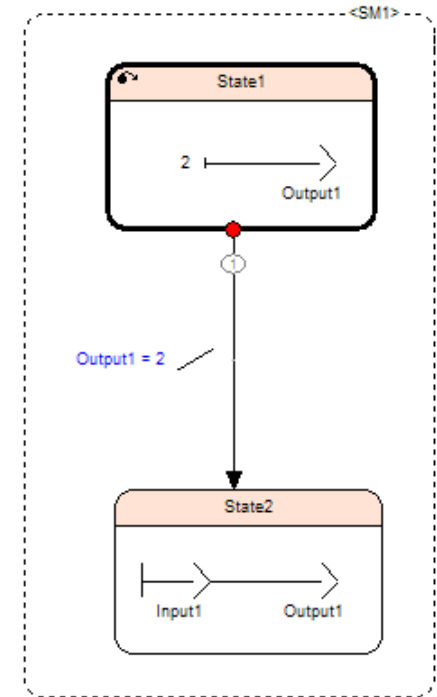
The error message indicates the list of flows, states, and transitions linked to the cyclic dependency

# State Machine: Causality Loops

Transition from State1 to State2 is triggered if  
Output1 = 2

Output1 is set to 2 by State1 exactly when the  
transition is not fired

Hence, transition firing depends on State1  
execution, which depends on the transition



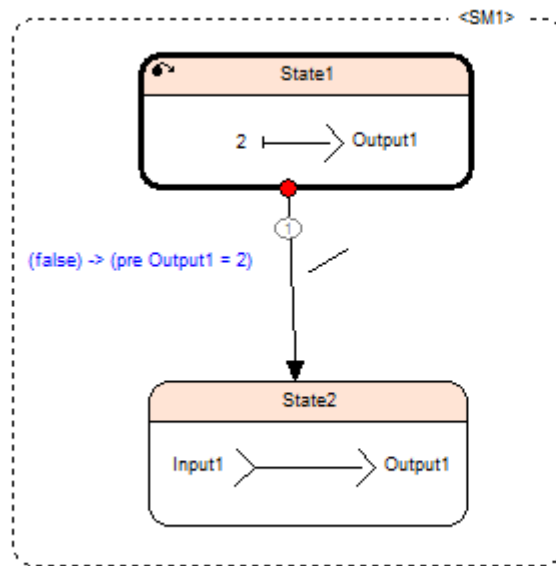
The Model Checker checks if such problems exist:

1 error(s) detected - 0 warning(s) detected		
Category	Code	Message
Semantic Error	ERR_400	<b>Causality at <a href="#">Package1::CyclicSSM/SM1:State1:</a></b> the definition of the strong guards of state State1 depends on flow Output1 ; ( <a href="#">Package1::CyclicSSM/SM1:State2:Output1=</a> ) the definition of flow Output1 depends on the state of automaton SM1 via the control context ; ( <a href="#">Package1::CyclicSSM/SM1:State1:</a> ) the definition of the state of automaton SM1 depends on the strong guards of state State1 ;

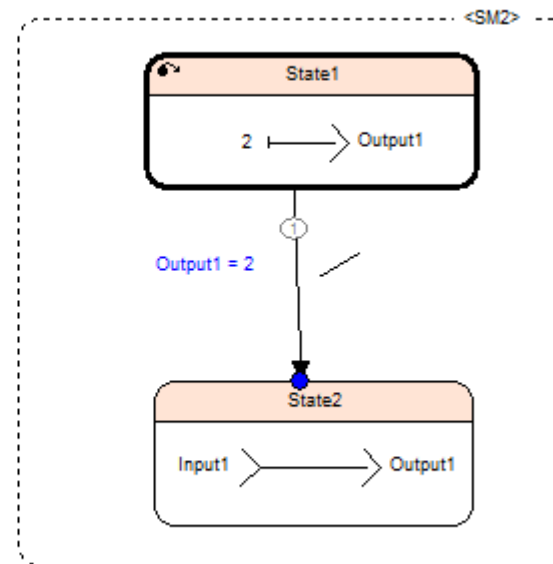
# Causality Loops

Causality loops can be solved by the introduction of additional memory to break the cycle

**Memory on one of the flows using pre or last**



**Memory on a transition using the weak delayed transition**



# SCADE Suite Checker

Semantic errors may be introduced while designing models and must be corrected before users can ever use a model for simulation or code generation

The SCADE Suite Checker analyses the selected items based on the active configuration set in the Code Generator toolbar:



Note: see details related to these options in the code generation section

# SCADE Suite Checker

Select an operator from the Scade view and *Right-click* > *Check*  
or  
Select Check on the Code Generator toolbar



# Lab 5 (1/2)

Use Lab Support p.65-74

## Objective:

Implement the `CruiseControl` (CC) operator:

- Use of states machines, mixed data and control flows

## Requirements:

Time: 35 min

Close all SCADE suite sessions and load a copy of “*Prerequisites\Lab 5\CruiseControl\CruiseControl.etp*”

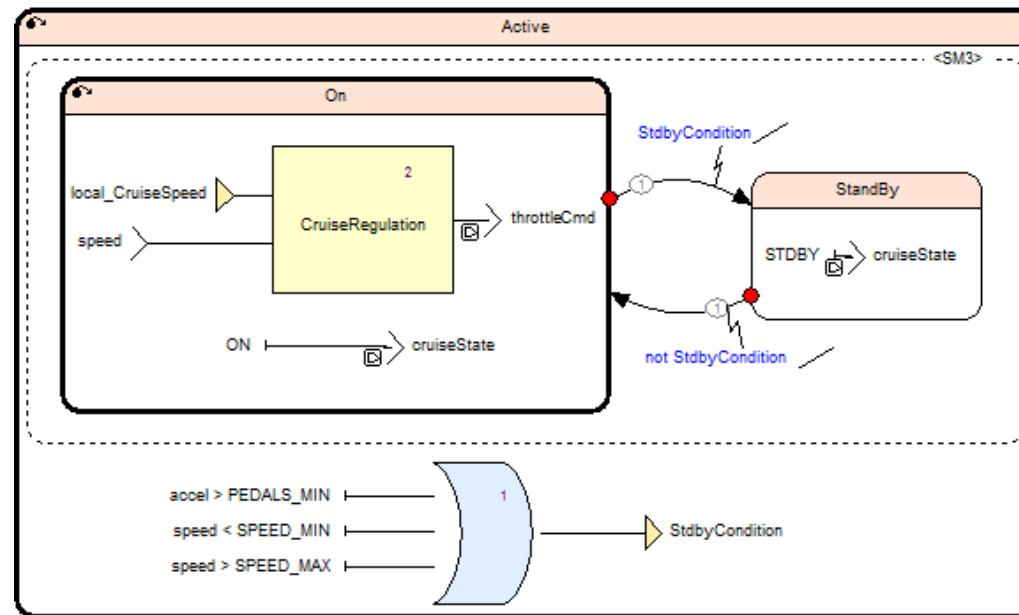
The `CruiseControl` operator calls:

- The `CruiseRegulation` operator computes the throttle command
- The `CruiseSpeedMgt` operator manages the value of the cruise speed according to the driver commands when the Cruise Control is enabled



# Lab 5 (2/2)

Modelling of the system when the Cruise Control is active (ON, STDBY states) in the Active state



Set default value of output

Output	Default Value
throttleCmd	accel
cruiseState	OFF

# Lab 6 (1/5)

Use Lab Support p.76-85

## Objective:

Simulate the `SystemSimul` operator which integrates both subsystems, the car and the `CruiseControl` operators to model the complete system

## Requirements:

Time: 30 min

The `SystemSimul` operator calls:

- The `CruiseControl` operator developed in the previous exercise
- The `CarModel` operator from the `Car` library

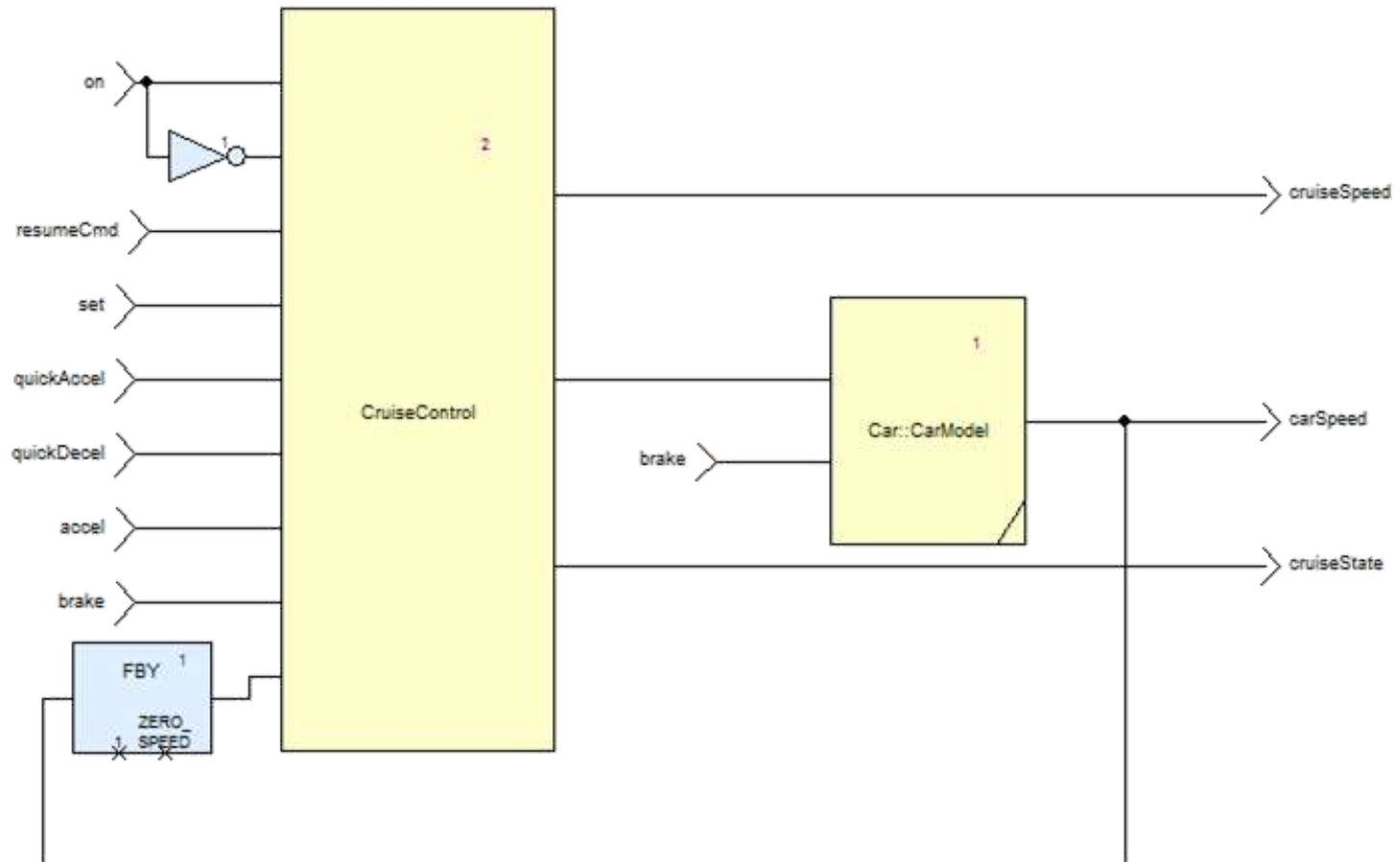
## Lab 6 (2/5)

Create the interface of `SystemSimul` operator

Name	Kind	Type
on	Input	bool
resumeCmd	Input	bool
Set	Input	bool
quickAccel	Input	bool
quickDecel	Input	bool
accel	Input	CarType::tPercent
brake	Input	CarType::tPercent
cruiseSpeed	Output	CarType::tSpeed
cruiseState	Output	CruiseControl::tCruiseState
carSpeed	Output	float64

# Lab 6 (3/5)

Design the SystemSimul diagram



# Lab 6 (4/5)

Simulate the scenario 1

Cycle duration	Input values	Output values to check
1	<code>on = false</code> <code>resumeCmd = false</code> <code>set = false</code> <code>quickAccel = false</code> <code>quickDecel = false</code> <code>accel = 0.0</code> <code>brake = 0.0</code>	-
2		<code>cruiseSpeed = CarSpeed = 0.0</code> <code>cruiseState = CruiseControl::OFF</code>
5	<code>Accel 50.0</code>	<code>cruiseState = cruiseControl::OFF</code>

Save the scenario1 (.sss format)

## Lab 6 (5/5)

Simulate the scenario 2 (after a reset of the simulation)

Cycle duration	Input values	Output values to check
1	-	-
1	accel = 100.0	-
1	on = true	-
1	-	cruiseState = CruiseControl::STDBY
12	brake = 10.0	cruiseState = CruiseControl::INT
1	Brake = 0.0 resumeCmd = true	cruiseState = CruiseControl::STDBY
1	accel = 1.0	cruiseState = CruiseControl::ON

Save the scenario 2 (.sss format)

You can reload and replay your scenarios

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Fax: +33 5 34 60 90 41

Submit questions to Technical Support: [scade-support@ansys.com](mailto:scade-support@ansys.com)

Contact one of our Sales representatives at: [scade-sales@ansys.com](mailto:scade-sales@ansys.com)

Direct general questions about Esterel Technologies to: [scade-info@ansys.com](mailto:scade-info@ansys.com)

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